## DISK SCHEDULING

Auess Time = Suk Time + Robational latency

Seek Time: The seek time is the time for the disk aim to more the heads to the yunder containing the desired sector.

Rotational eatency: The estational latency is the additional time for the disk to estate the desired sector to the disk head.

The disk bandwidth is the botal no of bytes transfined, divided by the total time blue the first request for service and the completion of the last transfer. We can improve both the access time and bandwidth by managing the order in which disk 1/0 requests are serviced.

a system and to the os. The request operaties:

- or whether this operation is input or output
- what the disk address for the transfer is
- I what the namony address for the transfu'is
- The desired disk drive and controller are anatolic, the request is suriced immediately observes any new request for surice were be placed in queue of pending requests for that drive. When request completed, os chooses pending request to surice next.

following an the algorithms used:

- 1. FCFS scheduling agouthm
- 2. 55 7F Sunduling Algorithm
- 3. SCAN Schoduling algorithm
- 4 C-SCAN Scheduling algorithm
- 5. LOOK Scheduling algorithm

of FCFS scheduling Algorithm

FCFS is the simplest disk scheduling algorithm. The algorithm the supplest disk scheduling algorithm. The algorithm the struction in the order truly online the the disk entertains enquests in the order truly online the the olisk quee. Appears to be fair and does not have stawation.

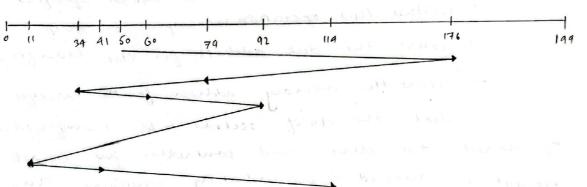
Algorithm

1. Let Request away represents an away storing indexes of tracks that have been requested in ascending order of their time of arrival. 'head' is the position of disk head.

2. Take one franks one by one in default order and calculate the absolute distance of the crack from the head.

- 3. Incument the total seek count with this distance
- 4. amonthy serviced track position now becomes new head post.
- 5. ho to sup 2 until all tracks in request anay have not been serviced.

eg: Sequence: [176, 79, 34, 60, 92, 11, 41, 114.]
Initial mad position: 50



Total seek count: (176-50) + (176-79) + (79-34) + (60-34) + (92-60) + (92-11) + (41-11) + (114-41) : 510

Bossie idea is the tracks which are closer to current disk head position should be serviced fust in order to nonimize the seek operations.

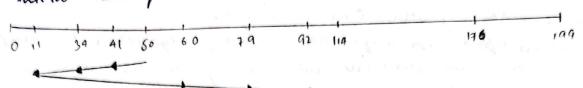
Angonithm:

1. Let Request away represents an away stowing incluses of exacts that have been requested, ineast is the position of deskhead find the positive distance of all tracks in the request away from head.

3. Find a track from the equested away which is not accessed / serviced get and was minimum distance

from head.

- total seek count with this distance. 4. Incument
- 5. convently serviced track position now becomes new headposts.
- 6. Go to step 2 until all tracks in request away have been not Swilld.
- eg: Seguma. [ 176, 79, 34, 60, 92, 11, 41, 114] Instrat head position: 50



Seek count: (50-41)+(41-34) + (34-11) +(60-11) +(79-60) +(92-79) + (114-92) + (176-114) = 204.

we sear sheduling Algorithm (Flevaror) the other end. servicing requests in 6/w one by one and wach me other end. Then direction of head is wrensed and the process continues as head continuously sun back and forth to ares the disk. Requests at midrange an surled more and those arriving behind the disk arm well have to wart.

Algorithm

- 1. Let request away upresents an away storing indexes of tracks that have been requested in ascending order of their
- 2. Let direction represents whether an head is moring towards left or right.
- 3. In the direction in which head is moving service all tracks one
- 4. Calculate Con absolute distance of the track from the head. 5. Incument the total seek count with this distance.
- country serviced track position non becomes new head poon.
- Cos to step 3 until we nach at one of the ends of the desk.
- If we wach at the ends of me disk
- If we reach the end of the disk severse the direction and goto depr untill all tracks in request away have not bun serviced.

Instal head position: 50 start direction: left

34 41 50 60 39 92 1A 136

Total seek count: (50- 41) + (41-36) + (34-11) + (11-0) + (60-0) + (79-60) + (176-114) = 226

O c-scan Algorithm (cincular (lunder)

→ modified vission of SCAN surduring algorithm that improves efficiency

→ more head from our and suricing all aquest to one and and
immediately uturns to beginning of disk althour surraing requests

on return tusp and starts again once maches beginning.

Algorithm

1. Invitalize request away and head.

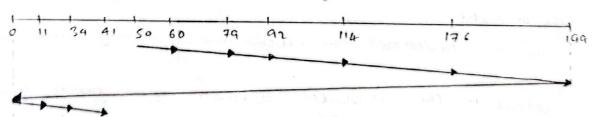
- ?. Head surius only in right direction from the size of disk
- 3. Which moving eff, no requests are surviced.

a. When we reach beginning, were duction.

- that moving light dherron calculate absolute distance of took from the head.
- 6. Indement total suk count and our cumunty suriced position as

7. Go to step 6 until we mach right end of disk.

- 8. If we wach right end, were direction and goto steps.
- G: sequena: 6176, 79, 34, 60, 97, 11, 41, 114} Initial head: 50 Direction: light.



Suk count- (60-50) + (79-60) + (92-79) + (114-92) + (176-114) + (199-176) + (199-0) + (11-0) + (34-11) + (41-34) = 389

Took Algorithm is back to reduce amount of remation of scanning fine to avers date on hard disk. Operates by scanning disk in a specific direction but instead of going all the way the end, it reverses direction as soon asit waches current objection.

digoutum

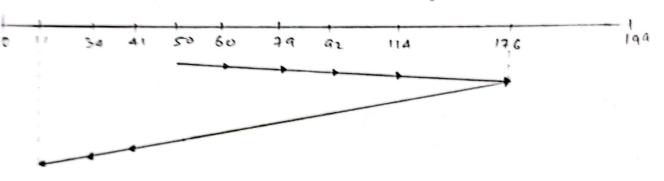
1. Desumine the inestial direction of disk head movement.

2-sort the personny dish requests in order to which they will be onvied.

3. San en disk in the Chosen distetion, servicing riquests

e when the last request is served twent obsertion and continue Scanning mutil all requests have been serviced.

Sequenu: \$176, 79, 34, 60, 92, 11, 41, 114} Inenat wend pontion: 50 Durction: right



Seck count: (60-50) + (79-60) + (92-79) + (114-92) + (176-114) + (176-41)
+(11-34) + (34-11) = 291.

\* clook is a modified algorithm of LOOK algorithm.