

Disproof by counterexample.

original claim: If x, y are rational, then xy is rational.

claim: If xy is rational, then x, y are rational.

$$x = \frac{1}{\sqrt{2}} \quad y = \frac{1}{\sqrt{2}}$$

$$xy = \frac{1}{\sqrt{2}} \cdot \frac{1}{\sqrt{2}} = \frac{1}{2}$$

xy is rational, but x, y are not rational.

Therefore, this claim does not hold.

$$x = 2 \quad y = 6 \quad xy = 12$$

Det

A disproof by counter example constructs an example which the claim is false and explains why it is false.

Det

Let n, m be integers, then n is divisible by m , if there exists an integer k , such that $n = m \cdot k$

Q: Is 10 divisible by 2?

$$10 = 2 \cdot 5 = 2 \cdot 5 \quad (5 \text{ is an integer})$$

Q: Is 11 divisible by 3?

$$11 = 3 \cdot \left(\frac{11}{3}\right) \leftarrow \text{is not an integer}$$

Q: Is 0 divisible by 4?

$$0 = 4 \cdot 0 \leftarrow 0 \text{ is an integer.}$$

Def

we say, if n is divisible by m ,
we say m divides n

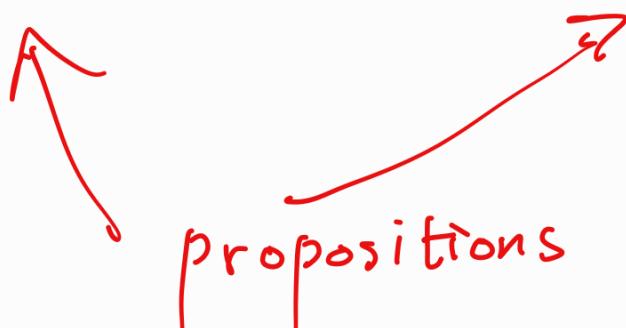
$$m \mid n$$

$$2 \mid 10$$

$$3 \nmid 10$$

2 divides 10

3 does not divide 10



claim! let n be an integer,
then $n \cdot (n+1)^2$ is even.

Step 01:

What is even? If n is divisible
by 2, then n is even.

Then I can write $n = 2 \cdot k$, where k
is an integer.

Step 02!: Let's do some examples.

n	$n \cdot (n+1)^2$	is $n \cdot (n+1)^2$ even?
0	$0 \cdot 1^2 = 0$	✓
1	$1 \cdot 2^2 = 4$	✓
2	$2 \cdot 3^2 = 18$	✓
-2	$-2 \cdot (1)^2 = -2$	✓
3	$3 \cdot 4^2 = 48$	✓

Observe that n can be any integer. Any integer has to be either odd integer or an even integer.

Therefore, If we show that when n is odd $n \cdot (n+1)^2$ is even and when n is even $n \cdot (n+1)^2$ is even, then we can claim that our proposition is true.

Step 03?

case 01: n is even

statements

reasoning

$n = 2 \cdot c$, c is an integer

by def. of even

$$n \cdot (n+1)^2 = 2 \cdot c (n+1)^2$$

by substitution

$$n \cdot (n+1)^2 = 2 \cdot (c \cdot (n+1)^2)$$

by rearranging

$$n \cdot (n+1)^2 = 2 \cdot K, \text{ where}$$

" "

$$K = (c \cdot (n+1)^2)$$

K is an integer

product of integers is an integer

$$n \cdot (n+1)^2 = 2 \cdot K$$

$$n \cdot (n+1)^2 \text{ is even}$$

by def. of even.

Case 2: n is odd

$n+1$ is even

because n is odd

$n+1 = 2 \cdot b$, b is an integer

by def. of even

$$n \cdot (n+1)^2 = n \cdot (2b)^2 = n \cdot 4 \cdot b^2 \quad \text{by substitution}$$

$$n \cdot (n+1)^2 = 4 \cdot b^2 \cdot n = 2 \cdot (2 \cdot b^2 \cdot n) \quad \text{by rearranging}$$

$2 \cdot b^2 \cdot n$ is an integer

product of ints. is an int.

$n \cdot (n+1)^2$ is even

by def of even.

Since n is either even or odd and in either case $n(n+1)^2$ is even, therefore our claim is true.

Det.

To give a proof by cases of a proposition φ , we identify set of cases, and then prove two different types of facts.

1. In any case φ holds
2. One of the cases must hold.

claim: Let x be a real value.
Then $-|x| \leq x \leq |x|$

Step 01:



$$|x| = \begin{cases} x & , \text{ if } x \geq 0 \\ -x & , \text{ if } x < 0 \end{cases}$$

$$x = 3$$

$$x = -5$$

$$x = 0$$

$$|x| = 3$$

$$|-5| = 5$$

$$|0| = 0$$

Step 02! $-|x| \leq x \leq |x|$

x	$ x $	$- x $	is $- x \leq x \leq x $
2	2	-2	✓
0	0	0	✓
-3	3	-3	✓

Case 01: $x \geq 0$

$$|x| = \begin{cases} x, & \text{if } x \geq 0 \\ -x, & x < 0 \end{cases}$$

$$|x| = x \quad \text{by def of } |x|$$

$$-|x| = -x \quad \text{by algebra}$$

$$-|x| \leq x \leq |x| \quad \text{by substitution}$$

case or!

$$x < 0$$

$$|x| = \begin{cases} x & , x \geq 0 \\ -x & , x < 0 \end{cases}$$

$$|x| = -x$$

by def of
 $|x|$

$$-|x| = x$$

by algebra

$$-|x| \leq x \leq |x|$$

by substitution

I have considered all possible cases for x , and we have proved that for each case the claim holds.

Therefore, $-|x| \leq x \leq |x|$



Claim! Let x, y, z be real numbers,
Then we say that

$$|x-y| \leq |x-z| + |y-z|$$

Step 01:

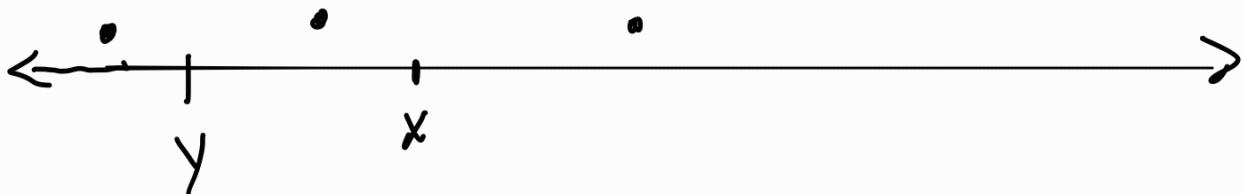
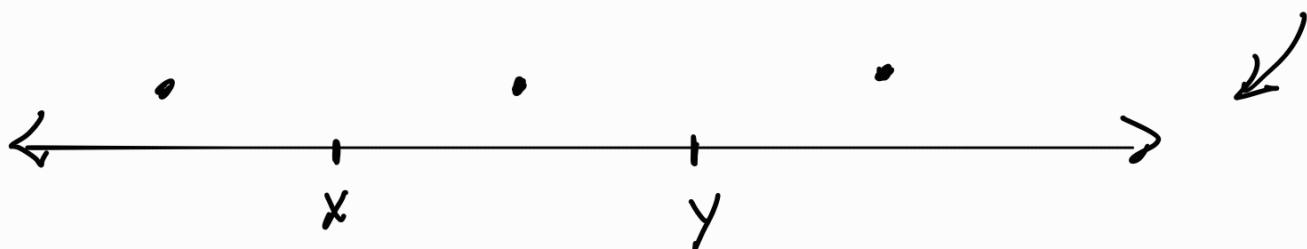
Step 02:

$$x = 4 \quad y = 5 \quad z = 2$$

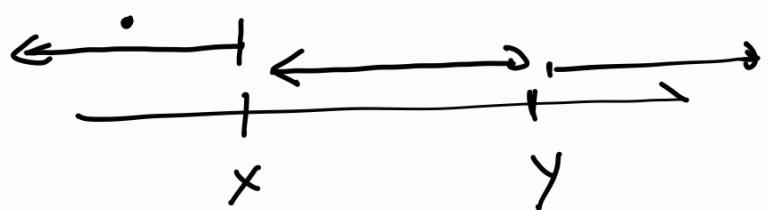
$$1 \leq 2 + 3 \quad \checkmark$$

$$x = -2 \quad y = -3 \quad z = 1$$

$$1 \leq 3 + 4$$



$$x \leq y$$

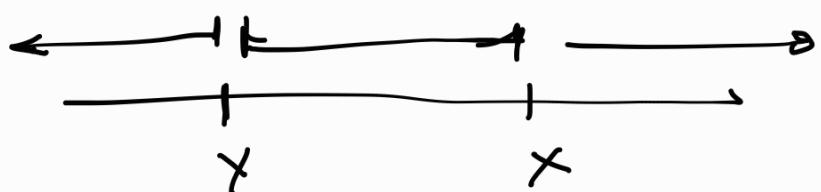


$$1. z \leq x$$

$$2. x < z \leq y$$

$$3. y < z$$

$$x > y$$



$$4. z \leq y$$

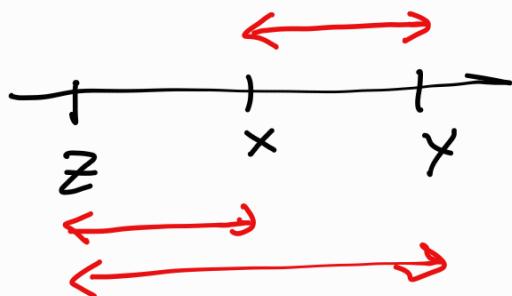
$$5. y < z \leq x$$

$$6. x < z$$

Proof:

case 01: $x \leq y, z \leq x$

$$|x-y| \leq |x-z| + |y-z|$$



$$|y-z| = |y-z|$$

$$|y-z| + |x-z| \geq |y-z|$$

$$|y-z| + |x-z| \geq y-z$$

$$|y-z| + |x-z| \geq y-z \geq y-x$$

$$|y-z| + |x-z| \geq y-x$$

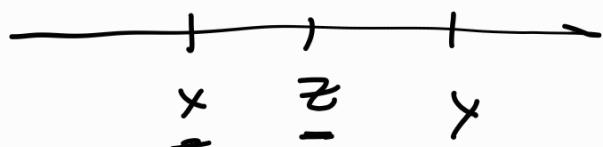
$$|y-z| + |x-z| \geq |y-x|$$

$$|y-z| + |x-z| \geq |-(x-y)|$$

$$|y-z| + |x-z| \geq |(x-y)|$$



Case 2! $x \leq y$, $x < z \leq y$



$$|x-y| \leq |x-z| + |y-z|$$

$$|x-z| + |y-z| = \cancel{x-z} + y - \cancel{z} \quad (\text{by def of } |\cdot|)$$

$$= y - x$$

$$= -(x-y)$$

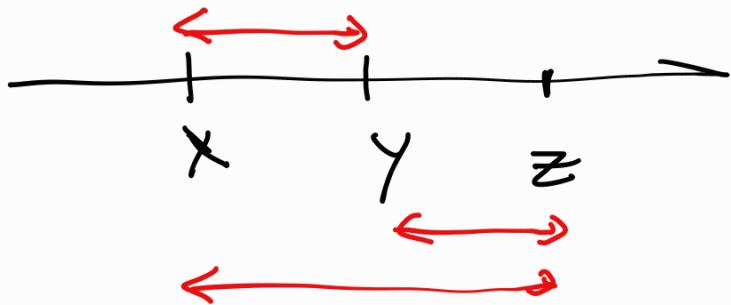
$$|x-z| + |y-z| = |x-y|$$

$$|x-z| + |y-z| \geq |x-y|$$

claim is true
for case 2.

case 3! $x \leq y, z > y$

$$|x-z| + |y-z| \geq |x-y|$$



$$|x-z| = |x-y| + |y-z|$$

$$|x-z| + |y-z| \geq \underbrace{|x-y|}_{z-x}$$

$$|x-z| + |y-z| \geq z-x$$

$$\begin{aligned} z &> y \\ z-x &> y-x \end{aligned}$$

$$|x-z| + |y-z| \geq z-x > y-x$$

$$|x-z| + |y-z| \geq y-x$$

$$\text{if } " \quad \geq -(x-y)$$

$$|x-z| + |y-z| \geq |x-y|$$

case is proven.

case 4, 5, 6

$$x > y$$

$$|x-y| \leq |x-z| + |y-z|$$

$$x = t \quad y = r \quad t > r$$

$$|t-r| \leq |t-z| + |r-z|$$

$$|t-r| \leq |r-z| + |t-z|$$

$$t = y$$

$$r = x$$

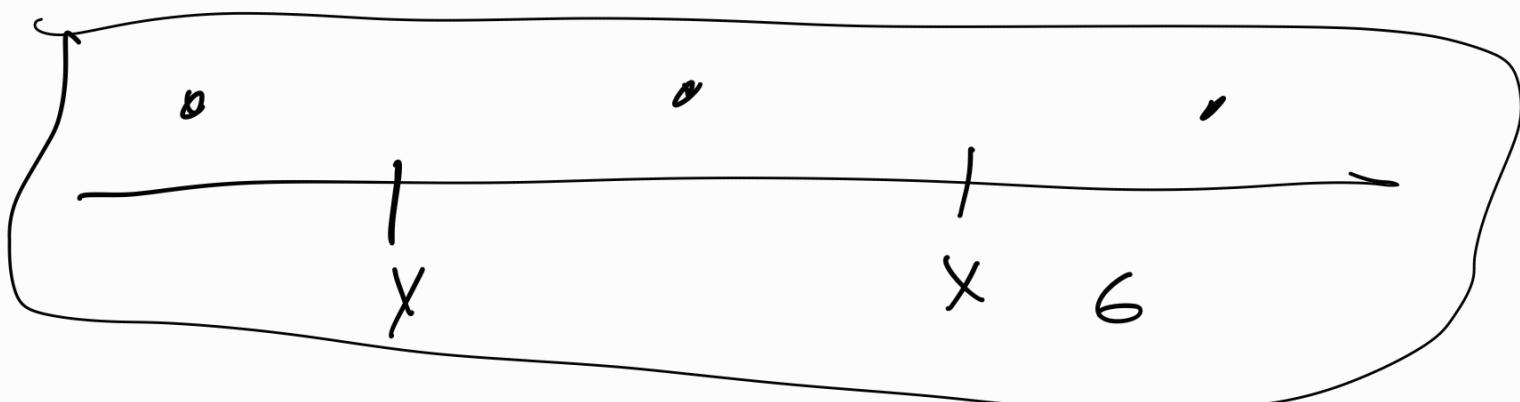
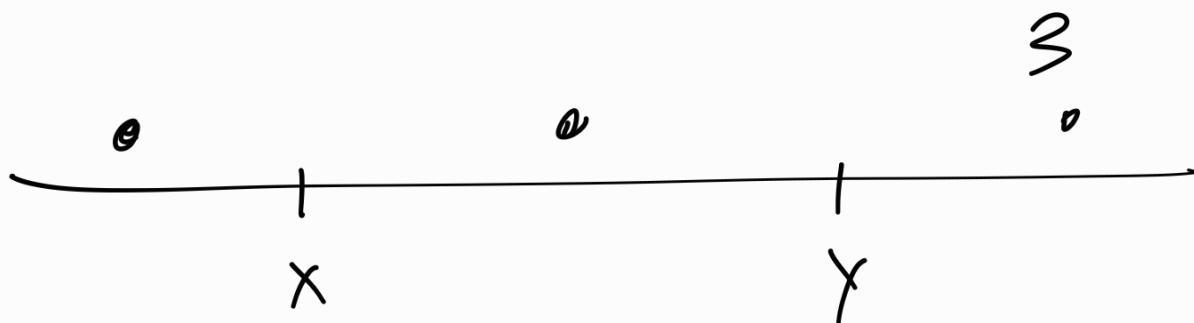
$$Y > X$$

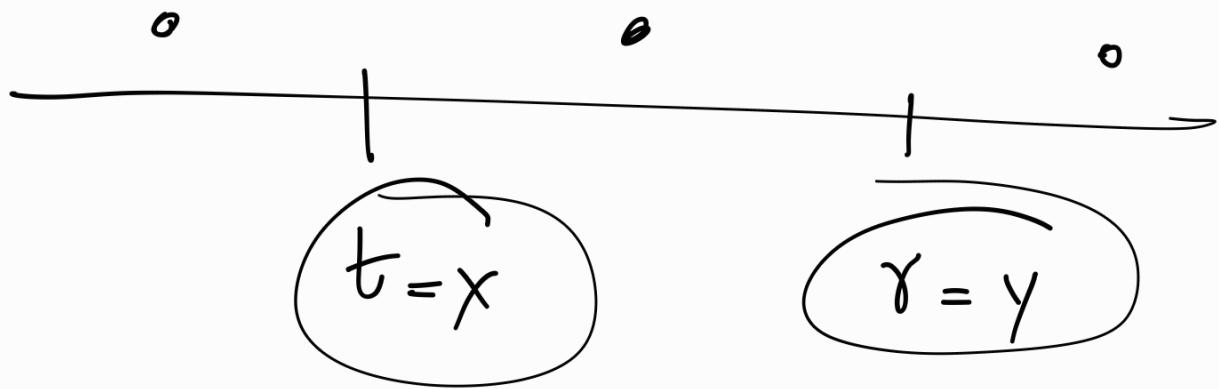
$$|y-x| \leq |x-z| + |y-z|$$

we proved
 $x \leq y$

$$x \leq y$$

case earlier.





case 4, 5, 6 are identical to cases 1, 2, 3. Therefore, the claim holds for cases 1, 2, 3, 4, 5, & 6.

Hence, the claim is true.

