

# CS 474

## Assignment 2

Name: Adharsh Kamath

UIN: 671259918

NetID: ak128

---

### Problem 1

#### **Soln:**

Let us use the symbol  $\psi$  to refer to the given formula.

$$\psi = (p \wedge (p \Rightarrow q)) \Rightarrow q$$

In order to show that  $\psi$  is valid, we can show that

$$\neg\psi = \neg((p \wedge (p \Rightarrow q)) \Rightarrow q)$$

is unsatisfiable. Rewriting the above formula:

$$\begin{aligned}\neg\psi &= \neg((p \wedge (p \Rightarrow q)) \Rightarrow q) \\ &\equiv \neg((p \wedge (\neg p \vee q)) \Rightarrow q) \\ &\equiv \neg(\neg(p \wedge (\neg p \vee q)) \vee q) \\ &\equiv (p \wedge (\neg p \vee q)) \wedge \neg q\end{aligned}$$

The last step is due to De Morgan's Law. We can now convert this to CNF, and construct a resolution refutation to show that it is unsatisfiable. To convert to CNF, we use the Tseitin transformation. We only need three new propositional variables,  $x_\psi, x_1, x_2$ , where  $x_\psi$  corresponds to  $\psi$ ,  $x_1$  corresponds to  $(\neg p \vee q)$  and  $x_2$  corresponds to  $(p \wedge x_1)$ . This gives us the following set of clauses:

$$\left\{ \begin{array}{l} \{\neg x_\psi\}, \\ \{\neg\neg x_\psi, x_2\}, \{\neg\neg x_\psi, \neg q\}, \{\neg x_\psi, \neg x_2, \neg\neg q\} \\ \{\neg x_2, p\}, \{\neg x_2, x_1\}, \{x_2, \neg p, \neg x_1\}, \\ \{x_1, \neg\neg p\}, \{x_1, \neg q\}, \{\neg x_1, \neg p, q\}, \end{array} \right\}$$

Simplifying the set by replacing  $\neg\neg p$  with  $p$  for all propositional variables, we get:

$$\left\{ \begin{array}{l} \{\neg x_\psi\}, \\ \{x_\psi, x_2\}, \{x_\psi, \neg q\}, \{\neg x_\psi, \neg x_2, q\} \\ \{\neg x_2, p\}, \{\neg x_2, x_1\}, \{x_2, \neg p, \neg x_1\}, \\ \{x_1, p\}, \{x_1, \neg q\}, \{\neg x_1, \neg p, q\}, \end{array} \right\}$$

We can now create a resolution refutation to show that this set is unsatisfiable:

1.  $\{\neg x_\psi\}$
2.  $\{x_\psi, x_2\}$
3.  $\{x_2\}$     Resolvent of 1 and 2
4.  $\{\neg x_2, p\}$
5.  $\{p\}$     Resolvent of 3 and 4
6.  $\{\neg x_1, \neg p, q\}$
7.  $\{\neg x_1, q\}$     Resolvent of 5 and 6
8.  $\{x_\psi, \neg q\}$
9.  $\{x_\psi, \neg x_1\}$     Resolvent of 7 and 8
10.  $\{\neg x_1\}$     Resolvent of 1 and 9
11.  $\{\neg x_2, x_1\}$
12.  $\{\neg x_2\}$     Resolvent of 3 and 11
13.  $\{\}$     Resolvent of 10 and 12

By creating this resolution refutation, we have shown that there is no valuation that can satisfy  $\neg\psi$ . Since  $\neg\psi$  is unsatisfiable, we can conclude that  $\psi$  is valid.

The above proof can be validated using the resolution tool provided as a part of the assignment. The initial set of clauses, and the resulting resolution proof object can be found at [https://github.com/adharshkamath/cs474-hw/tree/main/resolution\\_proofs/pl](https://github.com/adharshkamath/cs474-hw/tree/main/resolution_proofs/pl)

---

## Problem 2

### **Soln:**

We are given the formula:

$$\psi = (q \vee \neg r) \wedge (\neg p \vee r) \wedge (\neg q \vee r \vee p) \wedge (p \vee q \vee \neg q) \wedge (\neg r \vee q)$$

We can find resolvents in the following way: (Only resolutions that lead to clauses that do not already exist in the set of clauses are considered. Resolutions that result in "trivial" clauses are

also considered.)

1.  $\{q \vee \neg r\}$
2.  $\{\neg p \vee r\}$
3.  $\{q \vee \neg p\}$     Resolvent of 1 and 2
4.  $\{\neg q \vee r \vee p\}$
5.  $\{r \vee \neg r \vee p\}$     Resolvent of 1 and 4
6.  $\{p \vee q \vee \neg q\}$
7.  $\{p \vee q \vee \neg r\}$     Resolvent of 1 and 6
8.  $\{r \vee \neg q\}$     Resolvent of 2 and 4
9.  $\{r \vee \neg r\}$     Resolvent of 2 and 5
10.  $\{r \vee p \vee \neg p\}$     Resolvent of 2 and 5
11.  $\{r \vee q \vee \neg q\}$     Resolvent of 2 and 6
12.  $\{r \vee q \vee \neg r\}$     Resolvent of 2 and 7
12.  $\{p \vee q \vee \neg p\}$     Resolvent of 2 and 7
13.  $\{\}$

---

Problem 3

**Soln:**

---