

Example 04:

$$F = \{ww \mid w \in \{0,1\}^*\}$$

Show that F is not regular.

What if we pick $0^p 0^p$ as the string to prove F is not regular?

$$\Sigma = \{0,1\}$$

00 01 01 111 111

01 0101 X

$$s = \underline{\underline{0}}^p 0^p \in F \quad \checkmark$$

$$|s| = 2p \geq p$$

$$s = xyz \quad |y| > 0 \quad \underline{|xy| \leq p}$$

$$s = \begin{matrix} p-a & a & p \\ \overbrace{0}^x & \overbrace{0}^y & \overbrace{0}^z \\ x & y & z \end{matrix} \quad a \in \mathbb{N} \quad a > 0 \quad p - a + a = p$$

let a be a odd number

$$xy^az = 0^{p-a} 0^p = 0^{2p-a} \notin F$$

Let \underline{a} be a even number

$$\begin{aligned} xy^iz &= 0^{p-a} 0^a 0^p = 0^{p-a+a+i+p} \\ &= 0^{2p+a(i-1)} \\ &= 0 \end{aligned}$$

$2p+a(i-1)$ is even

because $2p$ & a is even.

Proof:

$$S = \underbrace{0^P}_1 \underbrace{0^P}_1 \in F$$
$$|S| = 2P+2 \geq P$$

Assume F is regular, then S must follow pumping lemma.

Then S can be decomposed into 3 components $S = xyz$ s.t

$$\textcircled{1} \quad xyz \in F$$

$$\textcircled{2} \quad |y| > 0$$

$$\textcircled{3} \quad |xy| \leq P$$

$$x = 0^{P-a} \quad y = 0^a$$

By $\textcircled{3}$, y must contain only 0's (before the first 1)

$$\text{Then, } xy^a z = 0^{P-a} 1 0^a 1, a > 0, P \neq a + p$$

Therefore $xy^a z \notin F$, Thus F is not regular.

Ex 5!

$$\Sigma = \{1\}$$

$\mathcal{D} = \left\{ 1^{n^2} \mid n \geq 0 \right\}$ is not regular.

$$n=0$$

ϵ

$$n=1$$

1

$$n=2$$

1111

$$n=3$$

11111111

$$s = 1^{p^2} \quad |s| = p^2 \geq p$$

$$1^{(p+1)^2} \swarrow$$

$$(p+1)^2 = p^2 + 2p + 1$$

$$s = \underbrace{1111111 \dots}_{p \text{ number of } 1's} \quad \overbrace{\dots \dots \dots}^{p^2-p \text{ number of } 1's}$$

$$p + p^2 - p = p^2$$

$$x = \epsilon \quad y = 1^p \quad z = 1^{p^2-p}$$

$$xy^2z = \epsilon \cdot 1^p \cdot 1^p \cdot 1^{p^2-p}$$

$$|xy^2z| = p^2 - p + 2p = p^2 + p =$$

Proof:

Assume D is regular.

Since D is regular D should follow Pumping lemma.

Suppose pumping length of D is p .

Select $s = 1^{p^2}$, $|s| = p^2 \geq p$

Since D is regular and $|s| \geq p$,
s should be able to decompose
into 3 strings $s = xyz$ s.t

1. $xyz \in D$ for $i \geq 0$

2. $|y| > 0$

3. $|xy| \leq p$

By ③, $|xy| \leq p$, therefore $|y| \leq p$

By ② $|y| > 0$ or $|y| \geq 1$

- we have

$$|x y^2 z| = |xyz| + |y| = p^2 + |y|$$

$$p^2 < |x y^2 z| = p^2 + |y| \leq p^2 + p \leq p^2 + 2p + 1$$

$$p^2 < |x y^2 z| < (p+1)^2$$

Basically this means even if we allocate max number of '1's to y , we still cannot generate the next string $1^{(p+1)^2}$, where

$$|1|^{(p+1)^2} = (p+1)^2.$$

Therefore, $xy^2z \notin D$, a contradiction.

$\therefore D$ is not regular.

1. Deterministic finite automata
2. Non-Deterministic finite automata
3. Regular expressions.

End of Regular languages.

Next, we are going to look at a more powerful computation model.