Let
$$e = \{u, v\}$$
 or (u, v)

- nodes u, v are adjacent or n
- In a directed graph,

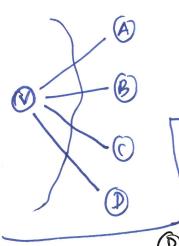
 v is an out-neighbour of

 u & u is an in-neighbour

 of v.
- · u, v are endpoints of e · u, v are incident to e

Let V be a node in a simple undirectel graph.

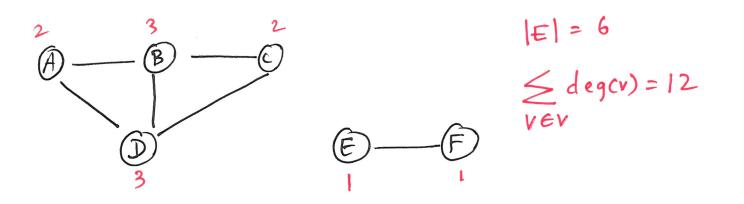
degree (v) = deg(v) = d(v) = # of neighbors of V. $= \left| \{ u \in V : \{u,v\} \in E \} \right|$



for directed graphs,
indeg(v) = # in-neighbors of v.
ontdeg(v) = # out-neighbors of v.
indeg(B) = 1 out-deg(B) = 2

In class Activity

For each graph, label each node V with degcv), and give \leq degcv), the total degree of the graph, and |E|, vev the number of edges in the graph.



What do you think about the relationship between $\leq \deg(v)$ & |E|

Theorem 11.8 Handshaking Lemma Let G=(V,E) be an undirected graph. $\leq \deg(v) = 2 \cdot |E|$ Proof: Let $G_1 = (V, E)$ an undirected graph. Notice that every edge is that connected to exactly two nodes, meaning that it contributes I to the degree of two nodes. ≤ deg(v) = 2.|E|

E = { { 1, 4 3, { 2, 4 3, { 2, 5 3, 4 3, { 4,5 }}} V = {1,2,3,4,5} G= (V, E)

Corollary > fact that follows simply from a previous lemma or theorem.

Let nodd denote the # of nodes whose degree is odd. Then hodd is even.

Proof: Aiming for a contradiction Assume the negation of the claim is true. sum of odd #odds nodd is odd. $\leq deg(v) = \leq deg(x)$ x ∈ {y ∈ V : deg(y) is odd? 215 deg(x) XE {ZEV! andeg(Z) is even } V = {YEV: deg(y) is odd} U {ZEV: deg(z) is even} even = odd teven This is a contradiction we cannot create an even number by Summing odd number and an even number Therefore, the initial claim is true.