

1341a

## Algoritmi de criptare bazati pe $\mathbb{Z}_n$

• Context:  $\mathbb{Z}_{29}$

A	B	C	D	...	$\mathbb{Z}_{26}$	!	$\mathbb{Z}_{29}$
0	1	2	3		Z	...	!
					25	26	27 28

$(\mathbb{Z}_{29}, +, \cdot)$  corp comutativ

• Cifru (coduri) flux (stream cipher)  
= o cheie / mesaj

Cifru bloc / pe blocuri (block cipher)  
= mesajul se imparte in blocuri,  
o cheie / bloc

→ cu padding = toate blocurile au  
aceeasi lungime

→ fara padding :  $\leq 1$  bloc mai scurt

### Caesar

$$C = m + K$$

Ec. de criptare: Cod = Mesaj + Cheie

Ec. de decriptare: Mesaj = Cod - cheie  
 $m = c - K$

Flux: Mesaj: MESAJ

cheia: 11

$$[M, E, S, A, J] \rightarrow [12, 4, 18, 0, 9] \xrightarrow{+K}$$

$$\rightarrow [23, 15, 29, 11, 20] \xrightarrow{\div 29} [23, 15, 0, 11, 20]$$

$$\rightarrow [X, P, A, L, U] \rightarrow XPALU$$

$$\text{decriptare } [X, P, A, L, U] \rightarrow [23, 15, 0, 11, 20] \xrightarrow[-K]{-11}$$

[MESAJ]

pe blocuri fără padding: Mesaj: MERE

bloc: 3  $\Rightarrow$  MER, E

$$K1 = 40; K2 = 71$$

$$[M, E, R] \rightarrow [12, 4, 17] \xrightarrow[+40]{+K1} [52, 44, 57] \xrightarrow{\div 29}$$

$$\rightarrow [23, 15, 28] \rightarrow [X, P, ?]$$

$$[E] \rightarrow [4] \xrightarrow[+71]{+K2} [75] \xrightarrow{\div 29} [17] = R$$

XP?R

$$\underline{MERE} \rightarrow \underline{XP?R}$$

## Cu padding

Mesaj: MERF

Bluc:  $b=3 \rightarrow MER$

$K_1=40, K_2=71$     ENS

$[M, E, R] \rightarrow XP?$

$[E, N, S] \rightarrow [4, 13, 18] \xrightarrow{+71} [75, 84, 89]$

$\xrightarrow{+40} [17, 26, 2] \rightarrow [R, L, C]$

MERENS  $\rightarrow XP? R, L, C$

## Afin: Ec. de gradul I

Criptare:  $Cod = Mesaj \cdot K_1 + K_2$

Decriptare:  $Mesaj = (Cod - K_2) K_1^{-1}$

## Hill: Ec. matriceala

Criptare:  $Cod = MC \cdot \begin{pmatrix} M \\ E \\ S \\ A \\ J \end{pmatrix}$

Decriptare:  $\begin{pmatrix} M \\ E \\ S \\ A \\ J \end{pmatrix} = M^{-1} \cdot Cod$

Ex: Mesaj: Noi  $\rightarrow \begin{pmatrix} N \\ 0 \\ i \end{pmatrix} = \begin{pmatrix} 13 \\ 14 \\ 8 \end{pmatrix}$

MC:  $\begin{pmatrix} 1 & -1 & 2 \\ 0 & 3 & -1 \\ 2 & -2 & 1 \end{pmatrix}$

Criptarea:  $\begin{pmatrix} 1 & -1 & 2 \\ 0 & 3 & -1 \\ 2 & -2 & 1 \end{pmatrix} \begin{pmatrix} 13 \\ 14 \\ 8 \end{pmatrix} = \begin{pmatrix} 15 \\ 34 \\ 6 \end{pmatrix} \text{ mod } 29$

$= \begin{pmatrix} 15 \\ 5 \\ 6 \end{pmatrix} \rightarrow PFG$

Decriptarea  $\begin{pmatrix} 1 & -1 & 2 \\ 0 & 3 & -1 \\ 2 & -2 & 1 \end{pmatrix}^{-1} \begin{pmatrix} 15 \\ 5 \\ 6 \end{pmatrix} = \begin{pmatrix} 13 \\ 14 \\ 8 \end{pmatrix} \begin{matrix} N \\ 0 \\ i \end{matrix}$

Hill afîn;

Ec. de criptare:  $\text{Cod} = MC_1 \cdot \begin{pmatrix} M \\ E \\ S \\ A \\ i \\ j \end{pmatrix} + MC_2$

Decriptarea:  $\begin{pmatrix} M \\ E \\ S \\ A \\ i \\ j \end{pmatrix} = MC_1^{-1} (\text{Cod} - MC_2)$