Other M/R Programming Interfaces

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PIG

- An alternative way to program analysis of large data sets that is declarative and ad-hoc
- Declarative programming is a programming paradigm that expresses the logic of a computation without describing its control flow.
 - in contrast to imperative programming which expresses an algorithm
- Ad-hoc: (adj)
 - Formed for or concerned with one specific purpose: an ad hoc compensation committee.
 - Improvised and often impromptu: "On an ad hoc basis, Congress has . . .



SQL (an aside)

SQL is declarative

- Users express queries using first order logic
- System evaluates queries in different fashions depending upon the size and cardinality of the data, types, constraints, etc.
- The program places no constraints on execution

SQL is ad hoc

- Allows arbitrary inspection, reuse of data, based on logical queries. The design of the database places no restriction on its future use.
- This is not true of most data structures, e.g. priority heaps support only insert and retrieve greatest



So why not SQL?

- Well, there is SQL for Map/Reduce. That's HIVEQL.
- But, we'll look at PIG first.



C. Olston *et al.* Pig Latin: A Not-So-Foreign Language for Data Processing. SIGMOD, 2008.

- Databases too expensive and don't conform to how programmers think
- Map/reduce is too low level

pensive at this scale. Besides, many of the people who analyze this data are entrenched procedural programmers, who find the declarative, SQL style to be unnatural. The success of the more procedural *map-reduce* programming model, and its associated scalable implementations on commodity hardware, is evidence of the above. However, the map-reduce paradigm is too low-level and rigid, and leads to a great deal of custom user code that is hard to maintain, and reuse.

We describe a new language called *Pig Latin* that we have designed to fit in a sweet spot between the declarative style of SQL, and the low-level, procedural style of map-reduce.



More on what's wrong with M/R

- Practical efforts at MR programs end up with multiphase confusing programs.
 - Need concepts such as joins

Unfortunately, the map-reduce model has its own set of limitations. Its one-input, two-stage data flow is extremely rigid. To perform tasks having a different data flow, e.g., joins or n stages, inelegant workarounds have to be devised. Also, custom code has to be written for even the most common operations, e.g., projection and filtering. These factors lead to code that is difficult to reuse and maintain, and in which the semantics of the analysis task are obscured. Moreover, the opaque nature of the map and reduce functions impedes the ability of the system to perform optimizations.



SQL versus PIG

- Example: find the average page rank of highly ranked pages in a big categories
- SQL is declarative

```
SELECT category, AVG(pagerank) FROM urls WHERE pagerank > 0.2 GROUP BY category HAVING COUNT(*) > 10^6
```

- PIG describes a computation as a sequence of steps
 - Imperative style programming



Dataflow

- PIG programs are akin to describing a dataflow execution plan
 - And the people they likey:

"I much prefer writing in Pig [Latin] versus SQL. The step-by-step method of creating a program in Pig [Latin] is much cleaner and simpler to use than the single block method of SQL. It is easier to keep track of what your variables are, and where you are in the process of analyzing your data." – Jasmine Novak, Engineer, Yahoo!



CREATE TEMPORARY TABLE rbAUtmp select rbm.Matter_Identifier, gpm.GP_PAM_Matter_Identifier, mbs.GP_PAM_Attorney_Identifier, count(mbs.GP_PAM_Attorney_Identifier) as cnt, mbs.GP_PAM_Country_Identifier FROM dbo_Matter_Billing_Statistics AS mbs JOIN dbo_Matter_GP_PAM_Matter_Identifier AS gpm ON (mbs.GP_PAM_Matter_Identifier=gpm.GP_PAM_Matter_Identifier AND mbs.GP_PAM_Country_Identifier=gpm.GP_PAM_Country_Identifier) JOIN rb_Matter as rbm ON (rbm.Matter_Identifier=gpm.Matter_Identifier) GROUP BY mbs.GP_PAM_Attorney_Identifier ORDER BY Matter_Identifier, cnt DESC:

set @num := 0, @mid := ";

CREATE TABLE rbMatAtt SELECT t1.Matter_Identifier, t2.Attorney_Identifier, t1.cnt, @num:=if(@mid=Matter_Identifier, @num + 1, 1) AS row_number, @mid:=Matter_Identifier AS dummy FROM rbAUtmp AS t1 LEFT OUTER JOIN dbo_Attorney_GP_PAM_Attorney_Identifier AS t2 ON (t1.GP_PAM_Attorney_Identifier=t2.GP_PAM_Attorney_Identifier AND t1.GP_PAM_Country_Identifier) ORDER BY t1.Matter_Identifier, t1.cnt DESC:

ALTER TABLE rbMatAtt ADD PRIMARY KEY (Matter_Identifier,row_number);
CREATE TEMPORARY TABLE rbATT1 SELECT Matter_Identifier, Attorney_Identifier AS
ATT1_ID, Attorney_AVP_Salary_USD as ATT1_Salary, Attorney_Seniority AS ATT1_Seniority,
Attorney_PLS as ATT1_Specialty, Attorney_Location as ATT1_Location FROM rbMatAtt LEFT
OUTER JOIN rbdp Attorney USING (Attorney Identifier) WHERE row number=1;

.

CREATE TEMPORARY TABLE rbATT1j SELECT m.Matter_Identifier, ATT1_ID, ATT1_Salary, ATT1_Seniority, ATT1_Specialty, ATT1_Location FROM rb_Matter AS m LEFT OUTER JOIN rbATT1 AS a1 USING (Matter_Identifier);

.

CREATE TABLE rbATTs SELECT * FROM rbATT1j JOIN rbATT2j USING (Matter_Identifier) JOIN rbATT3j USING (Matter_Identifier);



PIG Language Constructs

- FOREACH: allows parallel processing (in a mapper) for all inputs in a data set
- FILTER: discard unwanted data (in either mapper or reducer)
- GROUP/CO-GROUP: put related data together using the shuffle process.
- These constructs allow for database-style query optimization.



PIG Data Model

- Atom: simple value
- Tuple: sequence of values
- Bag: multiset with duplicates
 - flexible schema for elements

- Map: key/value data structure
 - Keys must be atoms for efficiency

$$\left[\begin{array}{c} \texttt{'fan of'} \to \left\{\begin{array}{c} \texttt{('lakers')} \\ \texttt{('iPod')} \end{array}\right\} \\ \texttt{'age'} \to 20 \end{array}\right]$$



PIG Expressions

Simple set that have to be parallelizable

$$\texttt{t = \left('alice', \left\{\begin{array}{c} ('lakers', 1) \\ ('iPod', 2) \end{array}\right\}, \left['age' \rightarrow 20\right]\right)}$$

Let fields of tuple t be called f1, f2, f3

Expression Type	Example	Value for t
Constant	'bob'	Independent of t
Field by position	\$0	'alice'
Field by name	f3	$\left["age" o 20" ight]$
Projection	f2.\$0	<pre>{ ('lakers') ('iPod') }</pre>
Map Lookup	f3#'age'	20
Function Evaluation	SUM(f2.\$1)	1 + 2 = 3
Conditional Expression	f3#'age'>18? 'adult':'minor'	'adult'
Flattening	FLATTEN(f2)	'lakers', 1 'iPod', 2

Table 1: Expressions in Pig Latin.



Bags and Co-Groupings

- Pig programming uses the patten of co-grouping data, applying aggregates, and then flattening the results
 - Allows SQL like functionality in sequenced programming
 - It's not super-intuitive (see the paper)

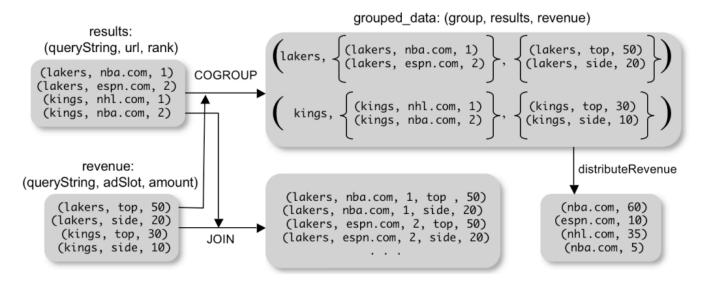




Figure 2: COGROUP versus JOIN.

Compiliing to MR

 Each PIG program compiles to several MR programs and is run in Hadoop!

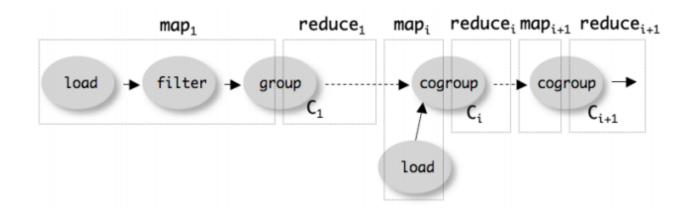


Figure 3: Map-reduce compilation of Pig Latin.



Compiling to MR (ii)

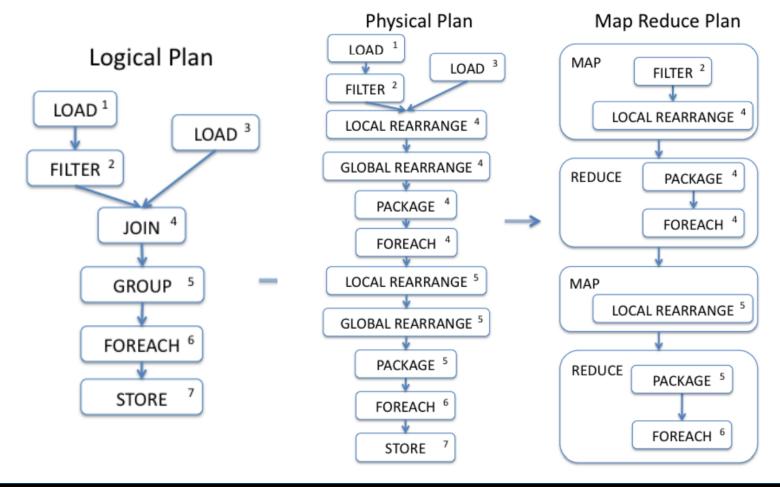
• From: Gates et al. Building a High-Level Dataflow System on top of Map-Reduce: The Pig Experience, VLDB 2009.

Logical Plan Pig Latin LOAD (x, y, z) LOAD (t, u, v) A = LOAD 'file1' AS(x, y, z); B = LOAD 'file2' AS (t, u, v); FILTER (x, y, z) C = FILTER A by y > 0;D = JOIN C BY x, B BY u;JOIN (x, y, z, t, u, v) E = GROUPDBYz;F = FOREACH E GENERATE group, COUNT(D); GROUP (group, {(x, y, z, t, u, v)}) STORE F INTO 'output'; **FOREACH** (group, count) **STORE** (group, count)



Compiling to MR (ii)

• From: Gates et al. Building a High-Level Dataflow System on top of Map-Reduce: The Pig Experience, VLDB 2009.





Hive

- Hive: data model and system for data warehousing in map/reduce systems.
- HiveQL: SQL programming for Map/Reduce
 - Not SQL 92 complete
 - No transactions, no materialized views, limited subquery support
- The definitive Hive paper
 - Thusoo et al. Hive A Warehousing Solution Over a Map-Reduce Framework. PVLDB, 2009.



Hive Example: Status Meme

Table schema:

```
status_updates(userid int, status string, ds string)
```

Load log files daily:

```
LOAD DATA LOCAL INPATH '/logs/status_updates'
INTO TABLE status_updates PARTITION (ds='2009-03-20')
```



Daily Statistics

 Join logs with profiles and figure out the number of tweets from men/women and by school

```
FROM (SELECT a.status, b.school, b.gender
FROM status_updates a JOIN profiles b
ON (a.userid = b.userid and
a.ds='2009-03-20')
) subq1
INSERT OVERWRITE TABLE gender_summary
PARTITION(ds='2009-03-20')
SELECT subq1.gender, COUNT(1) GROUP BY subq1.gender
INSERT OVERWRITE TABLE school_summary
PARTITION(ds='2009-03-20')
SELECT subq1.school, COUNT(1) GROUP BY subq1.school
```



How do it go?

 Hive puts tables on HDFS as files and runs queries as Hadoop! jobs

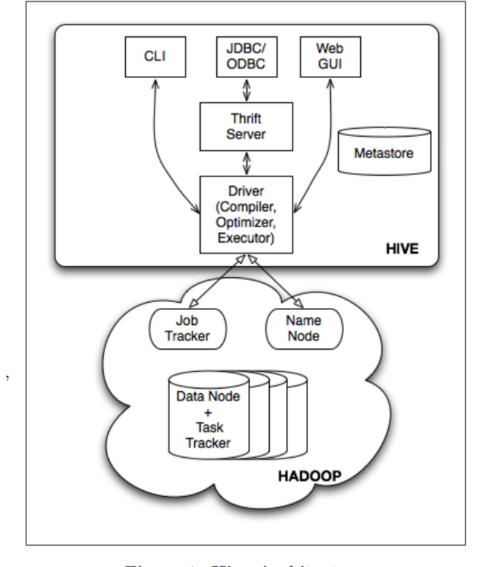
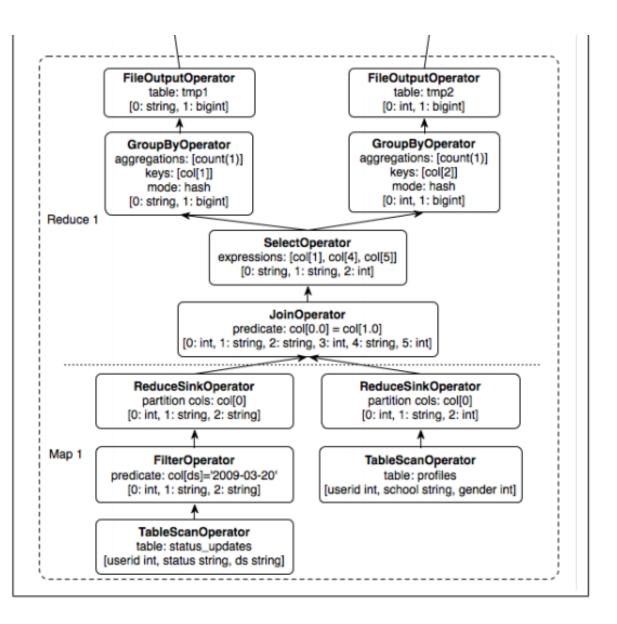


Figure 1: Hive Architecture



Resulting Query Plan (part 1)

 You don't need to understand. These are the MR jobs generated by the example





Resulting Query Plan (part 2)

```
FileOutputOperator
                                                                 FileOutputOperator
              table: school_summary
                                                                table: gender summary
                [0: string, 1: bigint]
                                                                    [0: int, 1: bigint]
                 SelectOperator 5 2 2
                                                                   SelectOperator
            expressions: [col[0], col[1]]
                                                              expressions: [col[0], col[1]]
                [0: string, 1: bigint]
                                                                   [0: int, 1: bigint]
Reduce 2
                                                                                           Reduce 3
                GroupByOperator
                                                                  GroupByOperator
             aggregations: [count(1)]
                                                                aggregations: [count(1)]
                   keys: [col[0]]
                                                                     keys: [col[2]]
               mode: mergepartial
                                                                  mode: mergepartial
                [0: string, 1: bigint]
                                                                    [0: int, 1: bigint]
              ReduceSinkOperator
                                                                ReduceSinkOperator
               partition cols: col[0]
                                                                 partition cols: col[0]
                [0: string, 1: bigint]
                                                                    [0: int, 1: bigint]
 Map 2
                                                                                           Map 3
               TableScanOperator
                                                                 TableScanOperator
                    table: tmp1
                                                                      table: tmp2
                [0: string, 1: bigint]
                                                                    [0: int, 1: bigint]
                 FileOutputOperator
                                                               FileOutputOperator
                     table: tmp1
                                                                    table: tmp2
```

```
FROM (SELECT a.status, b.school, b.gender
FROM status_updates a JOIN profiles b
ON (a.userid = b.userid and
a.ds='2009-03-20')
) subq1
INSERT OVERWRITE TABLE gender_summary
PARTITION(ds='2009-03-20')
SELECT subq1.gender, COUNT(1) GROUP BY subq1.gender
INSERT OVERWRITE TABLE school_summary
PARTITION(ds='2009-03-20')
SELECT subq1.school, COUNT(1) GROUP BY subq1.school
```

Take Aways

- Other ways to program M/R
 - More concise, easier to maintain
 - Particularly for data processing tasks that result in mutli-stage map/reduce programs
- Ethos: take the best from DBs
 - Declarative languages and optimization
 - Ad-hoc queries
- Ethos: and leave behind the stuff that's not parallel
 - Indexes, nested sub-queries



The (Un)Reasonable Debate

- Imperative programming
 - How humans think, step by step
 - Program encodes execution instructions
- Declarative programing
 - What! (Not how.)
 - Allows system to optimize execution
 - Non-intuitive (for many)
 - SQL != declarative programming. It is a specific instance that some love and some hate.
- PIG notable for trying to strike a happy balance
 - DB guys don't see the upside here

