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(Engineering and MCA)
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Lab Manual - SQL Joins, Sub-Queries, and Views

Experiment No. 3

Title

SQL Queries – Joins, Sub-Queries, and Views

Objectives

- To understand and implement different types of **Joins** (INNER, LEFT, RIGHT, FULL, SELF).
- To design **Sub-queries** (single row, multiple row, correlated).
- To create and use **Views** for simplified data access.
- To perform at least **10 SQL queries** on a given case study using DML statements.

Problem Statement

"Consider an Online Bookstore application. Write SQL queries that demonstrate various Joins, Sub-queries, and Views to retrieve meaningful information about customers, books, orders, and payments."

Software and Hardware Requirements

- Software: MySQL / Oracle / PostgreSQL / SQL Server, SQL Workbench / pgAdmin / Oracle **SQL** Developer
- Hardware:

Processor: Intel i3 or higher

RAM: 4 GB or higher

Disk Space: 500 MB for DBMS installation

OS: Windows/Linux/Mac

Theory - Concept in Brief

Joins

- **INNER JOIN:** Returns matching rows from both tables.
- LEFT JOIN: Returns all rows from left table and matching from right.
- **RIGHT JOIN:** Returns all rows from right table and matching from left.
- **FULL OUTER JOIN:** Returns rows from both tables (where matches exist or not).
- **SELF JOIN:** A table joins with itself.

Sub-Queries

- Single-row subquery: Returns a single value.
- Multiple-row subquery: Returns multiple values (used with IN, ANY, ALL).
- Correlated subquery: Subquery depends on outer query.

Views

• Virtual tables created by storing SQL queries for reuse and simplified data access.

Algorithm

- 1. Start DBMS and connect to the database.
- 2. Use the normalized schema from the case study (Customer, Book, Orders, Order_Details, Payment).
- 3. Write SQL queries for:
 - All types of joins.
 - Sub-queries (single-row, multiple-row, correlated).
 - Views.
- 4. Execute queries and observe results.
- 5. Store queries and outputs for reporting.
- 6. End.

Flowchart : Create one yourself

Test Cases

Query No.	Query Description	Expected Output	Status
1	INNER JOIN: Customer & Orders	List of customers with their orders	Pass/Fail
2	LEFT JOIN: Books not ordered yet	List of books without sales	Pass/Fail
13	RIGHT JOIN: Orders with/without customer details	Order details	Pass/Fail
4	FULL OUTER JOIN: Customers and orders	All customers + all orders	Pass/Fail
5	SELF JOIN: Books by same author	Book pairs	Pass/Fail
6	Sub-query: Customer who placed highest order	Name of customer	Pass/Fail
7	Sub-query: Books priced above avg price	List of books	Pass/Fail
llX I	Correlated Sub-query: Orders above customer's avg spend	List of orders	Pass/Fail
9	Create View: Order summary	Virtual table with customer & total bill	Pass/Fail
10	Query View: Select top 5 customers by spending	List of 5 customers	Pass/Fail

Test Data Set

Customer

Cust_ID	Name	Email	Phone
1	Ramesh	ramesh@gmail.com	9876543210
2	Sneha	sneha@gmail.com	9998887777

Book

ISBN Title Author Price

B101 DBMS Concepts Korth 550

B102 SQL Fundamentals Ram 350

B103 Data Structures Weiss 450

Orders

Order_ID Order_Date Cust_ID

1001 2025-09-01 1

1002 2025-09-02 2

Order_Details

Order_ID ISBN Quantity

1001 B101 2

1001 B103 1

1002 B102 3

Payment

Pay_ID Pay_Type Amount Order_ID

5001 UPI 1650 1001

5002 Card 1050 1002

Sample SQL Queries

1. INNER JOIN: Customers with their orders

SELECT c.Name, o.Order_ID, o.Order_Date

FROM Customer c

INNER JOIN Orders o ON c.Cust_ID = o.Cust_ID;

2. LEFT JOIN: Books not yet ordered

SELECT b.Title, od.Order_ID

FROM Book b

LEFT JOIN Order_Details od ON b.ISBN = od.ISBN

WHERE od.Order_ID IS NULL;

RIGHT JOIN: Orders and Customers
 SELECT o.Order_ID, c.Name
 FROM Customer c
 RIGHT JOIN Orders o ON c.Cust_ID = o.Cust_ID;

4. FULL OUTER JOIN: Customers and Orders SELECT c.Name, o.Order_ID FROM Customer c FULL OUTER JOIN Orders o ON c.Cust_ID = o.Cust_ID;

5. SELF JOIN: Books by same author SELECT b1.Title AS Book1, b2.Title AS Book2 FROM Book b1, Book b2 WHERE b1.Author = b2.Author AND b1.ISBN <> b2.ISBN;

6. Sub-query: Customer with highest total order value SELECT Name FROM Customer WHERE Cust_ID = (SELECT Cust_ID FROM Orders o JOIN Payment p ON o.Order_ID = p.Order_ID GROUP BY Cust_ID ORDER BY SUM(p.Amount) DESC LIMIT 1);

7. Sub-query: Books priced above average SELECT Title, Price FROM Book WHERE Price > (SELECT AVG(Price) FROM Book);

8. Correlated Sub-query: Orders above customer's avg spend SELECT o.Order_ID, p.Amount FROM Orders o JOIN Payment p ON o.Order_ID = p.Order_ID WHERE p.Amount > (SELECT AVG(p2.Amount) FROM Orders o2 JOIN Payment p2 ON o2.Order_ID = p2.Order_ID WHERE o2.Cust_ID = o.Cust_ID);

9. Create View: Order Summary
CREATE VIEW OrderSummary AS
SELECT c.Name, o.Order_ID, SUM(b.Price * od.Quantity) AS Total
FROM Customer c
JOIN Orders o ON c.Cust_ID = o.Cust_ID
JOIN Order_Details od ON o.Order_ID = od.Order_ID
JOIN Book b ON od.ISBN = b.ISBN
GROUP BY c.Name, o.Order_ID;

10. Query from View: Top 5 Customers by Spending SELECT Name, SUM(Total) AS GrandTotal FROM OrderSummary GROUP BY Name ORDER BY GrandTotal DESC LIMIT 5;

Mathematical Model (if applicable)

- Let C = {c1, c2, ... cn} be set of customers.
- Let **B** = {b1, b2, ... bn} be set of books.
- Let **O** = {**o1**, **o2**, ... **om**} be set of orders.
- Relation R1 (Customer–Order): $C \times O \rightarrow$ mapping between customers and orders.
- Relation R2 (Order-Book): $O \times B \rightarrow$ mapping between orders and books.
- Function f: TotalAmount(o) = Σ (Price(b) × Quantity(b)) for all b in order o.

Conclusion / Analysis

In this experiment, we successfully implemented Joins (INNER, LEFT, RIGHT, FULL, SELF), Subqueries (single row, multiple row, correlated), and Views for the selected database application. These concepts enable complex data retrieval, improved query modularity, and better readability in SQL-based database systems.