# Computer are dumb

(And why you care.)

### Outline

- Algorithms and scaling
- Heuristics in computation
- Some thoughts on hardware
- Are you right? (Or at least not wrong?)

# Consider: finding SNPs.

Given: reference genome, sequence reads, mapping.

The mapping contains a list of reads, mapped locations within reference, and the location of differences.

How can we find all single-nucleotide variation?

### Approach one: by genome

```
for location in genome:
    reference = genome[location]
    bases = get_overlapping(location)
    for base in bases:
        if base != reference:
        # count SNP
```

# Approach two: by read

```
for read in mapped_reads:
   for differences in read:
     # count SNP
```

### Approach one: by genome

```
for location in genome:
    reference = genome[location]
    bases = get_overlapping(location)
    for base in bases:
        if base != reference:
            # count SNP
```

How does this algorithm scale? Imagine:

increasing size of genome increasing number of reads

### Approach two: by read

```
for read in mapped_reads:
   for differences in read:
     # count SNP
```

How does this algorithm scale?

### Scaling and Big-O notation

 The first approach scales with both the size of the genome and the number of reads:

$$t \sim O(N * M)$$

 The second approach scales with just the number of reads:

$$t \sim O(M)$$

### Scaling and Big-O notation

 The first approach scales with both the size of the genome and the number of reads:

$$t \sim O(N * M)$$

- why would you want this??

 The second approach scales with just the number of reads:

$$t \sim O(M)$$

### What about a different problem?

- I am interested in locations X,Y, and Z.
- Give me all SNPs at or near those locations.

```
for location in list_of_locations:
   reference = genome[location]
   bases = get_overlapping(location)
   for base in bases:
      if base != reference:
        # count SNP
```

### Important note

• Algorithm scaling is independent of the actual time it takes to run.

 Scaling tells you how time-to-run scales as the problem size changes, nothing more.

# Easy-to-check vs easy-to-find

Given a number, factor it into only prime numbers.

This is hard.

Given a set of prime numbers, verify that they multiple to yield a particular number.

This is easy.

### Easy-to-check vs easy-to-find, #2

Suppose:

50 dorm rooms, two students per room 100 students can be admitted, of 400 total Dean has list of students that cannot be paired.

It is easy to check any particular list of student/room combinations for validity.

In general, it is extremely hard to quickly find a guaranteed solution.

### Heuristics

- "Heuristics" are short cuts that usually work (but occasionally go horribly wrong).
- Not all problems are amenable.
  - Prime numbers? No good, fast short cut.
  - Housing? Sure start with a random solution, eliminate one of each pair that conflicts, until you find a non-conflict..
- Heuristics rely on assumptions about the specific type of problem you're going to tackle, and don't always work.
  - If the Dean is evil, he can construct a list of incompatible roommates that breaks your process.
  - Or he can just gives you a really long list of incompatible roommates.

### Example: BLAST

BLASTN filters sequences for exact matches between "words" of length 11:

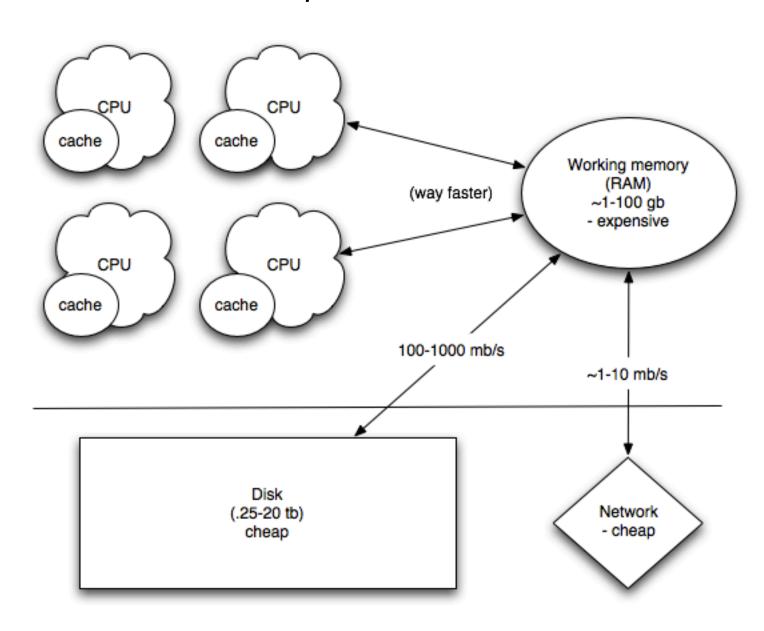
This results in a O(n log n) algorithm.

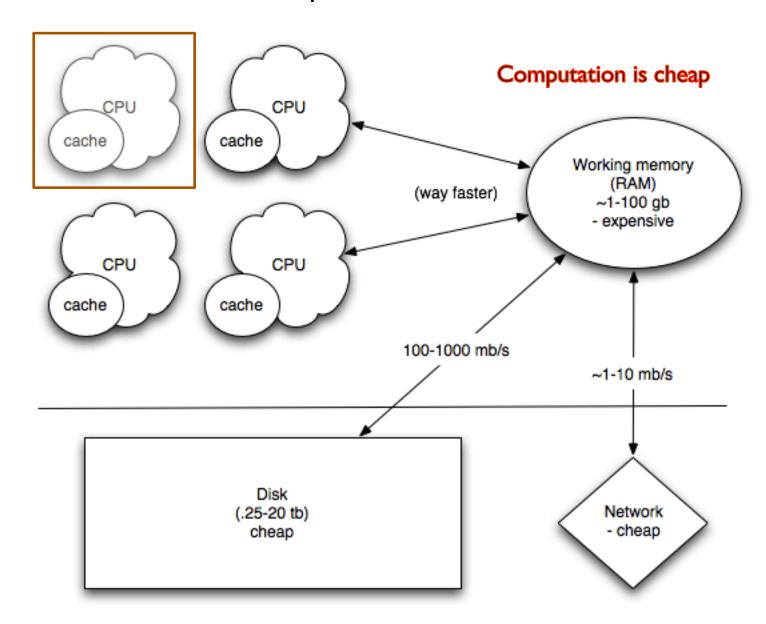
### Example: BLAST

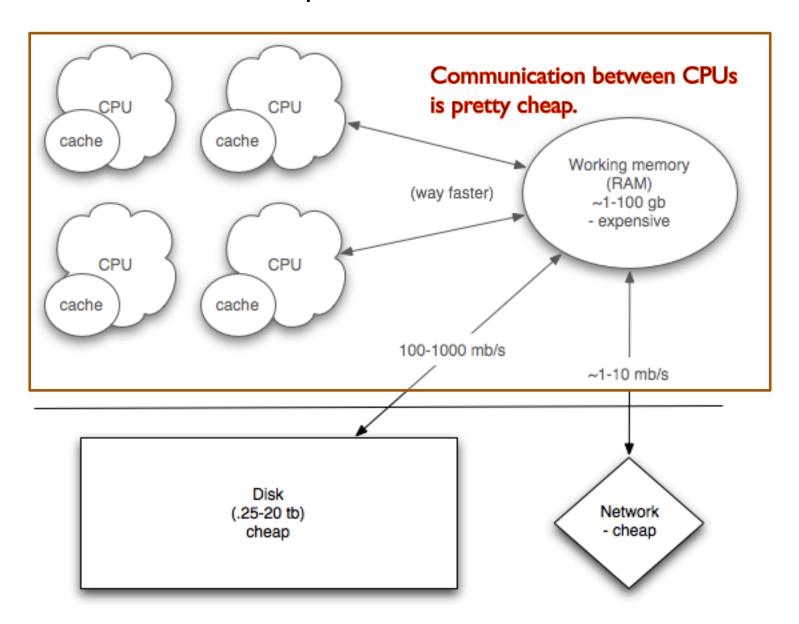
...but what about pathological situations?

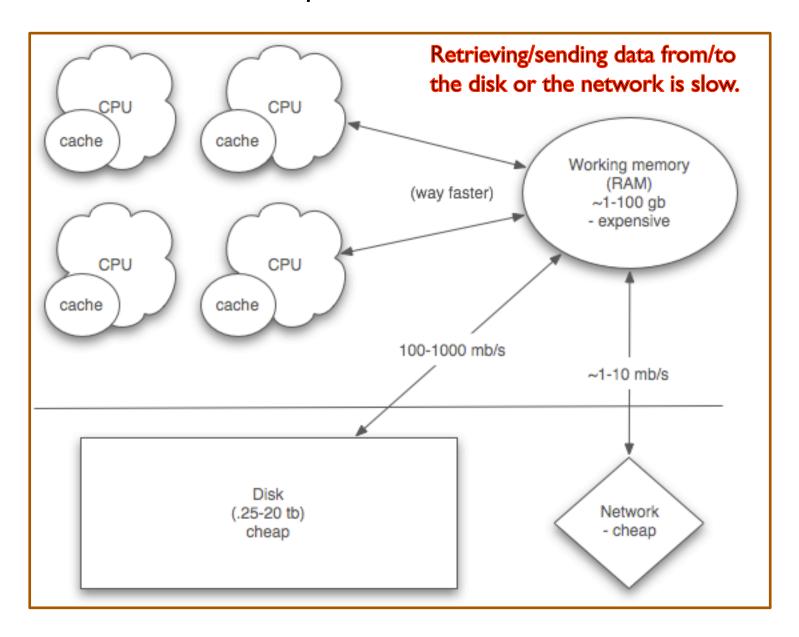
GAGGGTATGACGATATGGCGATGGAC | | x | | | | x | | | | | x | x | | | x | x | | x | x | | x | x | | x | x | | x | x | | x | x | | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x | x

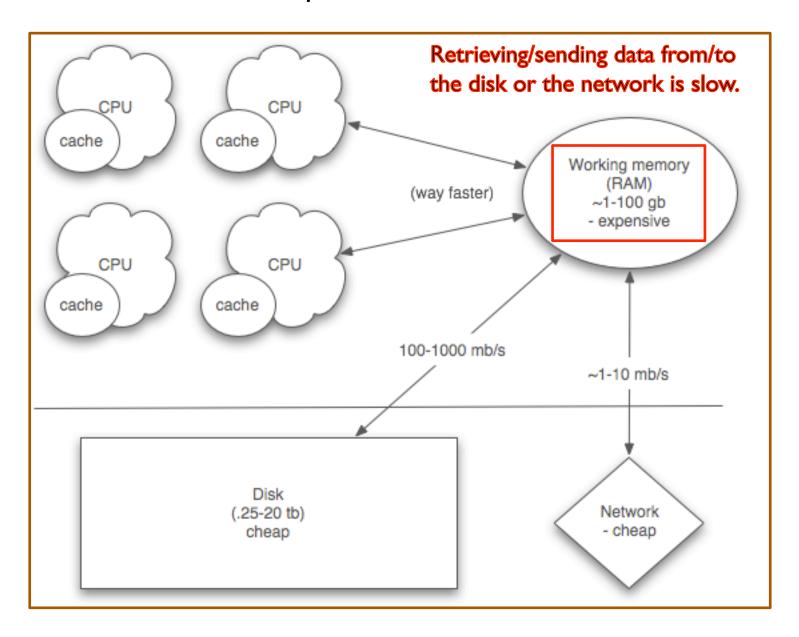
This will not be scored as a match, because BLAST *only* scores matches with a core "seed" match of 11 bases.











### Questions to ask

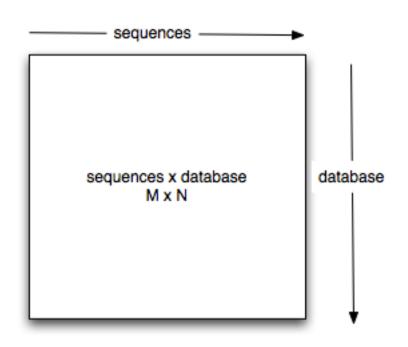
 Can I split my problem up into small chunks?

(because, if so, I can use more than one computer effectively.)

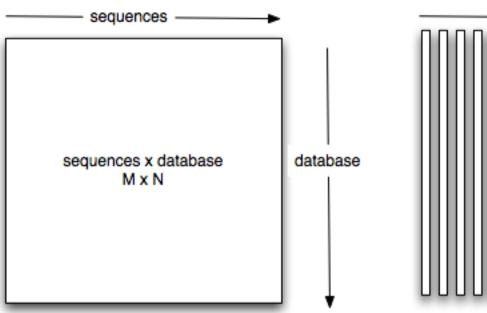
• How does my computation scale?

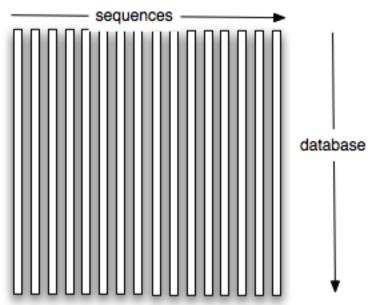
• How does my memory use scale?

# Sequence comparison: n^2



# ...but "embarrassingly parallel"





M problems, each 1 x N in size

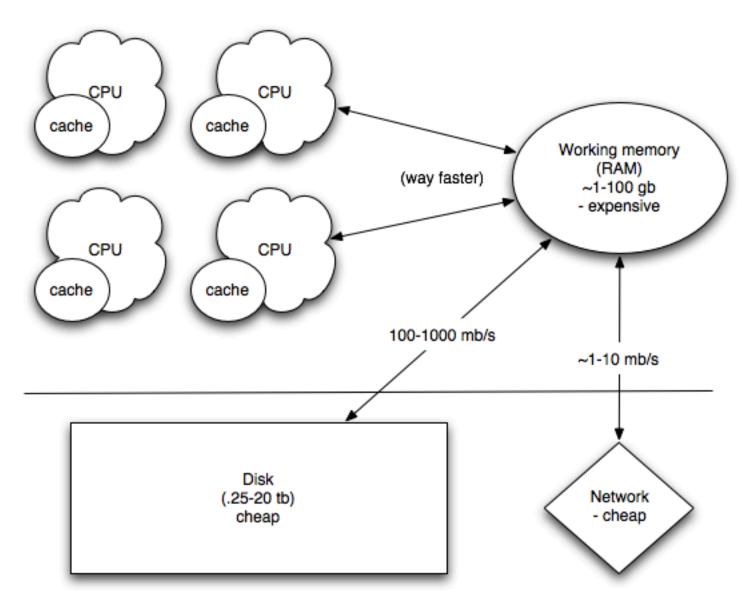
### Mapping is embarrassingly parallel.

You need to calculate individual overlaps.

# Assembly is not.

You need to calculate all overlaps.

Communication between CPUs is slow; this is main factor in splitting up tasks.



# None of this is the #1 problem you will face with bioinformatics.

Here is the #1 problem:

How would you know if your answer was right or wrong?

### Controls

- Just as with experiments, you can put negative and positive controls in your bioinformatics.
- e.g. with BLAST,
  - Do you see expected matches with the parameters and database you're using?
- Positive controls are often easier than negative, in "discovery-driven" science...

### Internal controls

 Molecules and sequences for which you have expectations.

 "I know this gene comes up, based on qPCR. I expect to see it in my mRNAseq."

### External controls

• Does the whole process work?

 "I can reproduce what this other person/ lab did, with their data, when I use my own software."

This is much more rarely done...

### Black box nature of algorithms

 When you listen to a computational biologist explain their clever algorithm...

• ...it's a mistake to think that they necessarily know what's going on.

 Software is full of bugs and unintended consequences.

### Tracking the process

 How do you know that your software today is doing the same thing it was doing last month? Or last year? Or in the hands of that other graduate student?

• There are some tools & techniques for dealing with change in software. We'll talk about them next week.

### Concluding thoughts

 Every step of the process needs to be critically thought about and controlled.

 Choice of algorithms can be important, but depends on your problem: the convenient tool in your toolbox may not be well suited to your problem.