

Ex: $L = \{a^n b^n c^n \mid n \geq 0\}$ is not a CFL.

Assume L is a CFL.

Then L must follow pumping lemma.
Suppose pumping length is p .

$$S = a^p b^p c^p, |S| = 3p \geq p$$

Then by pumping lemma, s must be able to be decomposed into 5 components where $S = uvxyz$, s.t

1. $uv^i xy^i z \in L, \forall i \geq 0$
2. $|vxy| > 0$
3. $|vxy| \leq p$

$S = aaaa \dots a bbbb \dots b cccc \dots c$

Diagram illustrating the decomposition of S into u, v, x, y, z :

- u consists of p 'a's.
- v and x consist of p 'b's.
- y is empty (not explicitly drawn).
- z consists of p 'c's.

Annotations:

- $p \#$ of a 's
- $p \#$ of b 's
- $p \#$ of c 's

case 1: vxy contains the same set of characters. Then v and y can only be allocated to the same type of character. Therefore, when we pump $uv^ix^iy^iz$, the # of one type of characters increases compared to other two types of characters,

$$\therefore uv^ix^iy^iz \notin L$$

case 2: vxy contains mix of characters.

Subcase 2.1, v and y contain the same character. We handled this in case 1.

ex: $vxy = aaa__a bbb__b$
 $v = aaa_a \quad y = aaa__a$

Subcase 2.2 v and y contain two different character types.

ex: $vxy = bbb__b ccc__c$
 $v = bbb_b \quad y = ccc__c$

Now, when we pump up, $uv^ix^iy^iz$ contains different number of type 3 character.

$$\therefore nv^ix^iy^iz \notin L$$

Subcase 2.3 either V or γ contains a mix of characters.

$$Vxy = aaa \dots a bbb \dots b$$

$$V = aa \dots ab \dots b \quad \gamma = bb \dots b$$

Now, when we pump up, uv^ixy^iz contains different order of a , b , and c 's.
 $\therefore uv^ixy^iz \notin L$

Now, we considered all the possible decompositions for S , and none of them are pumpable.

$\therefore L$ is not a CFL.

$C = \{ww \mid w \in \{0,1\}^*\}$ is not a CFL.

Assume C is a CFL.

Then C must follow pumping lemma.
Suppose pumping length is p .

Let $s = 0^p 1 0^p 1$ ← not good because it
 $\underbrace{\quad}_{p \# \text{ of } 0's}$ is pumpable
 $s = \underbrace{0000}_{n} \dots \underbrace{000}_{v} 1 \underbrace{0000}_{x} \dots \underbrace{0}_{z}$
 $uv^i xy^i z = 0^{p-3} 0^{3i} 1 0^{3i} 0^{p-3} 1$

Pick $s = 0^p 1^p 0^p 1^p$ $0 \dots 010 \dots 0 | 0 \dots 010 \dots 0$
 $s = 000 \dots 0 111 \dots 1 \quad | 000 \dots 0 (11\dots)$
 $\underbrace{\quad}_{vxz}$

$$|uv^2xy^2z| = 4p + |vy|$$

middle point would be $2p + \frac{|vy|}{2}$

↑
This is
pointing to 1 .

Since $|s| \geq p$, s must be able to be decomposed into $s = uvxyz$ s.t

1. $uv^ixy^iz \in C, \forall i \geq 0$
2. $|vy| > 0$
3. $|vxy| \leq p$

case 1: vxy does not straddle the mid point of s.

In this case $uv^2x^2y^2z \notin C$.

The reason is that, in $uv^2x^2y^2z$, the first character of string 1 and first character of string 2 does not match.

case 2: If vxy cross mid point, consider $uv^i x y^j z = u x z = 0^p 1^i 0^j 1^p$, where $i, j \neq p$.

$\therefore uv^i x y^j z \notin C$.

$\therefore C$ is not a CFL.



prove $L = \{a^n b^j \mid n = j^2\}$ is not a CFL.

Assume L is a CFL.

Then, let p be the pumping length.

Let $s = a^p b^p$

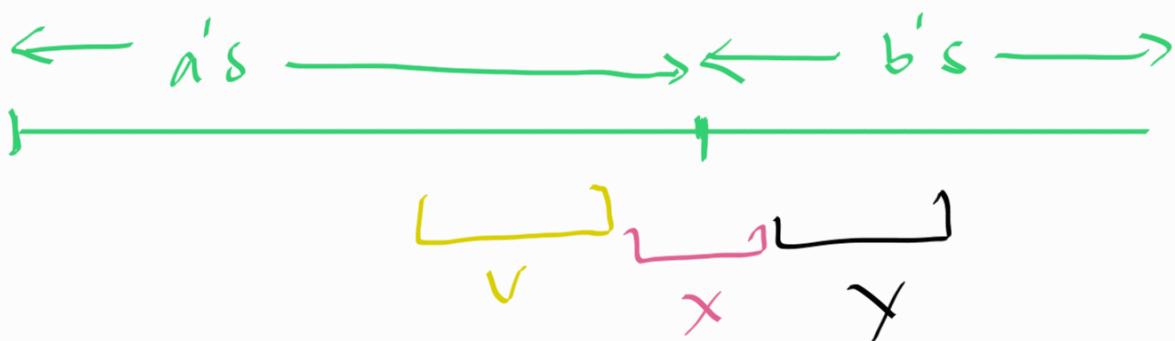
$s = aaaa \dots a bbbb \dots b$

$\underbrace{\qquad\qquad\qquad}_{p^2 \# a's} \qquad \qquad \qquad \underbrace{\qquad\qquad\qquad}_{p \# b's}$

case 1: Either v or y crosses the breakpoint. Thus, when pumping up, we will have out of order a 's and b 's.

$\therefore uv^i xy^i z \notin L$

case 2': Both v and y does not cross the break point.



k_1 # of a's to v

k_2 # of b's to y

Consider $uv^ix^iy^iz$.

$$|uv^ix^iy^iz| = p^2 + p + (i-1)k_1 + (i-1)k_2$$

subcase 2.1 $(k_1=0) \wedge (k_2=0)$

This violates condition a.

subcase 2.2 $(k_1=0) \vee (k_2=0)$

then $uv^0x^0y^0z \notin L$ since this would imbalance the # of a's and # b's.

$$\therefore uv^0x^0y^0z \notin L.$$

subcase 2.3 $(k_1 \neq 0) \wedge (k_2 \neq 0)$

consider $uv^i xy^i z$.

Now pumpdown, $i = 0$

$$|uv^0xy^0z| = p^2 + p - k_1 - k_2$$

Suppose $uxz \in L$

$$p^2 - k_1 = (p - k_2)^2$$

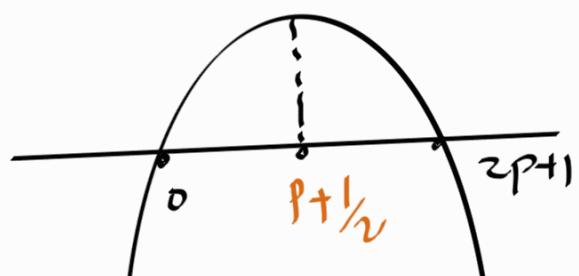
$$p^2 - k_1 = p^2 - 2pk_2 + k_2^2$$

$$k_1 = -k_2^2 + 2pk_2$$

I am going to consider $|vxy|$

$$\begin{aligned} |vxy| &= k_1 + k_2 + |x| \\ &= \underbrace{-k_2^2 + 2pk_2}_{f(k_2)} + k_2 + |x| \end{aligned}$$

$$f(k_2) = -k_2^2 + k_2(2p+1) = -k_2(k_2 - (2p+1))$$



$$f'(k_2) = -2k_2 + (2p+1)$$

$$f'(k_2) = 0 = -2k_2 + 2p+1, \quad k_2 = \frac{2p+1}{2} \\ = p + \frac{1}{2}$$

This means from $k_2=1$ to p $f(k_2)$ increases.

$$f(k_2) = -k_2^2 + k_2(2p+1)$$

$$f(1) = -1 + 2p + 1 = 2p$$

$$f(p) = -p^2 + p(2p+1) = -p^2 + 2p^2 + p \\ = p^2 + p$$

Here, we show $f(k_2)$ which is the length of $uvxyz$

$$\therefore \forall k_2 \in [1 \dots p] : f(k_2) > p$$

∴ There is no decomposition of $uvxyz$ which is pumpable

∴ L is not a cFL.

