# Discussion 03

**Floating Point; RISC-V Intro** 

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### Announcements <

## Agenda

- Floating Point
- RISC-V Intro

# Floating Point

# Floating Point Conversion

oating Points		
	8	23

1	8	23
Sign	Exponent	Significand/Mantissa

For normalized floats,

$$Value = (-1)^{Sign} * 2^{Exp+Bias} * 1. mantissa_2$$

For denormalized floats,

$$Value = (-1)^{Sign} * 2^{Exp+Bias+1} * 0. mantissa_2$$

$$Bias = -(2^{\# of \ exponent \ bits-1} - 1)$$

# **Decimal -> Floating Point**

$$Value = (-1)^{Sign} * 2^{Exp+Bias} * 1.mantissa_2$$

#### 10.75

- Convert to Binary w.r.t the floating point
  - 1010.11
- Shift floating point to match formula format
  - 1.01011 \* 2³
- Read Output
  - o Mantissa = 01100...0
  - $\circ$  Exp = 3 Bias

## **Step Size**

 Given a certain exponent, step size is the change in decimal value when we add 1 to the mantissa of binary FP

$$2^{exp+bias}*1.mantissa=>1M...M.MMMM$$

- Step Size is difference between 1MM...M.MMM0 and 1MM...M.MMM1
- 2<sup>-4</sup> in this example

## **Assembly Basics**

#### Assembly is...

- The direct output of compiled code
- Is not the final form of code
- Is still human-readable
- A set of instructions that can be directly understood by the system, after maybe some minor adjustments
- Built up of a single operation at a time
- Even more dumb than regular computer programs
- Read line by line when executed, except when told not to

# Storage: Registers

- On-chip memory
- RV32 has 32 of them numbered x0-x31
- They are all functionally the same but conventionally different
- They're all 32 bits wide
- Anything can be stored in them (no types)
- NOT A VARIABLE

	Name	Description	#	Name	Desc
:0	zero	Constant 0	x16		Args
x1	ra	Return Address	x17	a7	
<b>x</b> 2	sp	Stack Pointer	x18	s2	
<b>x</b> 3	gp	Global Pointer	x19	s3	
×4	tp	Thread Pointer	<b>x</b> 20	s4	ers
<b>x</b> 5	t0	Temporary Registers	x21	s5	gist
<b>x</b> 6	t1		x22	s6	Re
<b>x</b> 7	t2		<b>x23</b>	s7	Saved Registers
<b>x</b> 8	s0	Saved	x24	s8	S S
<b>x</b> 9	s1	Registers	<b>x2</b> 5	s9	
<b>x1</b> 0	a0	Function	<b>x</b> 26	<b>s10</b>	
x11	a1	Arguments or Return Values	<b>x</b> 27	s11	
<b>x12</b>	a2	Function Arguments	<b>x2</b> 8	t3	ies
<b>x13</b>	a3		<b>x29</b>	t4	rari
x14	a4		<b>x</b> 30	t5	Temporaries
<b>x1</b> 5	a5		<b>x31</b>	t6	7e
Caller saved registers					
Callee saved registers (except x0, gp, tp)					

# **Register Specifics**

There are 4 general categories (and some special!):

- 1. Argument registers: a0 a7
- 2. Return value registers: a0, a1
- 3. Saved registers: s0 s11
- 4. Temporary registers: to to Special registers
- Return address: ra NOT RETURN VALUE!!!

Zero: x0

Stack pointer: sp

Other pointers: tp, gp

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# RISC-V Greensheet!

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#### Loads

#### **I**<u>x</u> rd imm(rs1)

- Loads n bits worth of data from memory address: rs1 + imm where rs1 <u>should</u> already be a valid address.
- n will depend on the instruction: lb = 8 bits, lh = 16 bits, lw = 32 bits, ld = 16 bits, lw = 32 bits, ld = 16 bits
- If we have too few bits...
  - Sign-extend
- If we have too many bits...
  - Take the 32-most LSB bits!

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#### **Stores**

#### **s<u>x</u> rd imm(rs1)**

- Saves n bits worth of data to the memory address: rs2 + imm
- Truncates to n LSB bits to be stored

# Jump Instructions, jal, jalr, j, jr

	Saves return address	Jump and no return	
PC-relative address	jal ra, label	j label	
Address in a register	jalr ra, rs1, imm	jr rs1	

j and jr are shorthand for common combinations of certain instructions and register usages.

```
j label # jal x0, label
jr rs1 # jalr x0, 0(rs1)
```

Credits to Rosalie Fang

# Thank you