## a64strcpy

This documents details the strcpy function working for all alignments for ajit. We will introduce the basic flow of the program and then look at an example. We will then look at the testing methodology and the pseudocode will be given at the end.

## a64strcpy - introduction

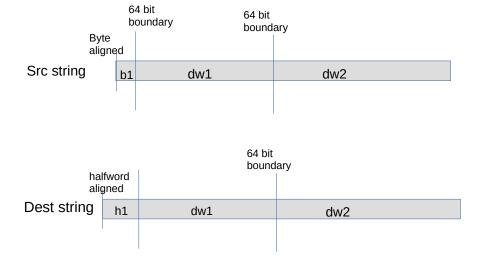
The a64strcpy function takes two inputs – a destination and a source string, and copies the source string into the destination string. The input can be of any alignment.

The logic used for the function can be summarized as follows -

- 1. We first check for alignment of source and destination
- 2. If both are aligned, we can move to direct 64 bit copies
- 3. Else if source is aligned we move to part of the function that deals with unaligned destination (step 6 onwards).
- 4. If the source is unaligned, we will first align it to 64 bits. We do this by copying the strng byte by byte till source is aligned. This will happen a maximum of 7 times (as there are only 8 bytes in a doubleword)
- 5. After we do this, destination alignement will have changed, check and if aligned go to 64 bit alignments else move ahead
- 6. We need to check destination alignment word, halfword or byte
- 7. If it is word aligned
  - 1. store a word, making it 64 bit aligned
  - 2. save the unstored word in another register
  - 3. load a new doubleword
  - 4. create a doubleword with the unstored word from previous doubleword, and the most significant word from recently loaded doubeword (use shifts and or operations)
  - 5. load this doubleword at once in destination
  - 6. repeat from step2
- 8. If it is halfword or byte aligned we follow similar steps as above, i.e., perform store till destination is aligned to 64 bit, save the rest of the part of doubleword. Load a new doubleword. Create the doubleword to be stored by combining required parts of previously saved and recently loaded doubleword and finally perform a doubleword store with it.

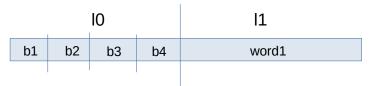
## a64strcpy – example

Assume – source address is byte aligned and destination address is halfword aligned



- Step 1: Store byte b1 from source string to destination string at h1.

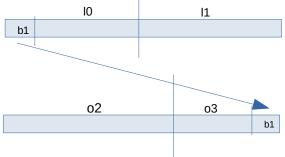
  Now, source is doubleword aligned and destination is byte aligned as b1 will only fill half of h1.
- Step 2: Load a new doubleword from source in register pair. And load b1 into destination. This makes the destination doubleword aligned.



Step 3: Shift 10-11 left by 8 bits and store in 12-13. As we have stored b1 already, we do not need it.



Step 4: Load a new doubleword into l0-l1. We need the most significant byte from the new doublword. Hence we shift l0-l1 right by 7 bytes and save it in another register pair, o2-o3.



- Step 5: Now we OR the two registers l2-l3 and o2-o3. We now have a doubleword we can store in the destination.
- Step 6: Store the doubleword and and repeat from step2 till we see null byte.

## a64strcpy - Testing

We need to test for all possible scenarios of alignment which will cover the whole code.

We start by defining a doubleword aligned source and destination. We then give the doubleword aligned source and destination as input.

For the second iteration, we input the doubleword aligned but we give source string address as doubleword aligned +1, i.e., byte aligned input. This will reduce the length of the string. In this case, as we first align the source string to 64 bits, we would have to store a byte and then destination would be byte aligned and we would be able to test that scenario.

Hence, the misalignment in the source will be eventually reflected in the destination. As there are 8 bytes, if we add one 8 times to the source string, we would cover all alignments for source, and as a result for destination too. This would cover the whole code atleast once.

#### Pseudocode -

- 1. Define 64 bit aligned destination and source strings. -
- 2. Initialize for loop to loop 9 times
- 3. For first iteration input aligned source and destination addresses
- 4. For further iterations input aligned destination address, but input source address is (aligned address + iteration count)

This excludes the test for finding null byte at various locations as that will be covered in a different test.

## a64strcpy - pseudocode

Pseudocode of the function is given below.

```
1. check alignment of src and dest
if both aligned
goto 64 bit string operation
elsif src unaligned
goto alignsrc
else dest unaligned
goto destuna
```

```
2. 64bit string operation - ()
load doubleword of src string
check for zero byte
if found
goto checkzero
else
store doubleowrd to dest string
goto 64 bit string operation
```

```
3. checkzero: (checks which byte has zero and stores in dest string accordingly) shift most significant byte to least significant byte position test if it is zero if yes

load it

return
repeat for all other bytes
```

4. alignsrc: (aligns src string address to doubleowrd)

```
load a src string byte
check if its zero
if yes
store it and return
check if source is doublword aligned
if no
store byte in destination
goto alignsrc
else check destination for doubleword alignment
if aligned
goto 64bit string operation
else
goto destinua
```

# 5. destinua: operation)

(here, we will have a doubleword aligned source and will perform the string

load a doubleword from src string into l0 - l1 check for zero bytes if yes

goto check zero add 8 to dest string

else

check alignment of destination

if

word aligned goto wordal

elsif

halfword aligned goto hwordal

else

goto bytealigned

6. wordal: (dest string is word aligned, source string is doublword aligned)

store the word in 10 to destination

add 4 to dest address

wordalag:

shift l1 to l6

load another doubleword from source string into l0-l1

check for zero byte if no zero byte

move 10 to 17

store 16-17 to destintaion

goto wordalag to repeat process

else

store the word in 16 to destination string

add 4 to destination address

goto check zero

7. hwordal: (here, destination string is halfword aligned and source string is double word aligned)

store the most significant halfword of l0 to destination string check if destination is dowrd aligned if yes
goto hword2al
else

store next word in destination address

hword1al:

shift the least significant halfword in 10 to most significant position in 16 load another doubleword in 10 shift 10 to the right by a halfword and store in o2 or 16-17 and o2-o3 and store in 16-17 check for zero if no store the double word in 16-17 goto hword1al

store halfoword in 16 into destination string goto check zero

hword2al:

else

shift the top three halfwords in l0-l1 to l6-l7 in the same position load another doubleword in l0-l1 shift the most significant hword in l0-l1 to least significant hword in o2-o3 or l6-l7 and o2-o3 and store in l6-l7 check for zero byte if no store l6-l7 into destination goto hword2al else store the three halfwords in l6-l7 into destination check for zero in l0-l1

8. bytealigned: (here destination string is byte aligned and source string is doubleword aligned)

store the most significant byte to destination string check if destination is doubleword aligned if yes

goto byte1
store the next halword to destination string check if destination is doubleword aligned if yes

goto byte3
store the next halword to destination string

check if destination is doubleword aligned if yes
goto byte5
store the next halfword to destination string

### byte7:

shift the least significant byte of l0 into l6 at most significant position load the next doubleword in l0 shift the top 7 bytes to l0-l1 by 8 and store into o2-o3 or l6-l7 and o2-o3 and store into l6-l7 check for zero byte if no store l6-l7 into destination and goto byte7 else

store the byte in 16 into destination check for zero byte in 10-11