

601.220 Intermediate Programming

Summer 2022, Meeting 9 (June 27th)

Today's agenda

- Exercise 15 review
- Midterm project overview
- "Day 17" material
 - Linked lists
 - Exercise 17
- "Day 18" material
 - More linked lists
 - Exercise 18

Reminders/Announcements

- Midterm project: due Friday, July 1st
- Midterm exam: in class on Wednesday, July 6th

review materials posted

```
(gdb) break endian.c:21
Breakpoint 1 at 0x1243: file endian.c, line 21.
(gdb) run
[...output omitted...]
Breakpoint 1, main () at endian.c:21
21    printf("%u\n", *p);
(gdb) print/x ((unsigned char *)p)[0]
$1 = 0x83
(gdb) print/x ((unsigned char *)p)[1]
$2 = 0x7e
(gdb) print/x ((unsigned char *)p)[2]
$3 = 0xa3
(gdb) print/x ((unsigned char *)p)[3]
$4 = 0x38
```

```
983
1005 105 15
```

0-9 A-F

In base-16, <u>950238851</u> is <u>38A37E83</u> Since we're seeing the bytes in order from least to most significant, the ugrad machines are <u>little</u> endian.

To negate a two's complement value:

- Invert all of the bits (the ~ operator is useful for this)
- Add 1

-- could also use (10 << 31)

Note that 0x80000000U is the unsigned int value with only the most significant bit set to 1. This is the sign bit, and values with this bit set are negative.

```
unsigned int magnitude(unsigned int value) {
  if ((value & 0x80000000U) == 0U) {
    return value; // value is non-negative
  }

  // value is negative, so invert bits and add 1
  value = ~value; // invert bits
  value += 1U; // add 1
  return value;
}
```

Generating a uniformly distributed pseudo-random integer in the range 0 (inclusive) to max_num (exclusive):

```
int gen_uniform(int max_num) {
  return rand() % max_num;
}
```

Generating 500 random values in range 0 (inclusive) to max_range (exclusive) and tallying them in the hist array:

```
for (int i = 0; i < 500; i++) {
  hist[gen_uniform(max_range)]++;
}</pre>
```

Generating normally-distributed integer values in the range 0 (inclusive) to max_range (exclusive):

```
int normal_rand(int max_num) {
  int result = 0;
  for (int i = 1; i < max_num; i++) {
    if ((rand() & 1) == 1) {
      result++;
    }
  }
  return result;
}</pre>
```

This is basically flipping a coin max_num-1 times and counting how many times it's heads.

Generating 500 normally-distributed values in the range 0 (inclusive) to max_range (exclusive) and tallying them in the hist array:

```
for (int i = 0; i < 500; i++) {
  hist[normal_rand(max_range)]++;
}</pre>
```

Day 17 recap questions

- 1 Describe the linked list structure by a diagram.
- 2 Compare arrays and linked lists. Write down their pros and cons.
- **3** What is a linked list's head? How is it different from a node? Explain.
- 4 How do you calculate length of a linked list?
- 6 How do you implement add_after on a singly linked list?

1. Describe the linked list structure by a diagram.

```
struct Node type:
struct Node {
  char payload; // payload could be any data type
  struct Node *next;
};
```

Example linked list

```
// code creating a linked list
- struct Node *head = malloc(sizeof(struct Node));
- head->payload = 'A';
- head->next = malloc(sizeof(struct Node));
- head->next->payload = 'B';
- head->next->next = malloc(sizeof(struct Node));
- head->next->next->payload = 'C';
- head->next->next->next = NULL;

Pay A Pay
```

A more concise representation



2. Compare arrays and linked lists. Write down their pros and cons. 1 = Helements

"random access" "constant time" Arrays:

- Pro(O(1) access to arbitrary element
- Con: O(N) to insert or remove element at arbitrary position
- Pro: better locality (fewer cache misses when iterating)
- Pro: more compact
- Con: fixed size, to reallocate must allocate new array and copy

existing data

Linked list pros and cons

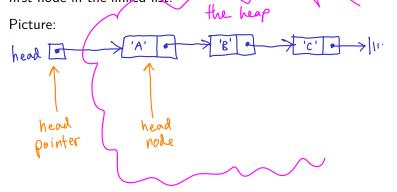


Linked list:

- Con: O(N) access to arbitrary element
- Pro: O(1) to remove element at arbitrary position
- Con: worse locality (more cache misses when iterating)
- Con: less compact (next pointers require space)
- Pro: can grow incrementally, nodes are allocated one at a time

3. What is a linked list's head? How is it different from a node? Explain.

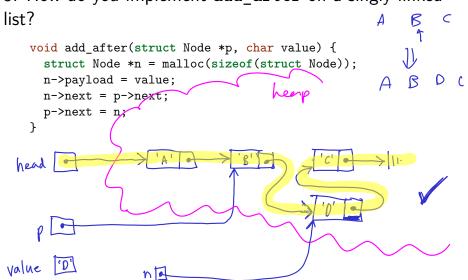
Contrast: *head pointer* vs. *head node*. The head pointer is a pointer variable storing a pointer to the first node. The head node *is* the first node in the linked list.



4. How do you calculate length of a linked list?

cur = head; A loop is required: struct Node *head = /* points to first node */; int count = 0; for (struct Node *cur = head; cur != NULL; cur = cur->next) { count++; count

5. How do you implement add_after on a singly linked



Exercise 17



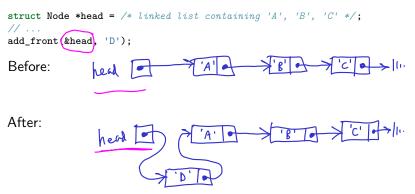
- Basic linked list functions
- Drawing pictures to reason about how linked lists operations should work is very helpful!
- Note that reverse_print is most easily implemented using recursion)
- Breakout rooms 1–10 are "social"
- Use Slack to let us know if you have questions

Day 18 recap questions

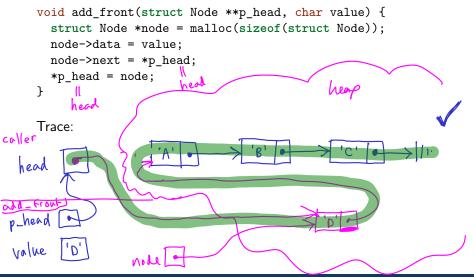
- How do you implement add_front on a linked list?
- How do you modify a singly linked list to create a doubly linked list?
- 3 How do you make a copy of a singly linked list?
- Why does add_after takes a struct Node * as input, but add front takes struct Node **?
- What cases should be handled when implementing remove_front?

4. Why does add_after takes a struct Node * as input, but add_front takes struct Node **?

Because add_after needs to change which node the head pointer points to. For example:



1. How do you implement add_front on a linked list?



2. How do you modify a singly linked list to create a doubly linked list?

Have each node store a pointer to the *previous* node in the list, in addition to the next node in the list. I.e.:

3. How do you make a copy of a singly linked list?

One way is to use recursion:

```
head Alli
```

```
struct Node *copy list(struct Node *n) {
  struct Node *result;
  if (n == NULL) {
    result = NULL:
  } else {
    result = malloc(sizeof(struct Node)):
    result->payload = n->payload;
    result->next = copy_list(n->next);
  }
  return result;
```

5. What cases should be handled when implementing remove_front?

There should not be any special cases.

```
void remove front(struct Node **p list) {
  assert(*p_list != NULL); head
  struct Node *succ = (*p_list)->next;
  free(*p_list); // free original head node
  *p list = succ; // make head pointer point to second node
```

Exercise 18

- More linked list operations (including ones requiring pointer to head pointer)
- Again, drawing diagrams is very helpful for reasoning about linked list operations
- Breakout rooms 1–10 are "social"
- Use Slack to let us know if you have any questions!