

# CSCI 347 Data Mining

## Graph Data

# Graph Data

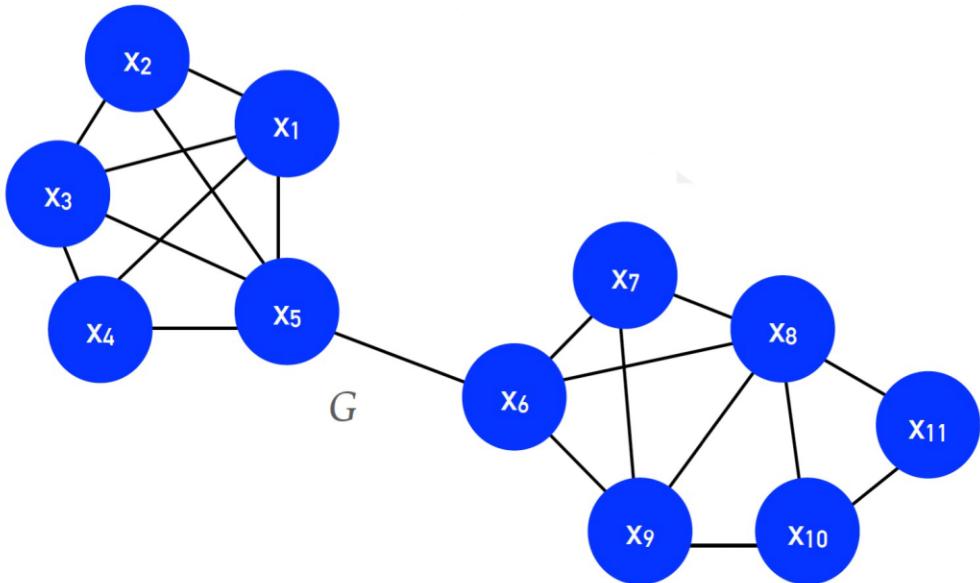
- Data instances are often not entirely independent
- they can be interconnected through various types of relationships.
- Graph data or networks are a data structure where instances are depicted as **nodes**, and the connections between these instances are represented by **edges**.

# Graph Data

$$G = (V, E)$$

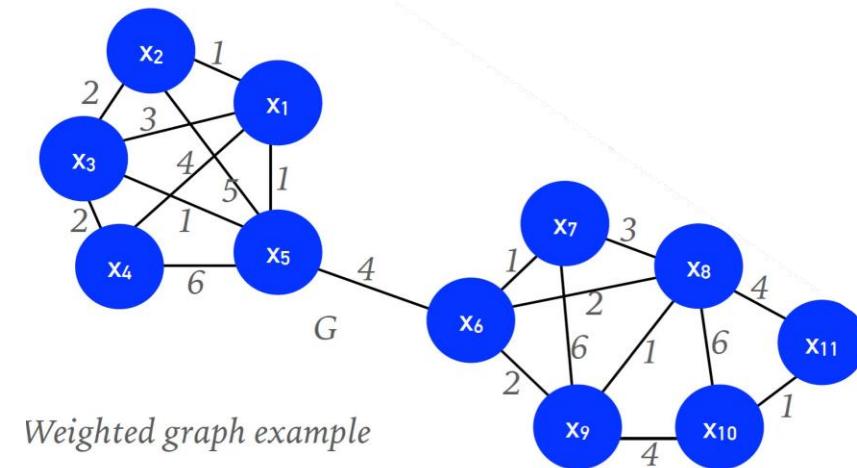
$V$  = set of vertices

$$E \subseteq \{\{u, v\} : u, v \in V\}$$



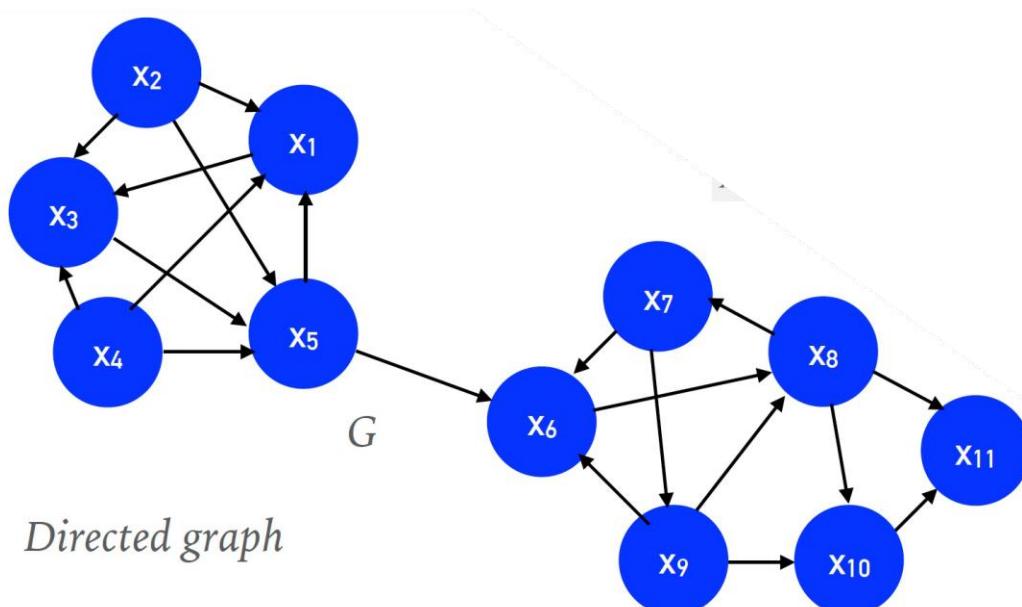
# Graph Data (Weighted graph)

- $G = (V, E, w)$
- $V$  = Vertices or Nodes
- $E$  = Unordered pairs of vertices with weights ( $w_{ij}$ )
- $w : E \rightarrow \mathbb{R}^+$  (usually positive real values)



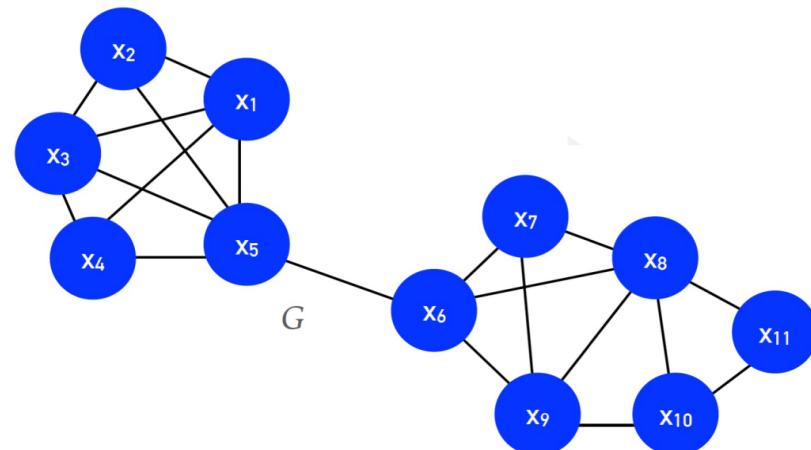
# Graph Data (Directed Graph)

- $G = (V, E)$
- $V$  = Vertices or Nodes
- $E \subseteq V \times V$  is the **ordered** pairs of vertices.



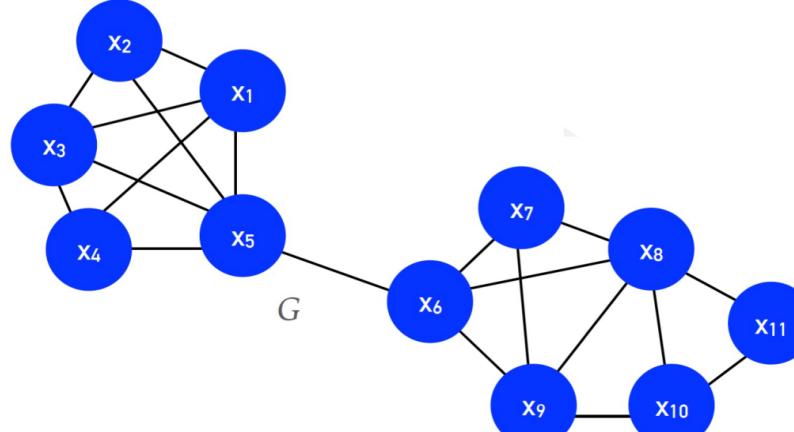
# Graph Data

- $G = (V, E)$
- $V = \text{Vertices or Nodes}$
- $E = \text{Unordered pairs of vertices}$
- Simple graph = Undirected graph without loops
- Edge,  $e = (v_i, v_j)$ ,  $v_i$  and  $v_j$  are adjacent or neighbors.
- Order:  $|V| = n$ , Size:  $|E| = m$



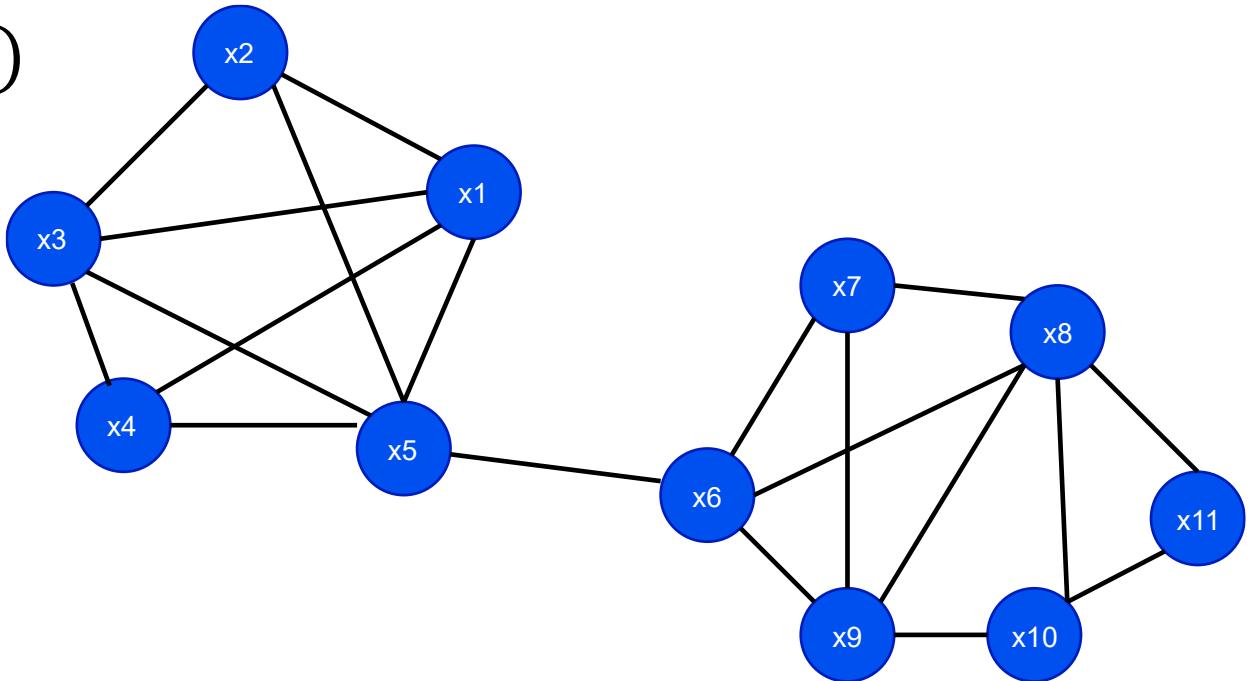
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- **Order**:  $|V| = n$ , **Size**:  $|E| = m$
- A graph  $H = (V_H, E_H)$  is called a subgraph of  $G = (V, E)$ , if  $V_H \subseteq V$  and  $E_H \subseteq E$  and the endpoints of edges in  $E_H$  is in  $V_H$ .



# Degree of a node

- The degree of a node  $v_i \in V$  is the number of edges incident with it and is denoted as  $d(v_i)$  or just  $d_i$ .

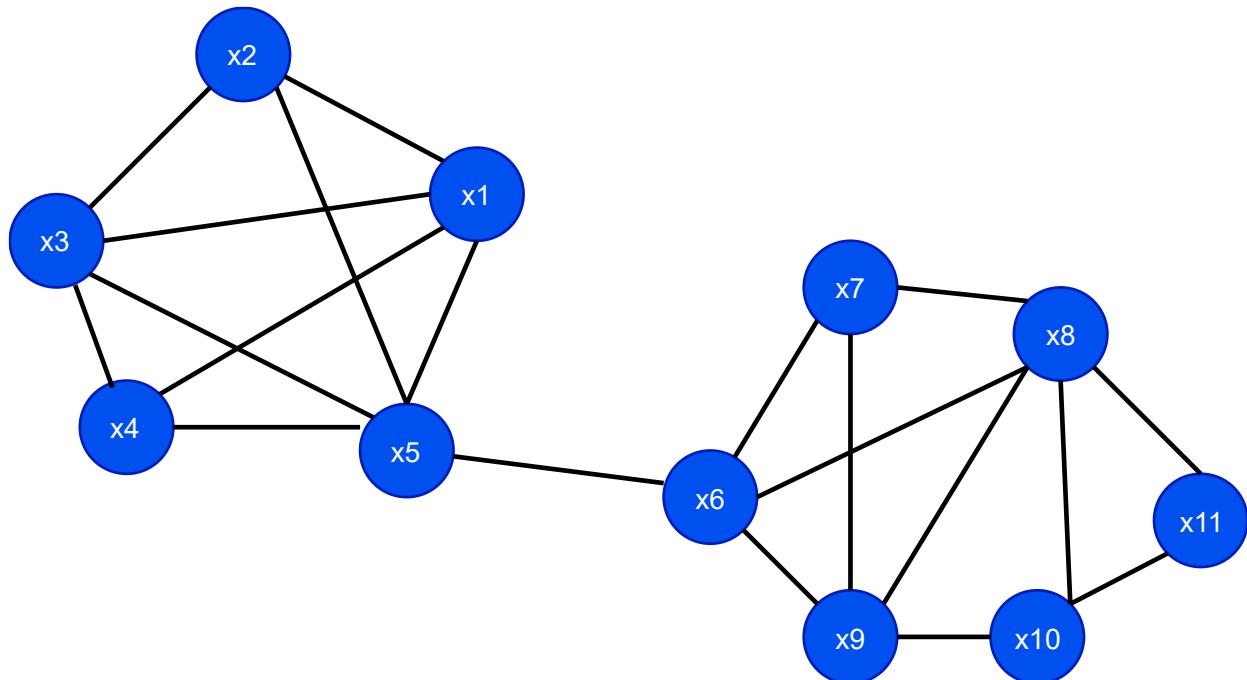


# Degree of a node

- The degree of a node  $v_i \in V$  is the number of edges incident with it and is denoted as  $d(v_i)$  or just  $d_i$ .

What is the degree of  $x_9$ ?

1. 3
2. 1
3. 4
4. 8



# Degree Distribution

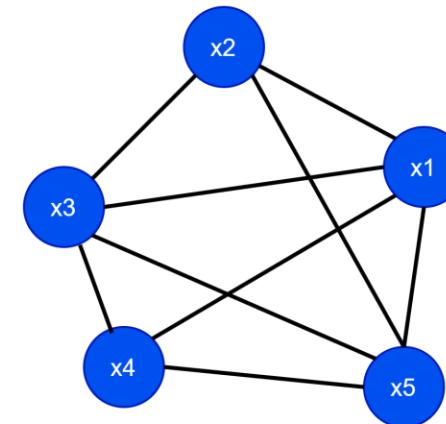
- Let  $N_k$  denote the number of vertices with degree  $k$ . The degree frequency distribution of a graph is given as  $(N_0, N_1, \dots, N_t)$ 
  - $t$  is the maximum degree of a node in the graph.

# Degree Distribution

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  - $t$  is the maximum degree of a node in the graph.

What is the degree distribution of this graph?

1.  $(4, 3, 4, 3, 4)$
2.  $(3, 3, 4, 4, 4)$
3.  $(0, 0, 0, 2, 3)$
4.  $(2, 3)$



# Degree Distribution

- The probability that a given node is of degree  $k$  is  $\frac{N_k}{n}$ .
- Suppose you have a random process of picking a node in a graph, and random variable  $X$  that assigns the degree of the picked node.

$$P(X = k) = \frac{N_k}{n}, n \text{ is the number of nodes}$$

Note: A random variable is a measurable function that assigns numerical values to outcomes of a random experiment.

$$X : \Omega \rightarrow \mathbb{R}, \quad \Omega \text{ is the sample space}$$

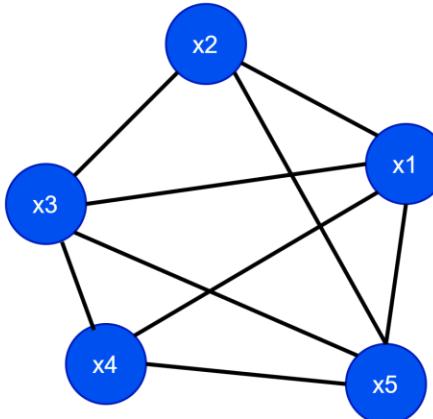
# Degree Distribution

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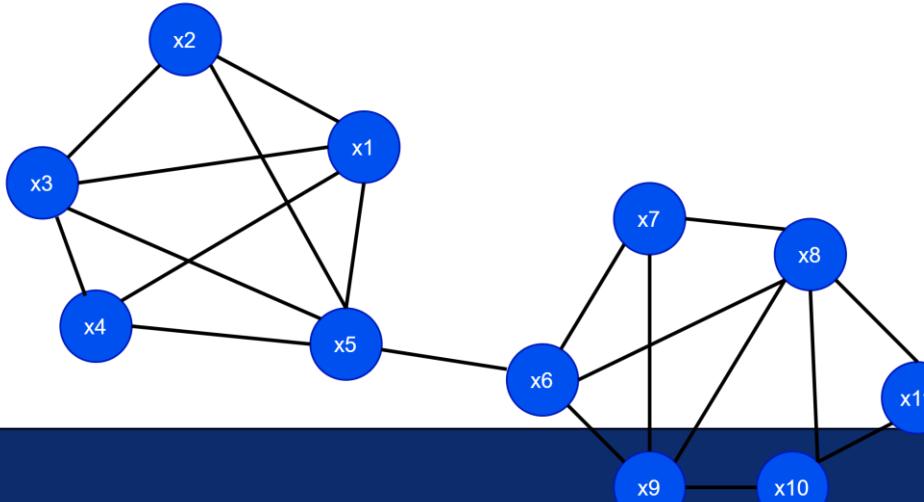
- Given the node distribution (0,0,0,2,3) what is the probability that a node is of degree 3?

1. 0
2. 2
3.  $3/5$
4. 3
5.  $2/5$
6. None of the above



# Walk, Path, shortest path

- A **walk** in a graph  $G$  between nodes  $x$  and  $y$  is an ordered sequence of vertices, starting at  $x$  and ending at  $y$ .  
 $Walk := < v_0, v_1, \dots, v_t >, v_0 = x, v_t = y, \forall i \in [0..t-1]: (v_i, v_{i+1}) \text{ exists}$
- **The length of the walk**  $t$ , is the number of edges along the walk.
- A **path** is a walk with distinct vertices.
- A path of minimum length between nodes  $x$  and  $y$  is called a **shortest path**.

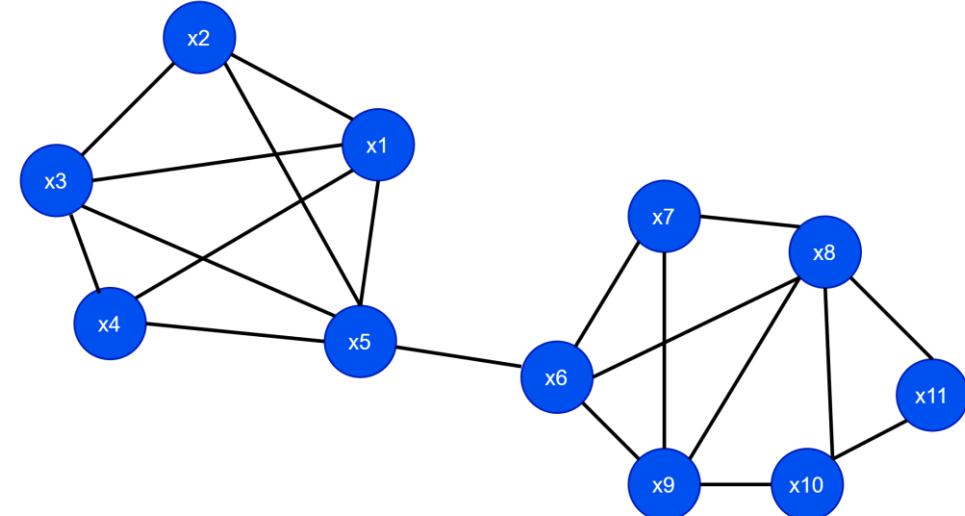


# Walk, Path, shortest path

- A path of minimum length between nodes  $x$  and  $y$  is called a **shortest path**.

What is the length of the shortest path between  $x_2$  and  $x_{10}$ ?

1. 6
2. 3
3. 4
4. 1
5. 0

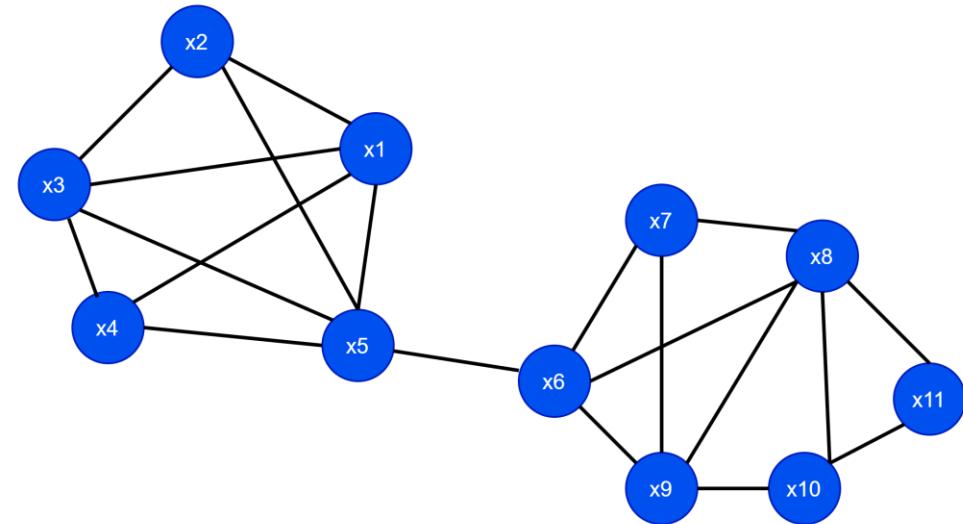


# Connectedness

- Two nodes  $v_i$  and  $v_j$  are said to be **connected** if there exists a **path** between them.
- A graph is **connected** if there is a path between all **pairs of vertices**.
- A **connected component**, or just **component**, of a graph is a **maximal connected subgraph**.
  - **maximal** means that the subgraph cannot be extended any further while still maintaining the property of being connected.

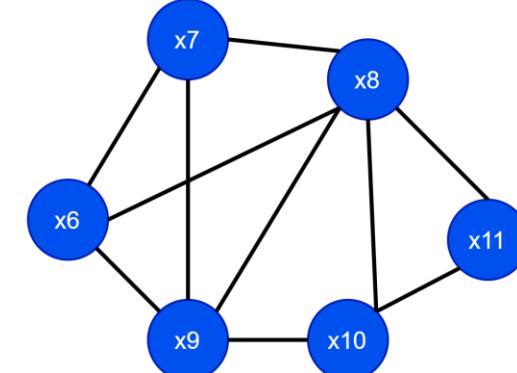
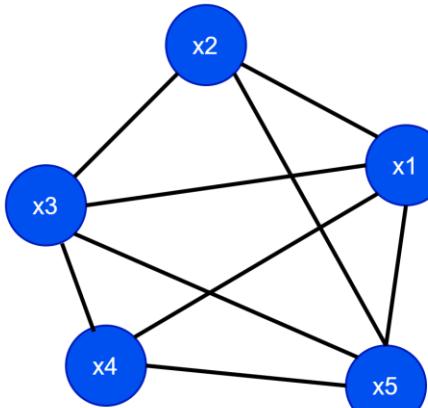
# Connectedness

- Is this graph connected?
  - Yes
  - No



# Connectedness

- Is this graph connected?
  - Yes
  - No



# Adjacency Matrix

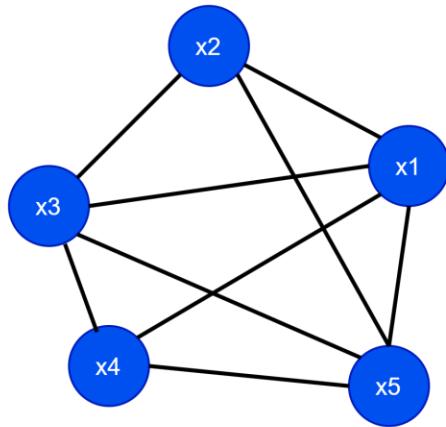
- A graph  $G = (V, E)$ , with  $|V| = n$  vertices, can be conveniently represented in the form of an  $n \times n$ , symmetric binary adjacency matrix,  $A$ , defined as:

$$A(i, j) = \begin{cases} 1 & \text{if } v_i \text{ is adjacent to } v_j \\ 0 & \text{otherwise} \end{cases}$$

- A weighted graph can be represented by  $n \times n$  weighted adjacency matrix.

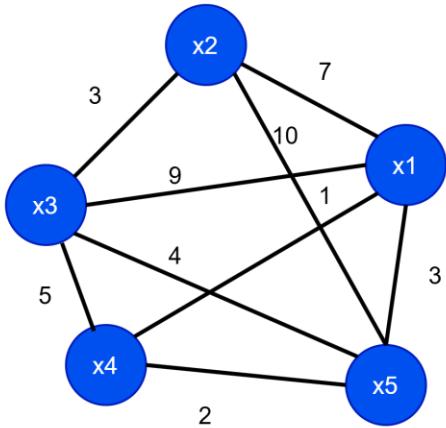
$$A(i, j) = \begin{cases} w_{ij} & \text{if } v_i \text{ is adjacent to } v_j \\ 0 & \text{otherwise} \end{cases}$$

# Adjacency matrix example.



	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$
$x_1$	0	1	1	1	1
$x_2$	1	0	1	0	1
$x_3$	1	1	0	1	1
$x_4$	1	0	1	0	1
$x_5$	1	1	1	1	0

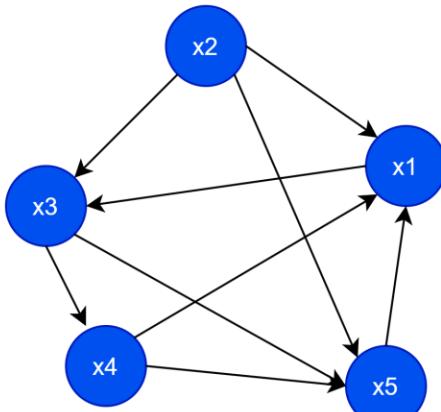
# Adjacency matrix (weighted) example.



	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$
$x_1$	0	7	9	1	3
$x_2$	7	0	3	0	10
$x_3$	9	3	0	3	4
$x_4$	1	7	5	0	2
$x_5$	3	10	4	2	0

# Adjacency matrix: directed graph

- In a **directed graph** adjacency matrix is **not symmetric**.



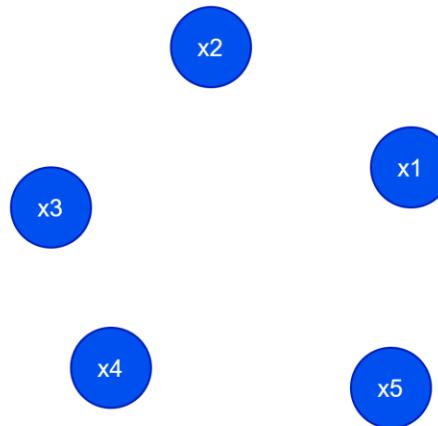
	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$
$x_1$	0	0	1	0	0
$x_2$	1	0	1	0	1
$x_3$	0	0	0	1	1
$x_4$	1	0	0	0	1
$x_5$	1	0	0	0	0

# Graphs from Data Matrix

- Given a dataset in the form of a matrix, can we create a graph?

	$X_1$	$X_2$	$X_3$
$x_1$	0.2	1	12.3
$x_2$	1.3	4	89.23
$x_3$	5.6	5	56.1
$x_4$	4.5	7	47.3
$x_5$	7.3	12	45.23

?



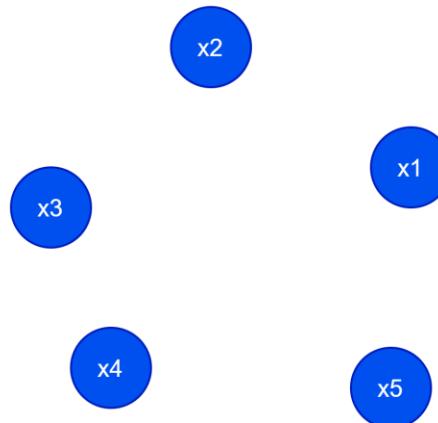
But what can we do about the edges?

# Graphs from Data Matrix

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	$X_1$	$X_2$	$X_3$
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$x_5$	7.3	12	45.23

?



How about we use a similarity measure  
and then use the similarity measure as the  
edge weights?

# How to create a graph from matrix?

- Define a weighted graph  $G = (V, E)$ .

$V = \{v_i \mid v_i \text{ represents the entity } x_i\}$

$$w_{ij} = sim(x_i, x_j)$$
$$sim(x_i, x_j)$$

*represents the similarity between points  $x_i$  and  $x_j$*

## Gaussian similarity

$$w_{ij} = sim(x_i, x_j) = e^{-\frac{\|x_i - x_j\|^2}{2\sigma^2}}$$

$\sigma$  is the spread parameter.

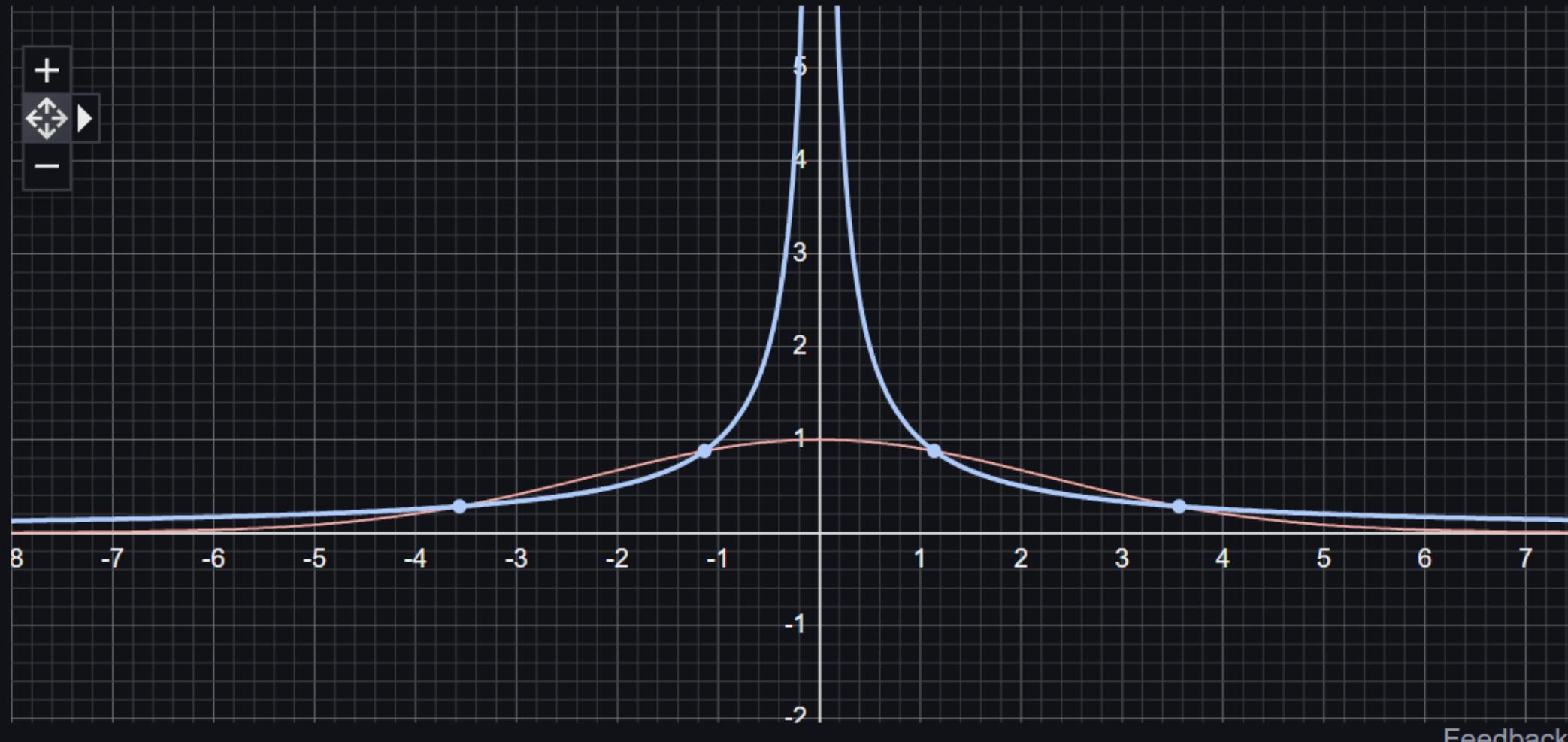
# Gaussian similarity

- Similarity is defined as being inversely related to the Euclidean distance.
- If two vectors are far apart, then we say it's less similar.
  - Therefore, we can put lower weight between them.
- But why do we use this?
  - We can use something like  $\frac{1}{\|x_i - x_j\|}$

# Gaussian similarity

- Exponential Decay
  - The similarity measure  $w_{ij}$  decays smoothly and asymptotically to 0 as the distance increases, ensuring that distant points contribute very little but not abruptly.
  - It is bounded between 0 and 1, which simplifies interpretation and normalization in algorithms.
- Inverse distance
  - $\frac{1}{\|x_i - x_j\|}$  decays too slowly as the distance increases, leading to non-negligible contributions from distant points.
  - It has an unbounded range  $(0, \infty)$ , which can create numerical instability and make it harder to interpret.

# Graph for $1/|x|$ , $e^{-|x|^2/10}$



# Gaussian similarity

- Handling zero distance:
  - When handling  $\|x_i - x_j\| = 0$ ,  $w_{ij}$  simplifies to  $e^0 = 1$ , but if we use  $\frac{1}{\|x_i - x_j\|}$ , then mathematically this is not defined.
- Sensitivity control
  - The parameter  $\sigma$  allows you to control the sensitivity to distance.
    - Smaller  $\sigma$ , similarity decays quickly.
    - Larger  $\sigma$ , similarity decays slowly.
    - We can tweak the graph by changing this parameter.
  - $\frac{1}{\|x_i - x_j\|}$  does not have this property.

# Gaussian similarity

- Robustness to outliers.
  - Gaussian similarity drops off quickly for large distances, effectively ignoring the outliers.
  - $\frac{1}{\|x_i - x_j\|}$ , even if far away, can have disproportionately large effect to slow decay of  $\frac{1}{d}$ .

# Graphs from Data Matrix

- Gaussian similarity with  $\sigma = 25$

$$e^{-\frac{\|x_i - x_j\|^2}{2\sigma^2}}$$

	$X_1$	$X_2$	$X_3$
$x_1$	0.2	1	12.3
$x_2$	1.3	4	89.23
$x_3$	5.6	5	56.1
$x_4$	4.5	7	47.3
$x_5$	7.3	12	45.23

	x1	x2	x3	x4	x5
x1	0	0.008709	0.207862	0.359302	0.366178
x2	0.008709	0	0.409967	0.241611	0.196716
x3	0.207862	0.409967	0	0.936019	0.87281
x4	0.359302	0.241611	0.936019	0	0.970737
x5	0.366178	0.196716	0.87281	0.970737	0

# Graphs from Data Matrix

- Gaussian similarity with  $\sigma = 50$

$$e^{-\frac{\|x_i - x_j\|^2}{2\sigma^2}}$$

	$X_1$	$X_2$	$X_3$
$x_1$	0.2	1	12.3
$x_2$	1.3	4	89.23
$x_3$	5.6	5	56.1
$x_4$	4.5	7	47.3
$x_5$	7.3	12	45.23

	x1	x2	x3	x4	x5
x1	0	0.30549	0.675218	0.774221	0.777899
x2	0.30549	0	0.800179	0.701098	0.665978
x3	0.675218	0.800179	0	0.983606	0.966562
x4	0.774221	0.701098	0.983606	0	0.992603
x5	0.777899	0.665978	0.966562	0.992603	0

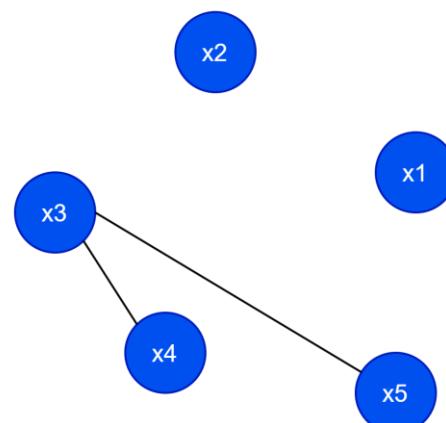
# Creating a graph from matrix

	$X_1$	$X_2$	$X_3$
$x_1$	0.2	1	12.3
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$x_3$	5.6	5	56.1
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	x1	x2	x3	x4	x5
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x3	0.675218	0.800179	0	0.983606	0.966562
x4	0.774221	0.701098	0.983606	0	0.992603
x5	0.777899	0.665978	0.966562	0.992603	0

$$\tau = 0.94$$

	x1	x2	x3	x4	x5
x1	0	0	0	0	0
x2	0	0	0	0	0
x3	0	0	0	1	1
x4	0	0	1	0	1
x5	0	0	1	1	0



$$sim(x_i, x_j) = e^{-\frac{\|x_i - x_j\|^2}{2\sigma^2}}$$

$$A(i, j) = \begin{cases} 1 & \text{if } sim(x_i, x_j) \geq \tau \\ 0 & \text{otherwise} \end{cases}$$

# Iris Similarity Graph: Gaussian Similarity

$\sigma = \frac{1}{\sqrt{2}}$ , edge exist if and only if  $w_{ij} \geq 0.777$

Order:  $|V| = n = 150$ , size:  $|E| = m = 753$

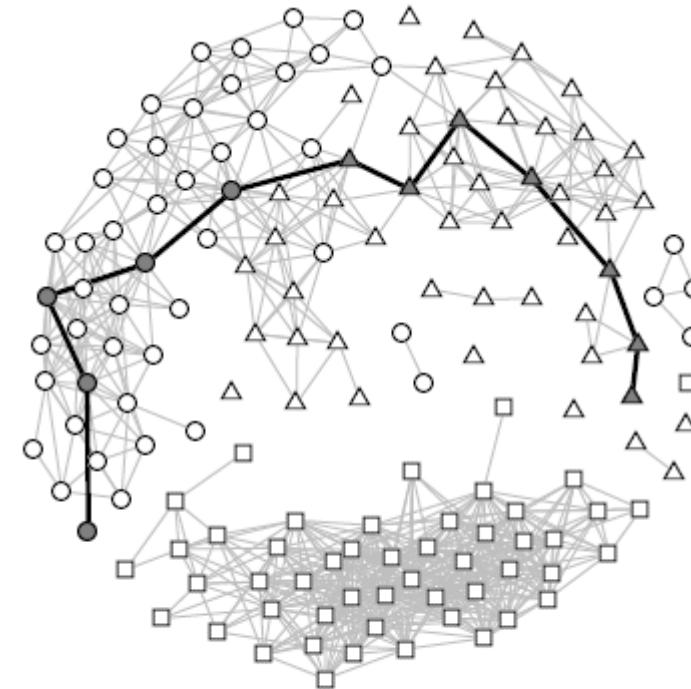


Figure 4.2: Iris Similarity Graph