

HANDS ON 3

1 Runtime given by:

$$T(n) = \sum_{i=1}^n \sum_{j=1}^n 1 = n^2$$

$\Theta(n^2)$, $\Omega(n^2)$ and $\mathcal{O}(n^2)$

2 (Plot1_Q2.jpg in repo)

3 Since n^2 denotes for large n , therefore there exists +ve constants c_1 and c_2 for all $n \geq n_0$

$$c_1 n^2 \leq T(n) \leq c_2 n^2$$

$\mathcal{O}(n^2)$ upper bound

$\Omega(n^2)$ lower "

$\Theta(n^2)$ tight "

4 For $n < 50$, timings deviate a lot from quadratic curve. For $n \geq 50$ data aligns well with quad polynomial indicates n^2 dominates and overhead is negligible. (Plot 2-Q4.jpg)

4 In modified function, inner loop performs 2 constant time operations. This increases constant factor of runtime.

5 Does not affect asymptotic analysis from part #1. Though actual runtime increases by a constant factor, asymptotic behavior remains same.