

# Graph Optimizer for Japanese Mahjong Tile-Discard-Finder

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Optimization Methods for Engineers



# Introduction: Problem description

- Japanese Mahjong (Richii Mahjong) is a very complex game with a lot of rules and mechanics
- For the sake of simplicity, *we will reduce the game to the following rulesets*: At the start of the game, we put 14 Tiles aside and the 4 players draw 13 Tiles each. On a players turn, that player draws a Tile and then discards one if he does not have a winning hand yet. A winning hand consists of 4 groups of three Tiles (sequences of the same colour or identical Tiles) and a pair of identical Tiles.
- Our **goal**: Given a random Hand consisting of 14 Tiles which all have a colour and a value, find the best Tile to discard and reach a winning hand.
- Tiles can either be grouped into Sequences or Groups of equals and when a Tile belongs to multiple groups, we need to decide which grouping is more beneficial for the Hand.
- Based on the unknown Tiles in the Stack we would like to take the probability of drawing a certain Tile into consideration
- *Our Optimization only covers a part of the real Japanese Mahjong*, We do not consider things like ron (a Player completing their Hand with another's discard) or the other players winning.

# Introduction: In mathematical Terms

- The Stack consists of 136 Tiles
- A Hand consists of 14 Tiles
- A Tile is a Tuple  $(type, value)$
- There are 34 distinct Tiles, all of which are 4 times in the Stack
- Tiles of type  $\{Pin, Character, Bamboo\}$  have values  $\{1, \dots, 9\}$
- Tiles of type  $\{North, West, South, East, Red, Green, White\}$  only have value  $\{1\}$
- **Shanten** is a positive Integer that denotes the *minimal edit distance* for a **ready Hand**
- A **ready Hand** consists of 13 Tiles such that the next draw could complete it to a **winning Hand**
- A **winning Hand** consists of 14 Tiles such that there are 4 Triples and one Pair

# The naïve Algorithm

## What our naïve Algorithm does:

- Check if there are still Tiles in the Stack; if not end the game in a loss
- Draw a Tile from the Stack
- Calculate Shanten for the current Hand  $H$
- If Shanten = -1 (a winning Hand); end the game as a win
- For all Tiles in our Hand  $H = \{t_1 \dots t_{14}\}$  check if they contribute to lowering the current Shanten
- Mark those who do not lower Shanten in the current hand as *possibleDiscard*
- From  $possibleDiscard \subseteq H$  choose one Tile  $t_x$  at random and discard it
- We end our turn

This is a **greedy** algorithm that uses the knowledge of whether a tile is currently part of Shanten or not to find a discard that for sure does not worsen ones Hand  $H$ . It runs in  $O(14) \leq O(1)$

# The DrawAnalyzer Algorithm

## What our improved Algorithm does:

- Generate a graph of distinct combinations of unordered draws up to *depth* draws
- Use the deepest layer of the graph to find the tile that can be discarded in the most scenarios

Our **goal** is to...

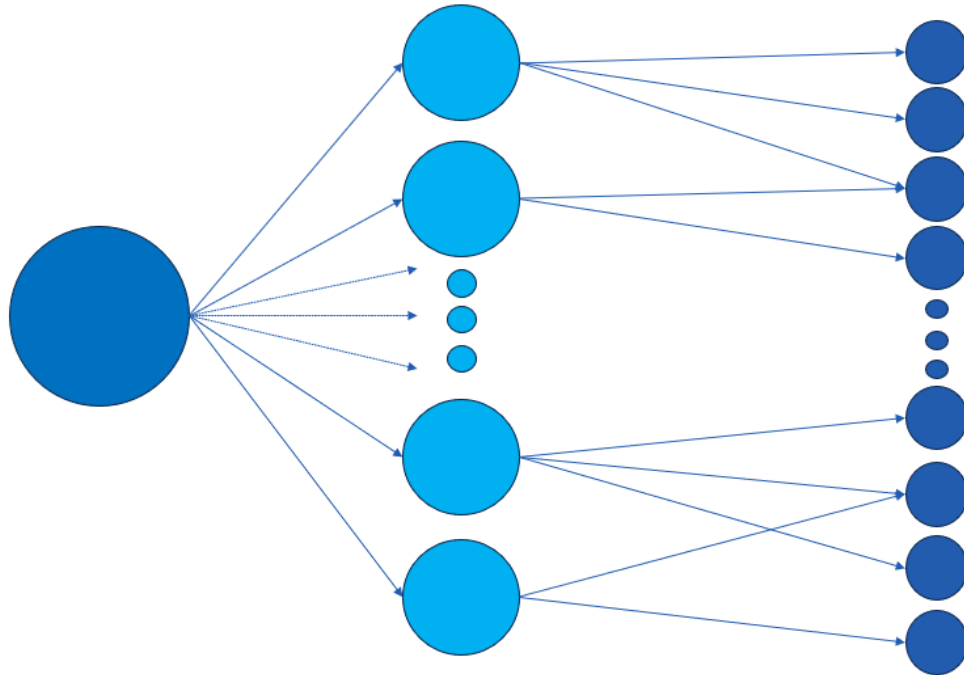
1. Improve the Win/Game ratio
2. Improve the reliability of our discard
3. Limit the average discard time to a maximum of 20s

# The DrawAnalyzer Graph

## How we build the graph:

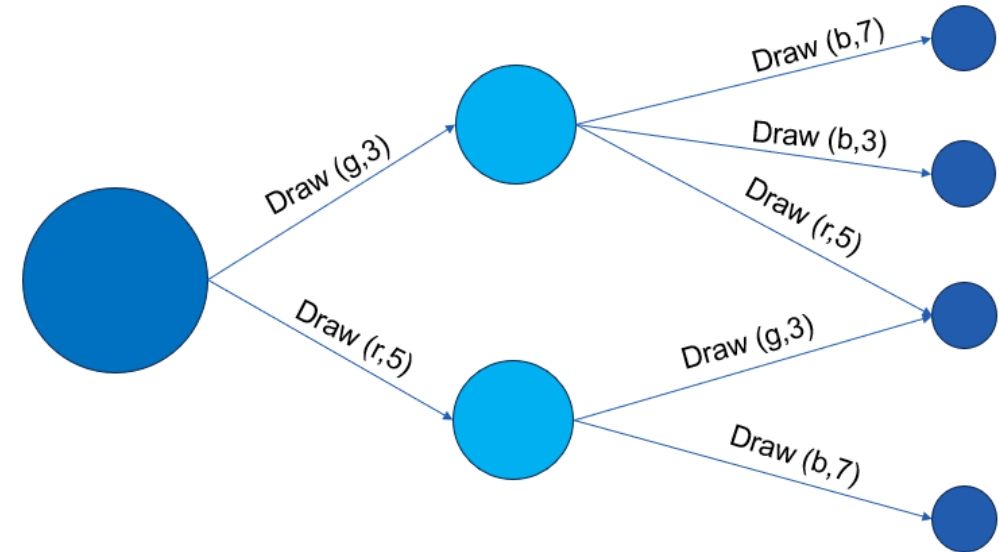
- Given a depth  $n$ :
- Generate the first  $\binom{34}{1}$  vertices which represent a draw of 1 unique tile and calculate the number of possibilities of drawing that Tile.
- Generate the next  $\binom{34+i-1}{i}$  for  $2 \leq i \leq n$  vertices of distinct combinations of  $i$  tiles, calculate the number of possibilities of drawing that set of tiles and add an edge from a previous vertex to the new one if the new vertex is a draw order reachable from the previous one
- A total of  $\sum_{i=1}^n \binom{34+i-1}{i}$  vertices will thus be generated
- Traversing the graph simulates following a draw order

# DrawAnalyzer visualization



## Example Graph with depth 2

The current Gamestate is connected to all 34 possible Tile-Draws, which in turn are connected to further draws.



## With some more Detail

Gamestates that have had the same draws in a different order are considered as equal but more likely than other states.

# DrawAnalyzer best discard

## How we calculate the best discard:

- Given our previously generated graph, we iterate over the deepest layer until  $depth > \#draws\ left$ , then we iterate over the first layer to save computation time
- For each vertex, we calculate the number of possibilities for this set of tiles to be drawn
- We then **simulate drawing the tiles** and calculate the Shanten
  - For  $n$  tiles drawn, we generate up to  $2(14 + n)^3 + 2(14 + n)^2$  groups. However, the average will be significantly lower at an estimated  $(14 + n)^2$ .
  - We then **generate all possible combinations** of up to 5 groups and calculate  $G = \{(t, a) \in T \times \mathbb{N} \mid t \notin \bigcap_{s \in S} s \cap H, a = comb\}$  where  $T$  is the set of all unique Tiles,  $S$  is the set of sets of unique tiles that form the largest possible groups in the given hand,  $H$  is the set of unique tiles in the hand and  $comb$  the possible number of combinations for drawing the set of tiles given by the vertex.
  - Finally, for each element in  $G$ , we add the value  $a$  to the tile  $t$  of the original hand and discard the tile with the highest number, thus the one that is discardable in most cases



# How we compare the Algorithms

## How much data

- We run 1000 Games on our naïve Algorithm and on DrawAnalyzer with depth 1
- We run 100 Games on DrawAnalyzer with depth 2 because it takes significantly longer

## Significant measurements

- We measure the time it takes to draw per round
- We measure the average Shanten per round
- We measure in what round a Game is won (if at all)

## Machine Specifics

- 2th Gen Intel(R) Core(TM) i5-12600K (16 CPUs), ~3.7GHz 32GB Ram NVIDIA GeForce RTX 2060

# Results: Naïve Algorithm

## 1000 Games with the Naïve Algorithm

Round	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
Total Time	54.40266																	
Winning Hands	109																	
Avg Discard Time	0.00309																	
Avg Shanten per Round	3	3	2	2	2	1	1	1	1	1	1	1	1	1	0	0	0	0
Avg Time per Round	0	0.00100	0.00118	0.00147	0.00165	0.00215	0.00246	0.00289	0.00296	0.00326	0.00379	0.00437	0.00411	0.00503	0.00504	0.00522	0.00603	0.00293
Ready Hand is reached	0	1	1	0	0	4	1	2	1	8	9	8	13	8	12	19	13	9

# Results: DrawAnalyzer with depth 1

## 1000 Games of DrawAnalyzer with depth 1

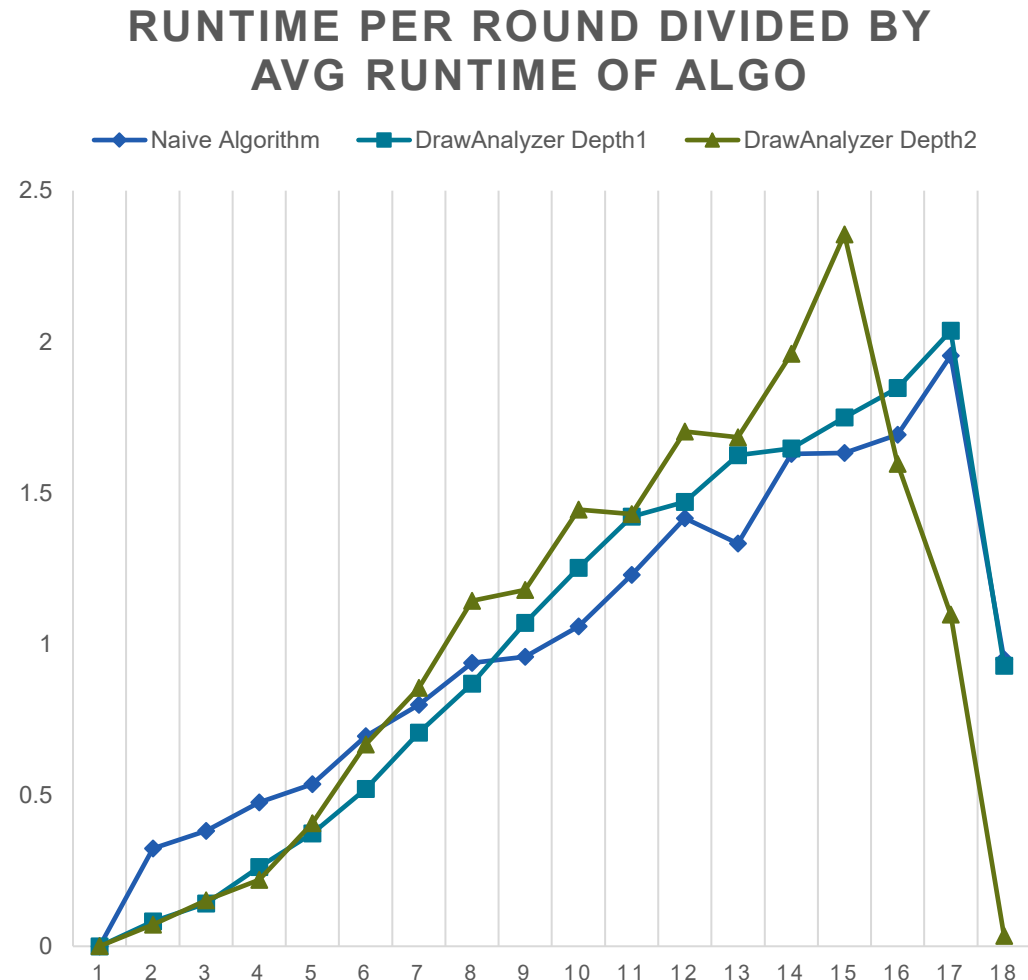
Round	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
Total Time	5661.174																	
Winning Hands	196																	
Avg Discard Time	0.307256																	
Avg Shanten per Round	3	2	2	2	1	1	1	1	1	1	0	0	0	0	0	0	0	0
Avg Time per Round	1.16E-05	0.02518	0.04354	0.08048	0.11459	0.15996	0.21706	0.26701	0.32887	0.38456	0.43677	0.45172	0.49916	0.50591	0.53749	0.56738	0.62565	0.28518
Ready Hand is reached	0	0	0	1	1	3	7	2	7	11	11	16	17	28	20	29	30	13

# Results: DrawAnalyzer with depth 2

## 100 Games of DrawAnalyzer with depth 2

Round	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
Total Time	19402.21																	
Winning Hands	29																	
Avg Discard Time	11.69604																	
Avg Shanten per Round	3	2	2	2	1	1	1	1	1	0	0	0	0	0	0	0	0	0
Avg Time per Round	0.00316	0.83298	1.77464	2.56912	4.76824	7.80642	9.99863	13.3688	13.7893	16.8931	16.7270	19.9140	19.6917	22.9336	27.5460	18.6651	12.8424	0.40397
Ready Hand is reached	0	0	0	0	0	1	0	0	1	5	0	1	2	2	6	6	2	3

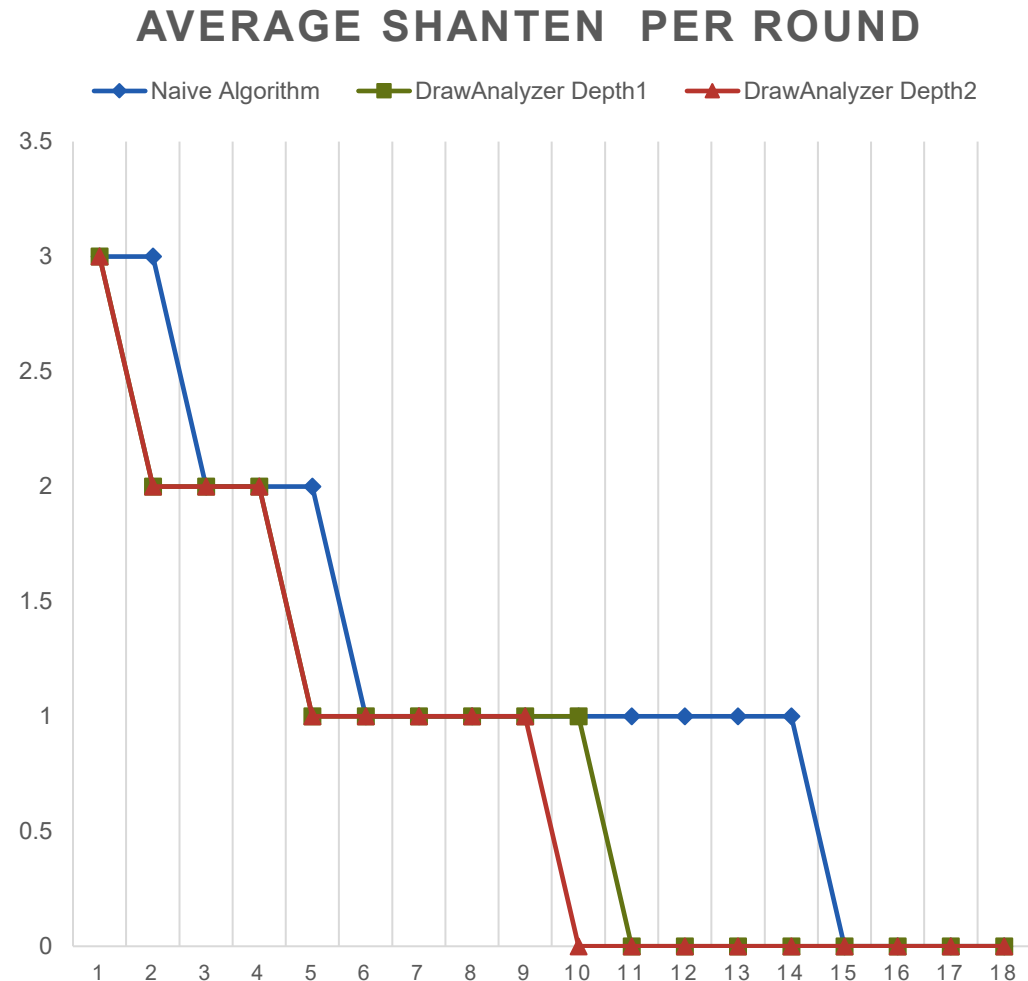
# Comparing Results I



## Comparison of times

- The naïve Algorithm takes an average of  $3 * 10^{-3}s$  per draw
- DrawAnalyzer with depth 1 takes an average of  $0.3s$  per draw
- DrawAnalyzer with depth 2 takes an average of  $11.7s$  per draw
- DrawAnalyzer with depth 1 takes on average 90times longer than the naïve algorithm
- DrawAnalyzer with depth 2 takes  $\sim 3500$ times longer than the naïve algorithm on average
- Runtimes tend to be longer the more the game progresses, unless in the last few rounds

# Comparing Results II



## Winning Hands reached

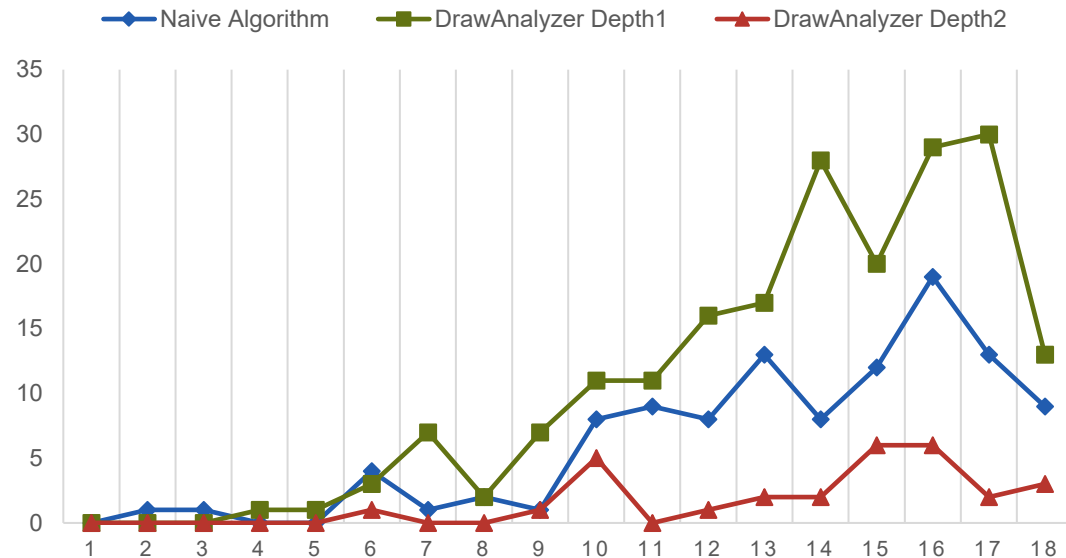
- With the naïve Algorithm ~11% of rounds end in a ready Hand
- With the DrawAnalyzer depth 1 ~20% of rounds end in a ready Hand
- With the DrawAnalyzer depth 2 ~30% of rounds end in a ready Hand

## Improving Shanten

- The average **Shanten** per round improves significantly earlier when using DrawAnalyzer of either depth, which implies faster **winning Hand** due to better decisions.

# Comparing Results III

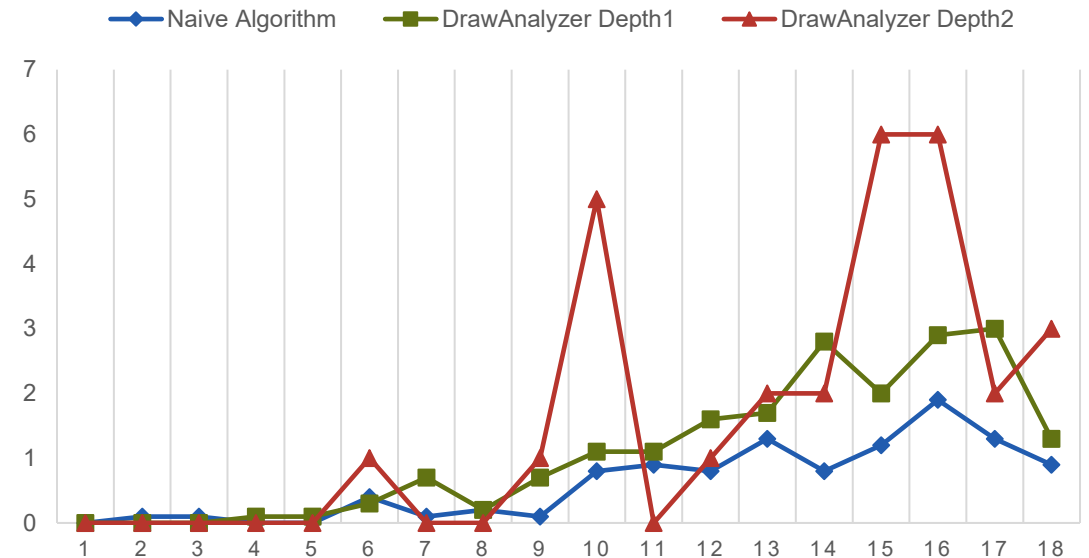
## IN WHAT ROUND IS READY HAND REACHED



### Ready Hands are reached more consistent

In this unadjusted Diagram we see that DrawAnalyzer for depth1 clearly outperforms our naïve algorithm

## IN WHAT ROUND IS READY HAND REACHED (ADJUSTED)



Here we have adjusted the Diagram for DrawAnalyzer with depth 2. We also see a trend here that it would perform significantly better than the other two algorithms.

# Discussion and Conclusion

- **We improve the success rate** from ~11% to >20% even with depth 1 of DrawAnalyzer practically doubling it at the cost of factor 90 in terms of calculation time
- With higher depth of DrawAnalyzer the performance seems to be improved even more; towards 30% success rate but with a cost of roughly factor 3500, which seems to be unproportionate.
- We see that the calculation time for DrawAnalyzer improves in the last few rounds because the depth is capped by the remaining tiles on the stack.
- **Shanten gets improved earlier** in the game with both iterations of DrawAnalyzer, which shows that the algorithm constantly performs goal-oriented and makes good decisions throughout the game.
- **DrawAnalyzer with depth 1 seems to be the current best solution** in terms of usability, because the DrawAnalyzer with depth 2 takes too long to calculate.



# Links

## GitHub Repository

- <https://github.com/adidell01/MahjongOptimization>

## Image1 / Game Explanation

- <https://www.tatlerasia.com/lifestyle/entertainment/hk-mahjong-beginners-guide-how-to-play>