

Answer of 1

In binary search, the sub-problem is searching for X in the search space of [lower. upper]. The problem is non-overlapping because we make the decision to search in either [lower, mid - 1] search space or [mid+1, upper] search space which will never repeat.

In fib(n), the sub-problem is calculating the fib(n-1) and fib(n-2) as $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$. This results in multiple overlapping calls as for fib(n) we need to find fib(n-1) and fib(n-2). Similarly, for fib(n-1), we need to calculate fib(n-2) and fib(n-3). Already we are calculating fib(n-2) twice. This is the overlapping subproblem.

Answer of 2

D	“”	“k”	“ka”	“kal”	“kale”
“”	0	1	2	3	4
“m”	1	1	2	3	4
“ma”	2	2	1	2	3
“map”	3	3	2	2	3
“mapl”	4	4	3	2	3
“maple”	5	5	4	3	2

Answer of 3

```
class Solution {
    public int climbStairs(int n) {
        if (n == 0 || n == 1) {
            return 1;
        }

        int prev = 1;
        int current = 1;

        for (int i = 2; i <= n; i++) {
            int next = prev + current;
            prev = current;
            current = next;
        }

        return current;
    }
}
```

Answer of 4

```
class Solution {
    public int longestCommonSubsequence(String text1, String text2) {

        int[][] cache = new int[text1.length() + 1][text2.length() + 1];

        for (int i = 0; i < cache.length; i++) {
            cache[i][0] = 0;
        }

        for (int j = 0; j < cache[0].length; j++) {
            cache[0][j] = 0;
        }

        for (int i = 1; i < cache.length; i++) {
            for (int j = 1; j < cache[i].length; j++) {
                if (text1.charAt(i - 1) == text2.charAt(j - 1)) {
                    cache[i][j] = cache[i - 1][j - 1] + 1;
                } else {
                    cache[i][j] = Math.max(cache[i - 1][j], cache[i][j - 1]);
                }
            }
        }

        return cache[text1.length()][text2.length()];
    }
}
```