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Data Structures Using C

Module 5 – PART 1



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01

Introduction to Graph

Introduction, Definition, Basic
Terminology & Types.

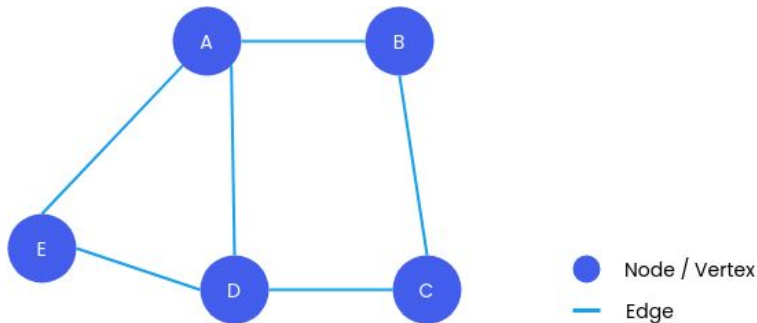


Definition

A **Graph** is a non-linear data structure consisting of nodes and edges.

Definition : A graph $G=<V,E>$ consist of two sets: V and E, where

- V is the set of **nodes** (vertices or points)
- E is the set of **edges**, identified with a unique pair of nodes in V, denoted by $e=[u, v]$.



Terminology

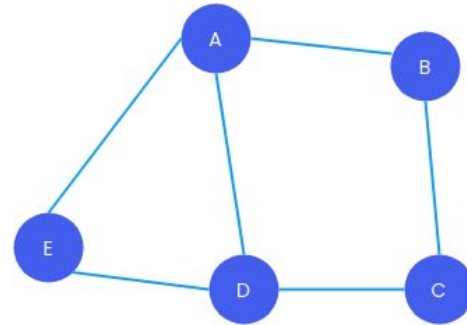
- Suppose $e = [u, v]$. Then nodes u and v are called **endpoints** of e , and u and v are called **adjacent nodes** or neighbours.
- The **degree** of node u , is the number of edges containing u .
- If u does not belong to any edge (degree 0)-then u is called **isolated edge**.
- A **path** in a graph is a finite sequence of edges which joins a sequence of vertices.
- **Path length** is equal to the number of edges present in the path.
- A path is said to be **closed** if it starts and ends at the same node.
- A path is said to be **simple** if all the nodes are distinct with exception that first node may equal last node .
- A **cycle** is a closed simple path with length 3 or more.
- A graph is said to be **connected** if and only if there is a path between any two of its nodes.
- A graph G is said to be **complete** if every node u in G is adjacent to every other node v in G .
- An edge e is called a **loop** if it has identical endpoints, that is, if $e = [u, u]$.

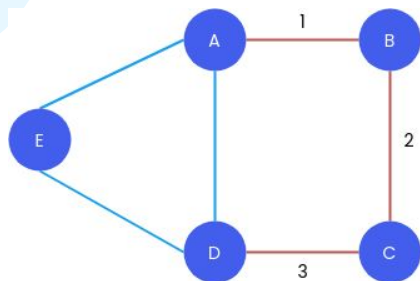
Terminology

- If there are multiple edges between the same pair of nodes in a graph, then the edges are called **parallel edges**.
- A graph containing self loop or parallel edges or both is called a **multigraph**.

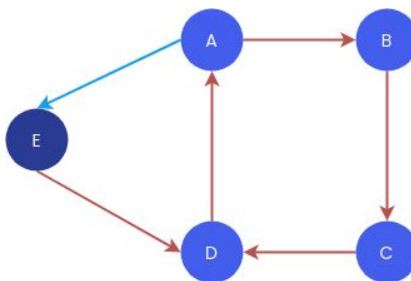
Example:

- A, B, C, D, E - **Nodes**.
- A and B are **adjacent nodes**.
- A-B, B-C, C-D, D-C, D-E, E-A - **Edges**.
- **Degree of A** = 3.
- A-B-C-D-E - **Path from A to E with length 4**.
- A-B-C - **Path from A to C with length 2**.
- B-C-D-A-B - **Closed path from B to B**.
- E-A-B-C-D-E is a **cycle**.

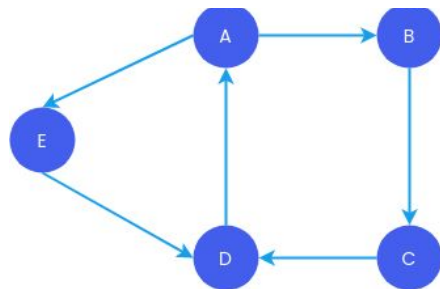
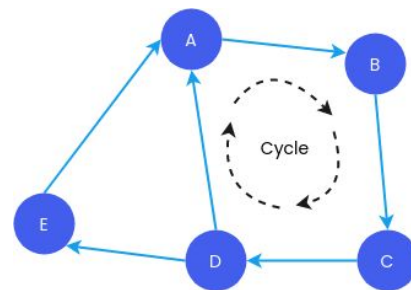




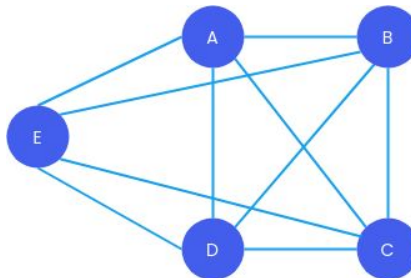
Path from A to D with length 3



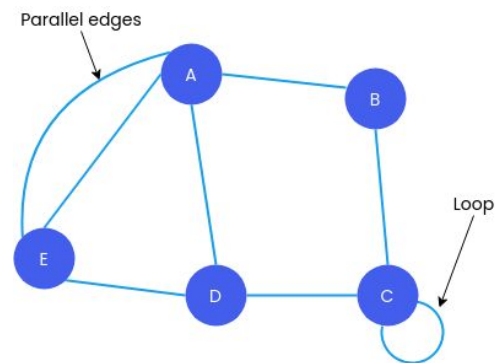
Not a simple path (E-D-A-B-C-D-A)



Connected Graph



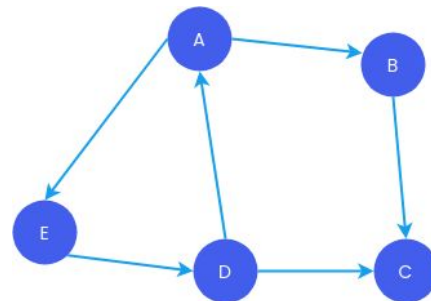
Complete Graph



Multigraph

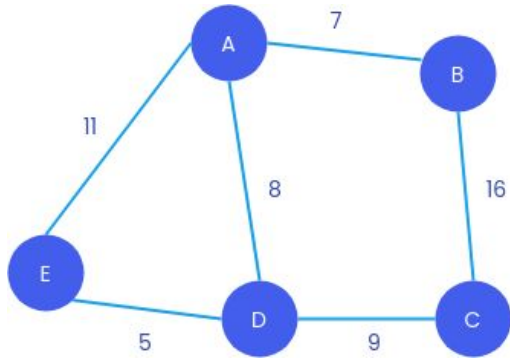
Directed Graph

- A **directed graph** is a graph with the property that its **edges** have **directions**.
- Each edge e is identified with an **ordered pair** $[u, v]$.
- They are also called a digraph.
- Suppose a digraph has a directed edge $e = [u, v]$
 - e begins at u , and ends at v .
 - u is the origin or initial point of e , and v is the destination or terminal point of e .
- The **outdegree** of a node u is the number of edges beginning at u .
- The **indegree** of a node u is the number of edges ending at u .
- A node u is called a **source** if it has a positive outdegree but zero indegree.
- A node u is called a **sink** if it has a zero outdegree but a positive indegree.
- A node v is said to be reachable from a node u if there is a directed path from v to u .
- A directed graph G is said to be connected (strongly connected) if for each pair u, v of nodes in G there is a path from u to v and there is also a path from v to u .



Weighted Graph

- A **weighted graph** is a graph with the property that its edges have a number associated with it.
- The number is called **weight** of the edge.
- A weighted graph can be undirected or directed.
- The **weight of a path** P is the sum of the weights of the edges along the path P.



02

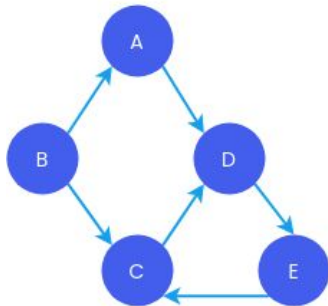
Graph Representation

Sequential & Linked list
representation of graphs.



Sequential Representation : Adjacency Matrix

- A graph having N nodes can be represented using a NxN square matrix where rows and column numbers represent the vertices.
- A cell at the cross-section of the vertices can have only values 0 or 1.
- An entry $(v_i, v_j) = 1$ only when there exists an edge from v_i to v_j .
- In case of weighted graph, the edge weight is stored instead of 1.
- Adjacency matrix will be a sparse, hence a great deal of space will be wasted.
- It may be difficult to add or delete new nodes.



	A	B	C	D	E
A	0	0	0	1	0
B	1	0	1	0	0
C	0	0	0	1	0
D	0	0	0	0	1
E	0	0	1	0	0

Path Matrix

- For a directed graph G with adjacency matrix A , the element A_{ij} of matrix A^k gives the **number of paths of length k** from vertex v_i to v_j .
- For example, an element A_{ij} of A^2 will denote the number of paths of length 2 from vertex v_i to v_j .
- Now we define a matrix B_r ,
$$B_r = A + A^2 + A^3 + \dots + A^r$$
- The ij entry of the matrix B_r gives the number of paths of lengths r or less from node v_i to v_j .
- If G has m nodes, a simple path from any v_i to v_j must have length $m-1$ or less.
- Similarly a cycle must have length m or less.
- That is if there is a non-zero ij entry in the matrix B_m , there is a simple path from v_i to v_j .
- A **path matrix** is obtained by replacing all non-zero elements in B_m by 1.
- Then $P_{ij} = 1$ if and only if there is a path from v_i to v_j .

Path Matrix

	A	B	C	D	E
A	0	0	0	1	0
B	1	0	1	0	0
C	0	0	0	1	0
D	0	0	0	0	1
E	0	0	1	0	0

A

	A	B	C	D	E
A	0	0	0	0	1
B	0	0	0	2	0
C	0	0	0	0	1
D	0	0	1	0	1
E	0	0	0	1	0

A²

	A	B	C	D	E
A	0	0	1	0	0
B	0	0	0	0	1
C	0	0	1	0	0
D	0	0	0	1	0
E	0	0	0	0	1

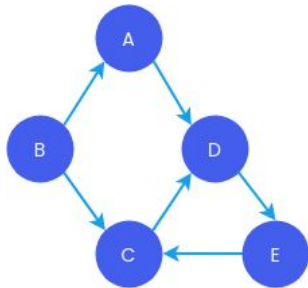
A³

	A	B	C	D	E
A	0	0	0	1	0
B	0	0	2	0	0
C	0	0	0	1	0
D	0	0	0	0	1
E	0	0	1	0	0

A⁴

	A	B	C	D	E
A	0	0	0	1	0
B	0	0	0	2	0
C	0	0	0	0	1
D	0	0	1	0	0
E	0	0	0	1	0

A⁵



	A	B	C	D	E
A	0	0	1	3	1
B	1	0	3	4	1
C	0	0	1	2	2
D	0	0	2	1	3
E	0	0	2	2	0

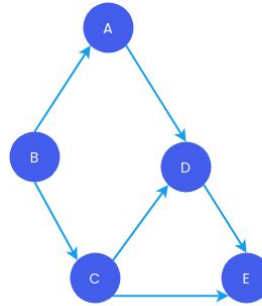
$$B = A + A^2 + A^3 + A^4 + A^5$$

	A	B	C	D	E
A	0	0	1	1	1
B	1	0	1	1	1
C	0	0	1	1	1
D	0	0	1	1	1
E	0	0	1	1	0

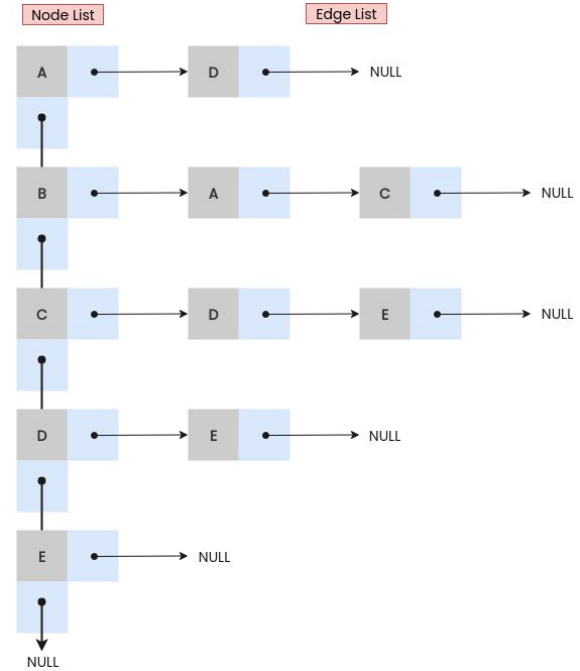
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Linked Representation

- The linked representation contains two lists, a **NODE** list and an **EDGE(Adjacency)** list.
- Each element in the NODE list will correspond to a node in the graph.
 - **Name** of the Node.
 - Pointer to **adjacent nodes**.
 - Pointer to **next node**.
- Each element in the edge list will represents an edge of the graph.
 - **Destination node** of the edge.
 - **Link** that link together the edges with the same initial node.



Node	Adjacency List
A	D
B	A, C
C	D, E
D	E
E	-



03

Traversal of Graph

Depth First Search(DFS) and
Breadth First Search(BFS).



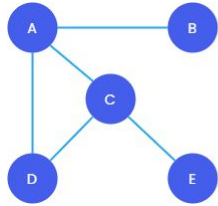
Depth First Search (DFS)

- In DFS method, the travel starts from a node then follows its outgoing edges.
- Within successors, it travels their successors and so on.
- The search goes deeper into the search space till no successor is found.
- DFS implementation puts each nodes of the graph in following lists
 - Visited
 - Not Visited (Stack)
- This marking is used to avoid cycles while traversing.

Algorithm: DFS

1. Push Starting Node A on the Stack.
2. Repeat steps 3 and 4 until stack is empty.
3. Take the top item N from the Stack and add it to the Visited List.
4. Push all unvisited neighbours of N onto Stack.
5. Exit.

Depth First Search (DFS): An Example

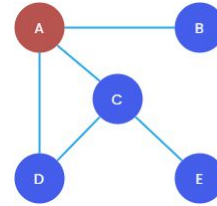


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Visited

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Stack

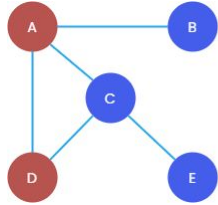


A				
---	--	--	--	--

Visited

D	C	B		
---	---	---	--	--

Stack

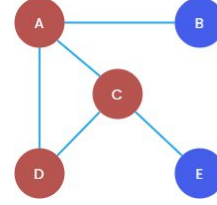


A	D			
---	---	--	--	--

Visited

C	B			
---	---	--	--	--

Stack

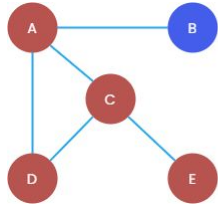


A	D	C		
---	---	---	--	--

Visited

E	B			
---	---	--	--	--

Stack

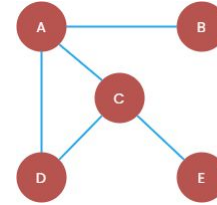


A	D	C	E	
---	---	---	---	--

Visited

B				
---	--	--	--	--

Stack



A	D	C	E	B
---	---	---	---	---

Visited

--	--	--	--	--

Stack

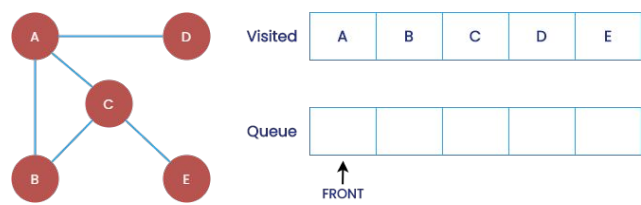
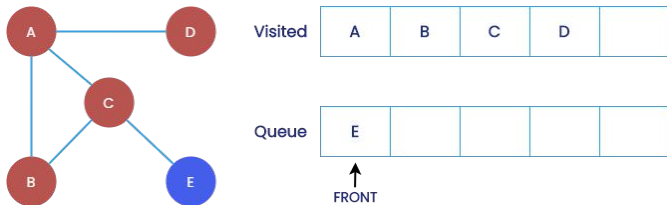
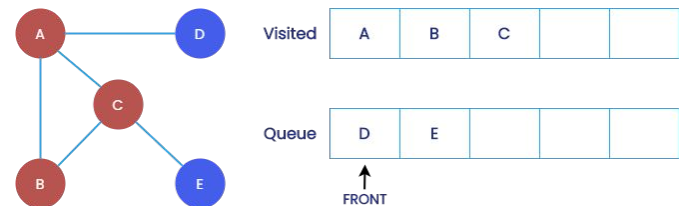
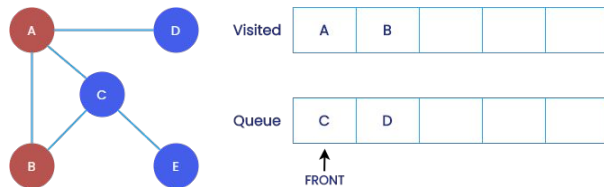
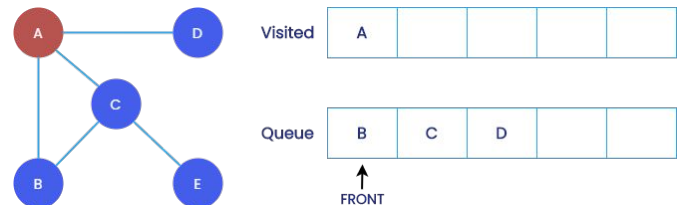
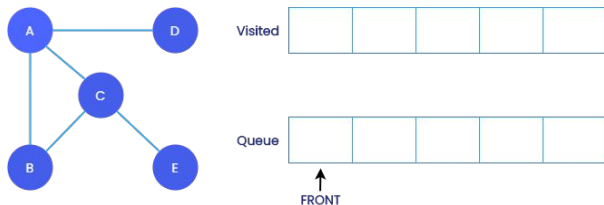
Breadth First Search (BFS)

- In BFS method, the travel starts from a node then carries onto its adjacent nodes.
- It traverse all the sibling nodes within a level and then moves to next level.
- This continues until all the vertices are visited.
- BFS implementation puts each nodes of the graph in following lists
 - Visited
 - Not Visited (Queue)
- This marking is used to avoid cycles while traversing.

Algorithm: BFS

1. Put the starting node A to the back of the Queue.
2. Repeat Steps 3 and 4 until Queue is empty.
3. Remove the front node N of Queue and add it to the Visited List.
4. Add all unvisited neighbours of N to the rear of the Queue.
5. Exit

Breadth First Search (BFS): An Example



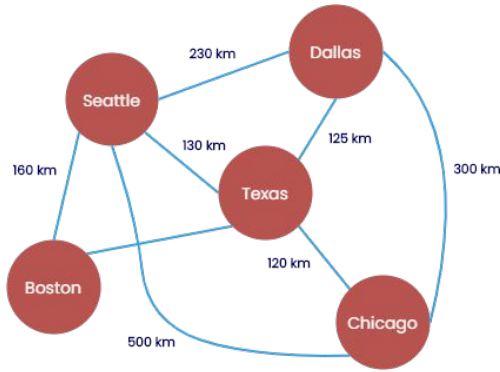
03

Applications

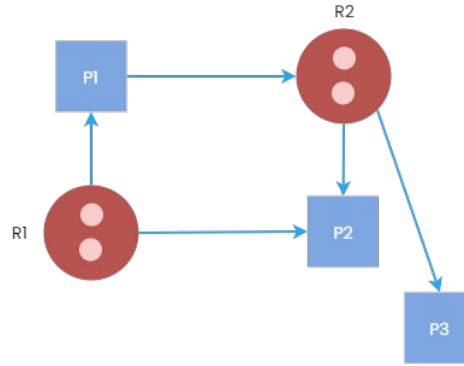
Applications of Graph data structure.



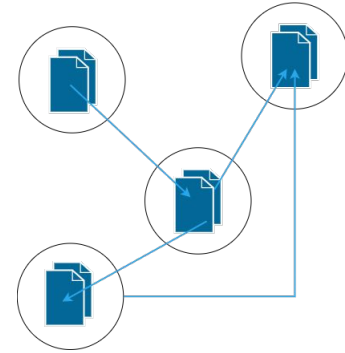
Applications



Google Map as a Graph












Resource Allocation Graph



World Wide Web as a Graph

- In Facebook, Users, Pages, Groups, etc. - nodes and edges - connection/relation.
- Path optimization Algorithms.
- Scientific Computations.
- Recommendation Algorithms.

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