A function dependency A → B means for all instances of a particular value of A, there is the same value of B.

Α	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b2	c3	d2	e3
a3	b3	c4	d2	e4
a4	b4	c5	d4	e5

 A function dependency A → B means for all instances of a particular value of A, there is the same value of B.

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a1	<u>b1</u>	c1	d1	e1
a2	b2	c2	d2	e2
a1	<u>b1</u>	c3	d2	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

**Function dependency** 

# 1.What about B→A ???????

Α	В	С	D	E
<u>a1</u>	b1	c1	d1	e1
a2	b2	c2	d2	e2
<u>a1</u>	b1	c3	d2	e3
<u>a3</u>	b1	c4	d2	e4
a4	b4	c5	d4	e5

1.What about

C→A ????????

1.What about

C→D ????????

1.What about

**D**→A ???????

A	В	С	D	E
a1	b1	<b>c1</b>	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d2	e3
a3	b1	c4	d2	e4
a4	b4	<b>c</b> 5	d4	e5

**Function dependency** 1.What about C→A holds 1.What about  $C \rightarrow D$  holds 1.What about  $D \rightarrow A$ does not hold

Α	В	C	D	E
a1	b1	c1	d1	e1
a2	b2	<u>c2</u>	d2	e2
a1	b1	c3	d2	e3
a3	b1	<u>c4</u>	d2	e4
a4	b4	c5	d4	e5

# Candidate Key

- Any attribute(attributes) that can uniquely identify a tuple in a relation
- According to functional dependency, any attribute (set of attributes) can functionally determine all other attributes in a relation, that attribute(attributes) is called candidate key

Function dependency
1.What about A??

A→B holds A→C does not hold

A can't be the candidate key

A	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d2	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

# Function dependency 1.What about B??

B→A does not hold

B can't be the candidate key

Α	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d2	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

# Function dependency 1.What about B??

### B→A does not hold

B can't be the candidate key

Α	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d3	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

### 1.What about C??

**C→A**???

C→B ???

 $C \rightarrow C ???$ 

**C**→**D**???

**C→E** ???

Α	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d3	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

### 1.What about C??

**C→A** ??? yes

 $C \rightarrow B ??? yes$ 

 $C \rightarrow C$  ??? yes

 $C \rightarrow D???$  yes

 $C \rightarrow E$  ???yes

C is a candidate key

A	В	С	D	E
a1	b1	<b>c1</b>	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d3	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

### 1.What about D??

**D**→A ??? **NO** 

**D**→**B** ??? **NO** 

 $D \rightarrow C$  ??? NO

 $D \rightarrow D$ ??? yes

 $D \rightarrow E$  ??? No

D is a not a candidate key

A	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d3	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

# Any other candidate key possible?????

We can take all two attribute combinations after excluding C& E

**Ie:AB??** 

**AD??** 

BD?

Α	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d3	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

**AB?????** 

AB-→C does not exist

So AB is not a key

Α	В	С	D	E
a1	b1	<b>c1</b>	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d3	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

**AD?????** 

 $AD \rightarrow A$ ?

 $AD \rightarrow B$ ?

 $AD \rightarrow C$ ?

 $AD \rightarrow D$ ?

 $AD \rightarrow E$ ?

Α	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d2	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

**AD?????** 

 $AD \rightarrow A? YES$ 

**AD-→B?YES** 

AD-→C?YES

AD-→D? YES

**AD-→E?YES** 

AD is a key

A	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d2	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

**BD?????** 

BD-→A? NO

BD-→B?YES

BD-→C?NO

BD-→D? YES

BD-→E?NO

**BD IS NOT A KEY** 

Α	В	С	D	E
a1	b1	c1	d1	e1
a2	b2	c2	d2	e2
a1	b1	c3	d2	e3
a3	b1	c4	d2	e4
a4	b4	c5	d4	e5

## Prime attributes

- Prime attribute-The constituent attributes of a relation are called prime attributes.
- Conversely, an attribute that does not occur in ANY candidate key is called a non-prime attribute.
- In the previous relation ,R(A,B,C,D,E)
   ,Candidate keys are {C,E,AD}
- So prime apptributes are {A,C,D,E}
- Non prime attribtes are {B}

## Armstrong's Axioms

Armstrong's Axioms: Let X, Y be sets of attributes from a relation T.

```
[1] Inclusion rule: If \underline{Y} \subseteq \underline{X}, then \underline{X} \to \underline{Y}.
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- [2] Transitivity rule: If  $\underline{X} \to \underline{Y}$ , and  $\underline{Y} \to \underline{Z}$ , then  $\underline{X} \to \underline{Z}$ .
- [3] Augmentation rule: If  $\underline{X} \rightarrow \underline{Y}$ , then  $\underline{XZ} \rightarrow \underline{YZ}$ .
- Other derived rules:
  - [1] Union rule: If  $\underline{X} \to \underline{Y}$  and  $\underline{X} \to \underline{Z}$ , then  $\underline{X} \to \underline{YZ}$
  - [2] Decomposition rule: If  $X \to YZ$ , then  $X \to Y$  and  $X \to Z$
  - [3] Pseudotransitivity: If  $\underline{X} \to \underline{Y}$  and  $\underline{WY} \to \underline{Z}$ , then  $\underline{XW} \to \underline{Z}$