

LP RIC

Randomized Incremental d -Dimensional Linear Programming

Input: A set $H = \{h_1, \dots, h_n\}$ of $(d - 1)$ -dimensional halfspaces, such that the first d define an initial feasible vertex v_d , and the objective vector c .

Output: The optimum vertex v or an error status indicating that the LP is infeasible.

- (1) If the dimension is 1, solve the LP by brute force in $O(n)$ time.
- (2) Let v_d be the intersection point of the hyperplanes bounding h_1, \dots, h_d , which we assume define an initial feasible vertex. Randomly permute the remaining halfspaces, and let $\langle h_{d+1}, \dots, h_n \rangle$ denote the resulting sequence.
- (3) For $i = d + 1$ to n do:
 - (a) If $(v_{i-1} \in h_i)$ then $v_i \leftarrow v_{i-1}$.
 - (b) Otherwise, intersect $\{h_1, \dots, h_{i-1}\}$ with the $(d - 1)$ -dimensional hyperplane ℓ_i that bounds h_i and project onto \mathbb{R}^{d-1} . Let c' be the projection of c onto ℓ_i and then onto \mathbb{R}^{d-1} . Solve the resulting $(d - 1)$ -dimensional LP recursively.
 - (i) If the $(d - 1)$ -dimensional LP is infeasible, terminate and report that the LP is infeasible.
 - (ii) Otherwise, let v_i be the solution to the $(d - 1)$ -dimensional LP.
- (4) Return v_n as the final solution.

Analysis for randomized LP!

Let $T_d(n)$ be the expected run time of RIC LP alg

we show by induction on n
that $T_d(n) \leq \delta_d d! n$

(\curvearrowleft)

wy δ_d is dependent on the dim

for $i \in \{d+1, \dots, n\}$ then we remove dim dependency

let p_i be the prob that h_i changed the opt vertex

~~IN ℓ_i the
 v_i~~

Case 1: prob ($1 - p_i$) that there is no change
 $\Rightarrow O(d)$ time to determine

case 2: prob p_i that there is a change in opt vertex

- 1) project C onto ℓ_i - $O(d)$ time
 - 2) intersect $i-1$ halfspaces w/ ℓ_i - $O(d(i-1))$ time
 - 3) recurse to solve $(d-1)$ -dim LP on $i-1$ halfspaces
 - by induction hyp it takes $T_{d-1}(i-1)$ time
- takes $O(d)$ time

cobine (and omit constant factors)

$$\begin{aligned} T_d(n) &\leq \sum_{i=d+1}^n \left((1-p_i)d + p_i (d + T_{d-1}(i)) \right) \\ &\leq \sum_{i=d+1}^n \left(d + p_i (d + T_{d-1}(i)) \right) \end{aligned}$$

What is p_i ? ... Backwards analysis

Let S_i be an arbitrary set of i hyperplanes from the input
how many permutations $\rightarrow i!$

Want to count: Number of permutations in which the opt vert changes
on the i^{th} step?

Let v_i be the opt vertex of i -halfspaces
... Note v_i depends only on S_i (not the order)

Observations:

- there are d halfspaces that define v_i
- if none of them are $h_i \Rightarrow v_{i-1} = v_i$
- if any of them are $h_i \Rightarrow$ opt vertex changed

\Rightarrow opt vert changed \Leftrightarrow definer of v_i is h_i

Since any halfspace could be last
and there are d halfspaces that can define v_i

$$\Rightarrow p_i = \frac{d}{i}$$

plug our p_i into our bound for $T_d(n)$

$$T_d(n) \leq \sum_{i=d}^n (d + p_i (d^i + T_{d-1}(i)))$$

assume $T_{d-1}(i) = \gamma_{d-1} (d-1)! i$ (by induction hyp.) side note

$$\leq \sum_{i=d}^n \left(d + \frac{d}{i} (d^i + \gamma_{d-1} (d-1)! i) \right)$$

$$\leq \sum_{i=d+1}^n (d + d^2 + \gamma_{d-1} d!) \leq (d + d^2 + \gamma_d d!) n$$

Find a const γ_d so that
 $\gamma_d d! n \geq d + d^2 + \gamma_{d-1} d!$

$$\text{choose } \gamma_d = \frac{d+d^2}{d!} + \gamma_{d-1}$$

$$\Rightarrow T_d(n) = (d + d^2 + \gamma_{d-1} d!) n = \left(\frac{d+d^2}{d!} + \gamma_{d-1} \right) d! n \leq \gamma_d d! n$$

What about our const
for large enough d

$$\frac{d+d^2}{d!} \leq \frac{1}{2^d} \quad \text{and} \quad \sum_{i=1}^{\infty} \frac{1}{2^i} = \sum_{i=1}^{\infty} \left(\frac{1}{2}\right)^i \quad \text{is geometric series}$$



$$\Rightarrow T_d(n) = O(d! n)$$

since d is const $d!$ is a const

wy $r < 1$ (e.g. $\frac{1}{2}$)

\Rightarrow geometric series converges to a const !!

So sidel: the prob of time as 6-times the expected running time
is $O\left(\frac{1}{c} e^{6d!}\right)$ for any fixed const c

e.g. in 2^d 10x expected running time is at most $6.5e^{-12}$