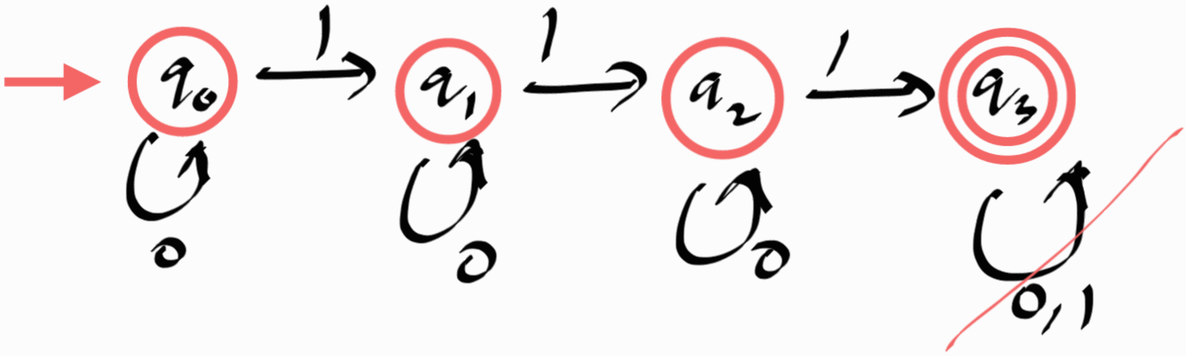
01/31/2025

clarifications.

Q1: $A = \{w \mid w \text{ contains at least } 3 \text{ 1s}\}$ $\Sigma = \{0,1\}$



111 ✓

11111 × ← not accepted

What happens if we do not have this transition?

Q2: $\underline{A \circ B} = \{xy \mid x \in A, y \in B\}$

How important is the order xy?

Let $\underline{A = \{0^*\}}$ $\quad$ $\underline{B = \{1^*\}}$

$\underline{00001111}$, $001111 \in A \circ B$

$111111 \in A \circ B$ why 0^0

But what happens

$\underbrace{1111}\,\underbrace{0000} \notin A \circ B$

$1111000 \in B \circ A$

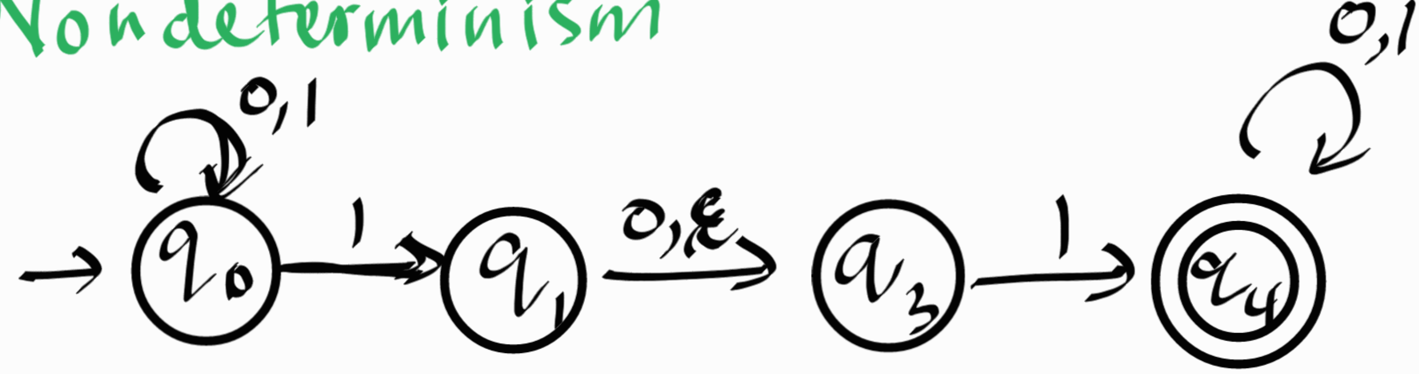
$A \circ B \neq B \circ A$ when $A \neq B$

Let's do some examples

Next week

Deterministic - Given the current state and the input, we know exactly which state to reach to.

Nondeterminism



Language?

$(0|1)^* 1 (0|\varepsilon) 1 (0|1)^*$

$(0|1)^* 1 0 1 (0|1)^*$

$(0|1)^* 1 1 (0|1)^*$

10

$q_0 \quad q_1$

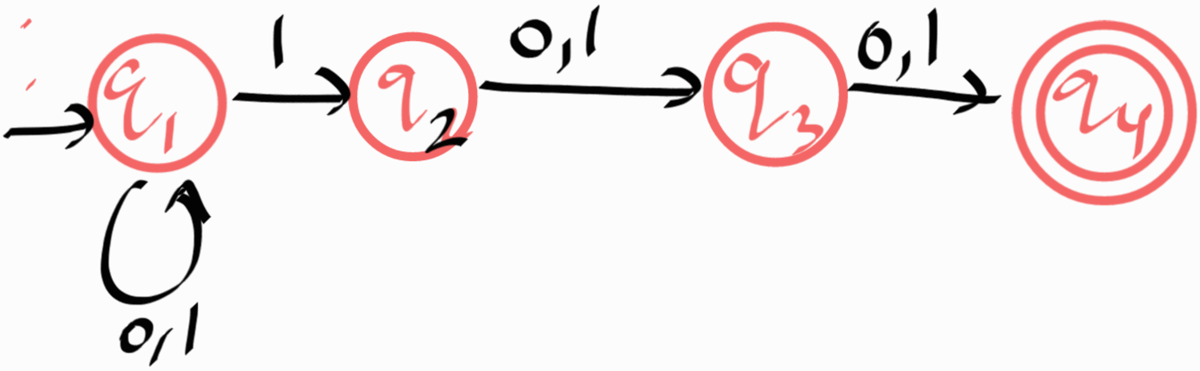
$\downarrow \quad \downarrow$

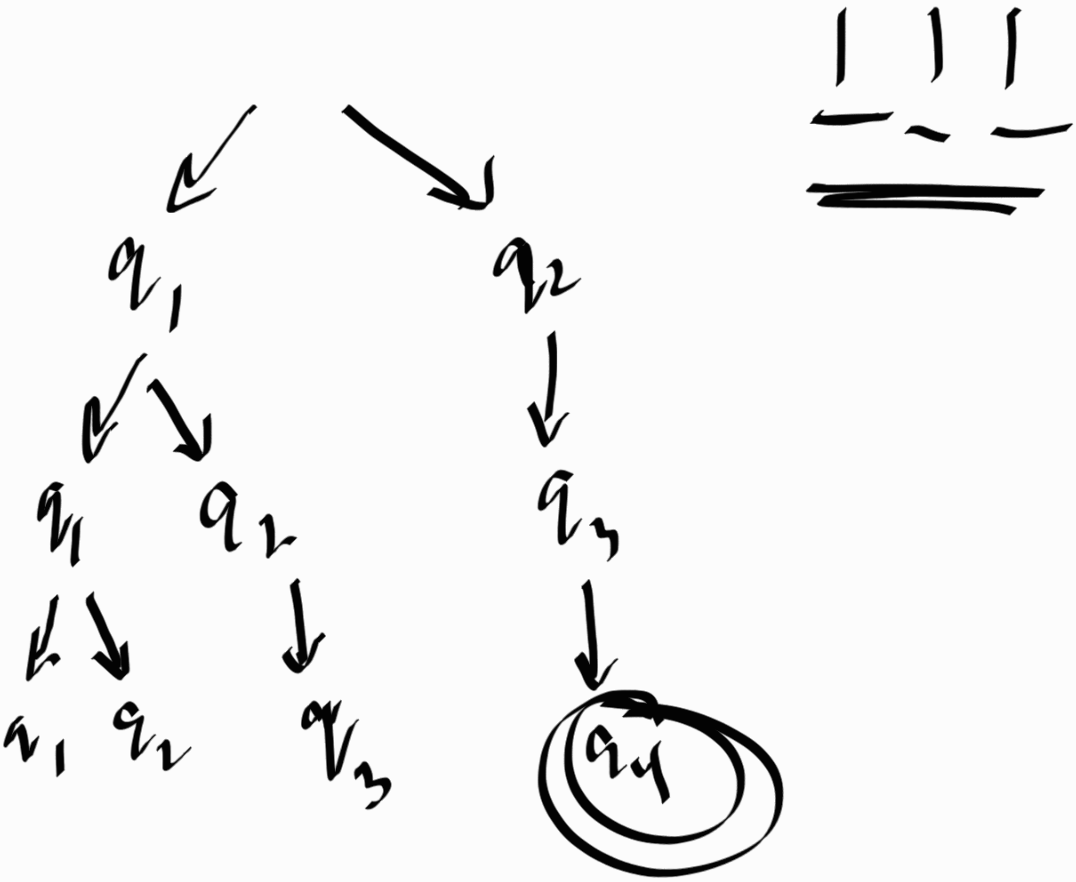
$q_0 \quad q_3$

How does NFA compute?

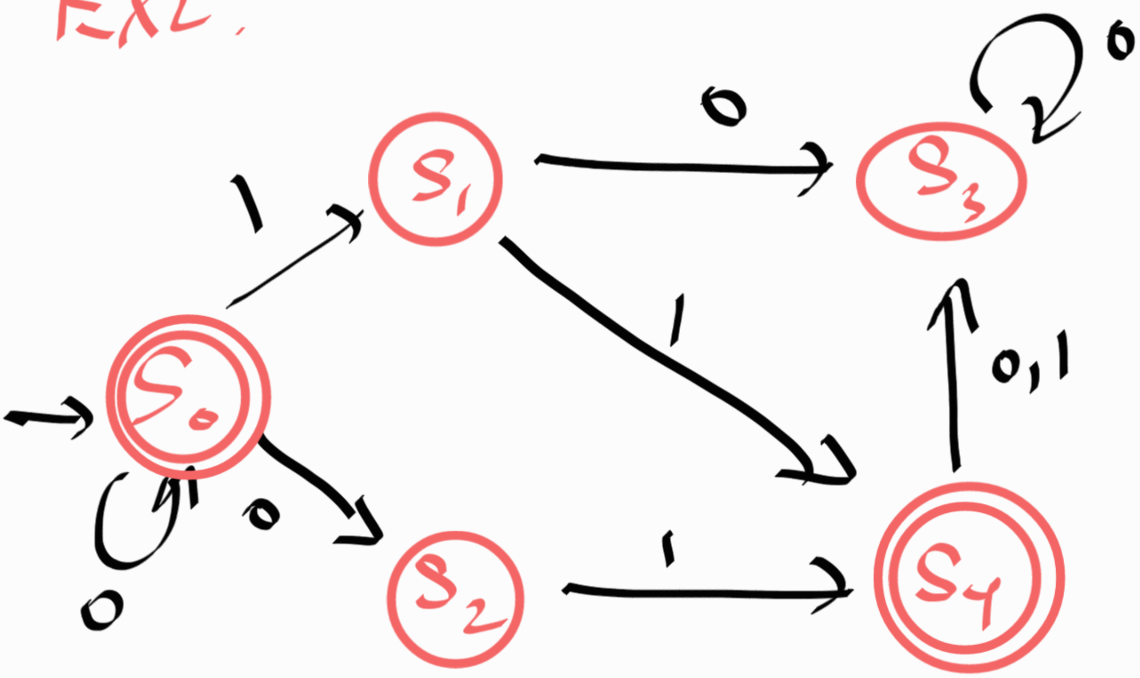
- If any of the accept state can be reached by reading the input x in any way, we say NFA accepts x.

Ex:





Ex2:



0^* ✓

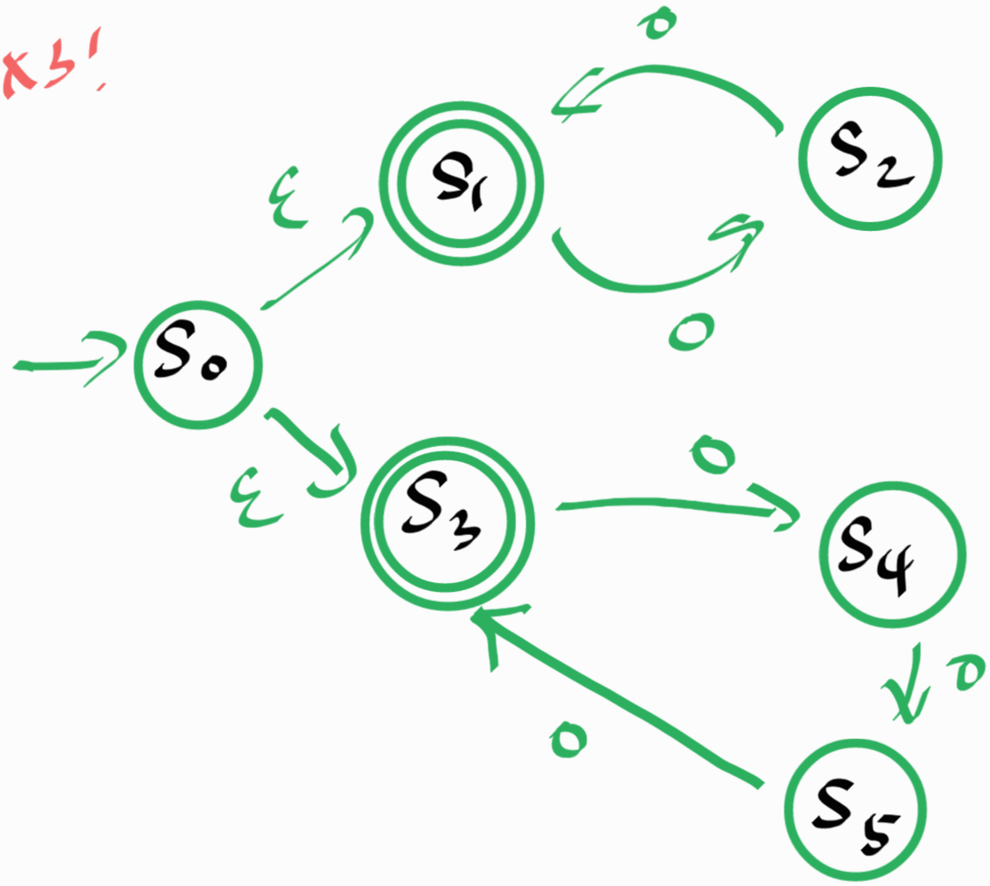
$0^* 1 0$ ✗

$0^* 1 1$ ✓

$0^* 1 1 (0|1)$ ✗

$0^* 0 1$ ✓

Ex3!



Proper definition of NFA

is a 5-tuple

$(Q, \Sigma, \delta, q_0, F)$, where

01. Q is the set of states.
02. Σ is a finite alphabet.
03. $\delta : Q \times \Sigma_\varepsilon \longrightarrow P(Q)$

$$\Sigma_\varepsilon = \Sigma \cup \{\varepsilon\}$$

$P(Q)$ = power set of Q

04. $q_0 \in Q$ is the starting state.
05. $F \subseteq Q$ is the accepting state.

Ex: $\delta(q_1, \varepsilon) = \{q_2, q_3, q_4\}$

$\delta(q_2, 1) = \{q_4\}$

Let $N = (Q, \Sigma, \delta, q_0, F)$ be an NFA, and $w = \langle w_1, w_2, w_3, \ldots, w_n \rangle$

$w_i \in \Sigma$

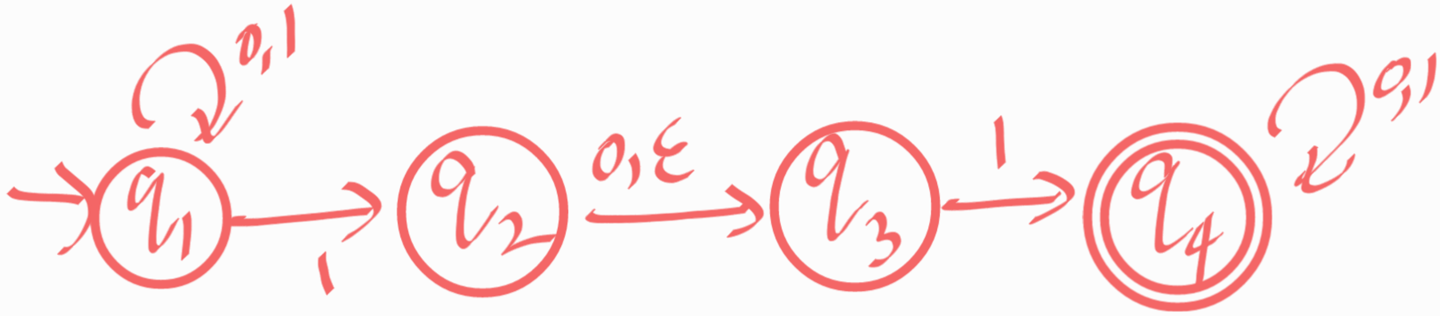
N accepts w if

$\exists r_0, r_1, r_2, \ldots\ldots, r_n \in Q$ s.t

1. $r_0 = q_0$
2. $r_{i+1} \in \delta(r_i, w_{i+1})$, $\forall i \in [0 \ldots n-1]$
3. $r_n \in F$

$$r_0 \xrightarrow{w_1} r_1 \xrightarrow{w_2} r_2 \xrightarrow{w_3} \cdots \xrightarrow{w_n} r_n$$

TWO MACHINES ARE EQUIVALENT if they accept the same language



1. $Q = \{q_1, q_2, q_3\}$
2. $\Sigma = \{0,1\}$, $\Sigma_\varepsilon = \{0,1,\varepsilon\}$
3. δ is given as follows

$\{\} = \emptyset$

	ε	0	1	
q_1				
q_2				
q_3				
q_4				

4. $q_1 \in Q$ is the start state.
5. $\{q_4\} \subseteq Q$ is the accepting states