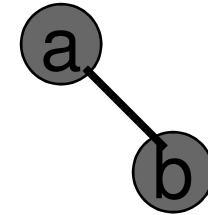
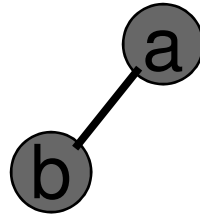


Binary Tree Construction

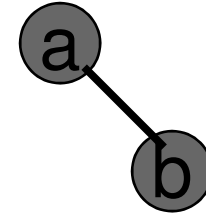
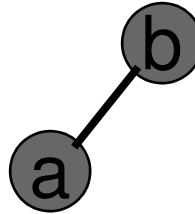
- Suppose that the elements in a binary tree are distinct.
- Can you construct the binary tree from which a given traversal sequence came ?
 - When a traversal sequence has more than one element, the binary tree is not uniquely defined.
 - Therefore, the tree from which the sequence was obtained cannot be reconstructed uniquely.

Examples: binary tree is not uniquely defined with a given traversal

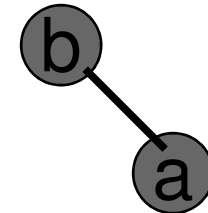
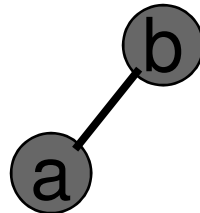
preorder = ab



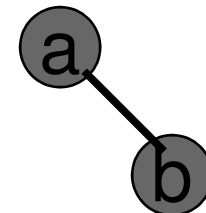
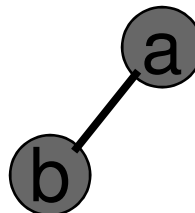
inorder = ab



postorder = ab



level order = ab



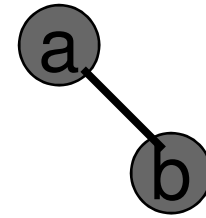
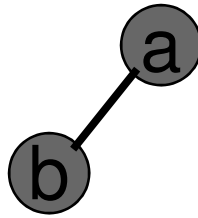
Binary Tree Construction

- Can you construct the binary tree, given two traversal sequences?
 - Depends on which two sequences are given.

Preorder and Postorder

preorder = ab

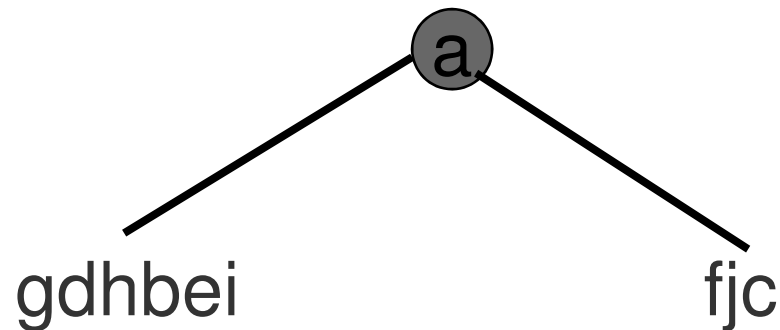
postorder = ba



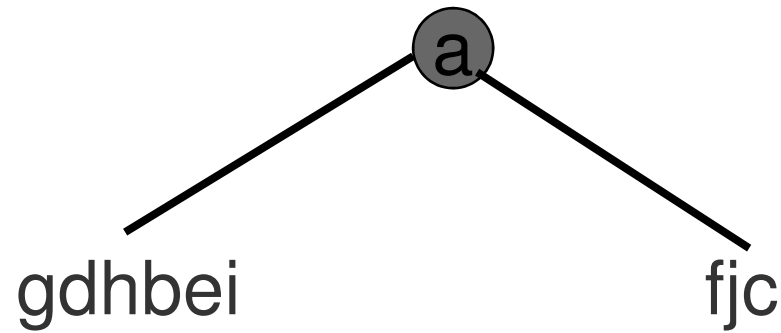
- Preorder and postorder do not uniquely define a binary tree.
- Nor do preorder and level order (same example).
- Nor do postorder and level order (same example).

Inorder and Preorder

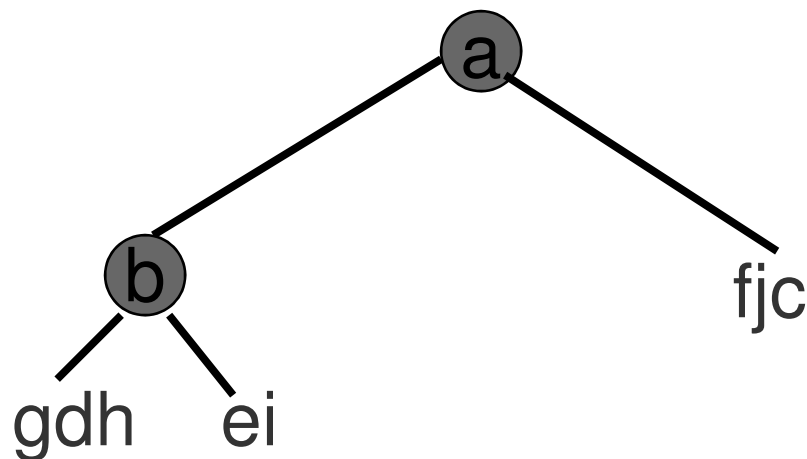
- inorder = g d h b e i a f j c
- preorder = a b d g h e i c f j
- Scan preorder from left to right using the inorder to separate left and right subtrees.
- a is the root of the tree; gdhbei are in the left subtree; fjc are in the right subtree.



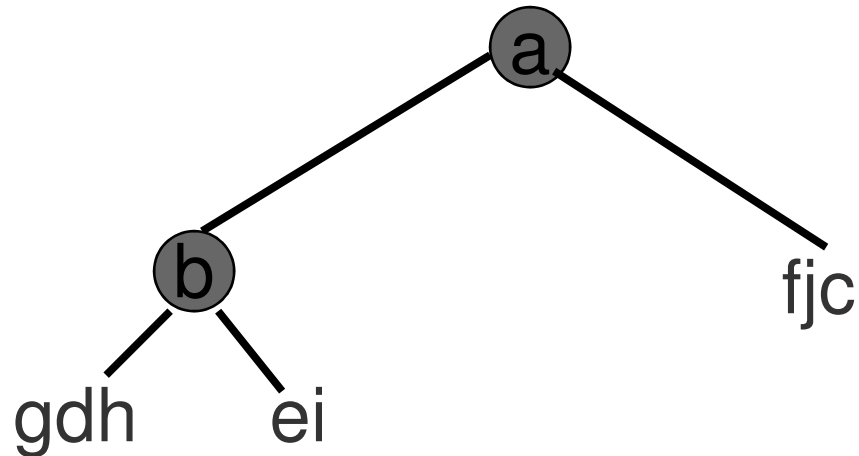
Inorder and Preorder



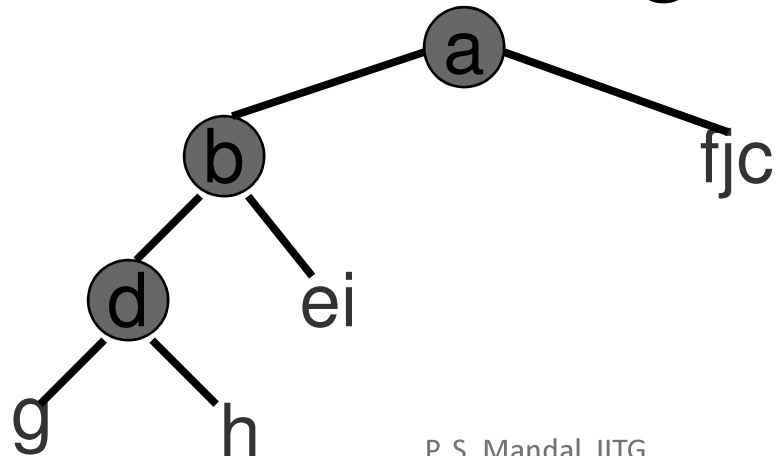
- preorder = a b d g h e i c f j
- b is the next root; gdh are in the left subtree; ei are in the right subtree.



Inorder and Preorder



- preorder = a b d g h e i c f j
- d is the next root; g is in the left subtree; h is in the right subtree.



Inorder and Postorder

- Scan postorder from right to left using inorder to separate left and right subtrees.
- inorder = g d h b e i a f j c
- postorder = g h d i e b j f c a
- Tree root is a; gdhbei are in left subtree; fjc are in right subtree.

Inorder and Level Order

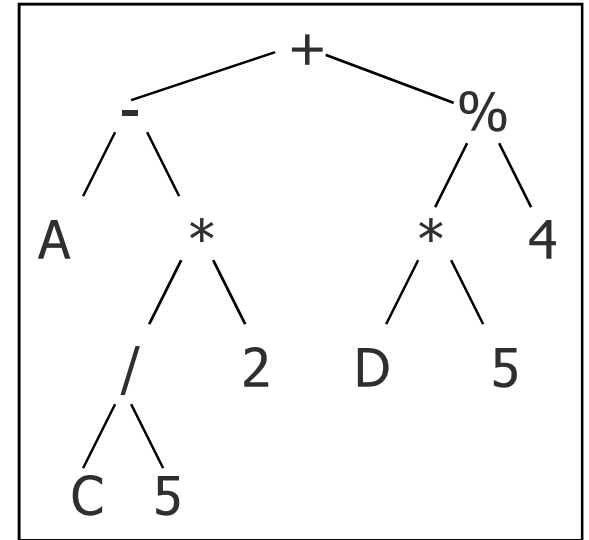
- Scan level order from left to right using inorder to separate left and right subtrees.
- inorder = g d h b e i a f j c
- level order = a b c d e f g h i j
- Tree root is a; gdhbei are in left subtree; fjc are in right subtree.

Uniquely define a binary tree

- for given two traversal sequences
 - Preorder and postorder X
 - Preorder and level order X
 - Postorder and level order X
 - Indoder and preorder ✓
 - Inorder and postorder ✓
 - Inorder and level order ✓

Expression Trees

- Expressions, programs, etc can be represented by tree structures
- It's a binary tree
- The leaves of the expression tree are operands
- The other nodes are contain operators
 - **E.g. Arithmetic Expression Tree**
 - **$A - (C / 5 * 2) + (D * 5 \% 4)$**

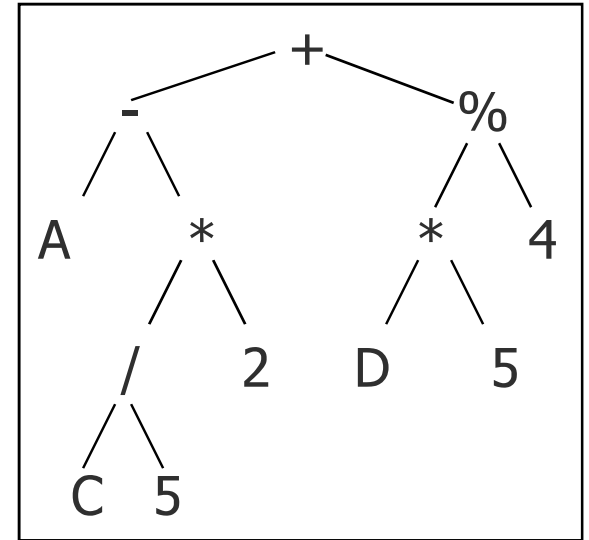


Tree Traversal

- Goal: visit every node of a tree
- **in-order** traversal

```
void Node::inOrder () {  
    if (left != NULL) left->inOrder();  
    cout << data << endl;  
    if (right != NULL) right->inOrder()  
}
```

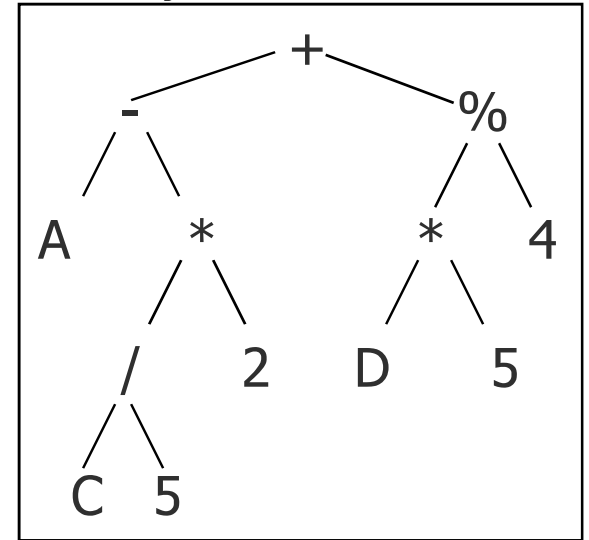
Output: $A - C / 5 * 2 + D * 5 \% 4$



Tree Traversal (contd.)

- **pre-order and post-order:**

```
void Node::preOrder () {  
    cout << data << endl;  
    if (left != NULL) left->preOrder ();  
    if (right != NULL) right->preOrder ();  
}
```



Output: + - A * / C 5 2 % * D 5 4

```
void Node::postOrder () {  
    if (left != NULL) left->preOrder ();  
    if (right != NULL) right->preOrder ();  
    cout << data << endl;  
}
```

Output: A C 5 / 2 * - D 5 * 4 % +

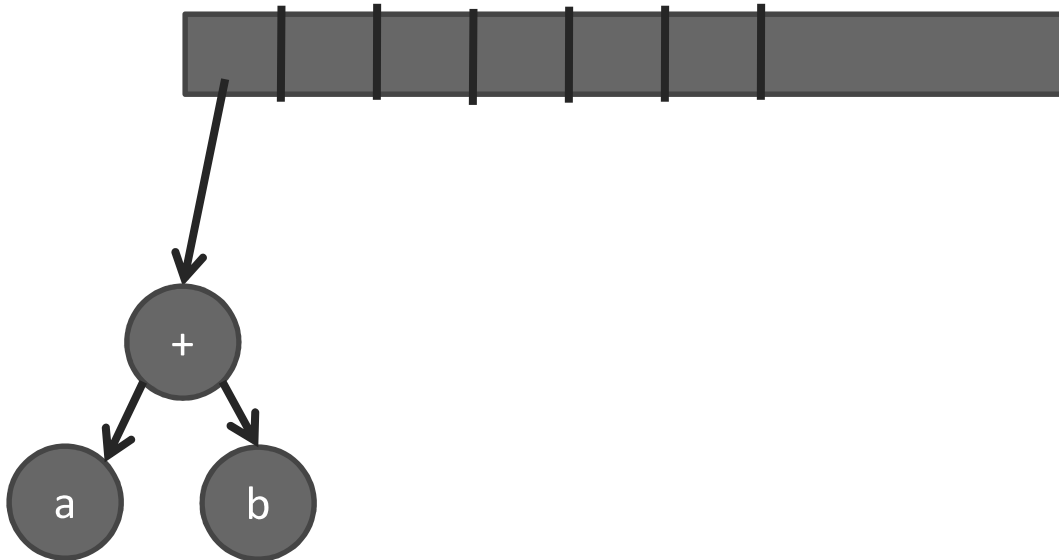
Constructing an expression Tree

- Convert postfix expression to expression tree
- $a\ b\ +\ c\ d\ e\ +\ * \ *$



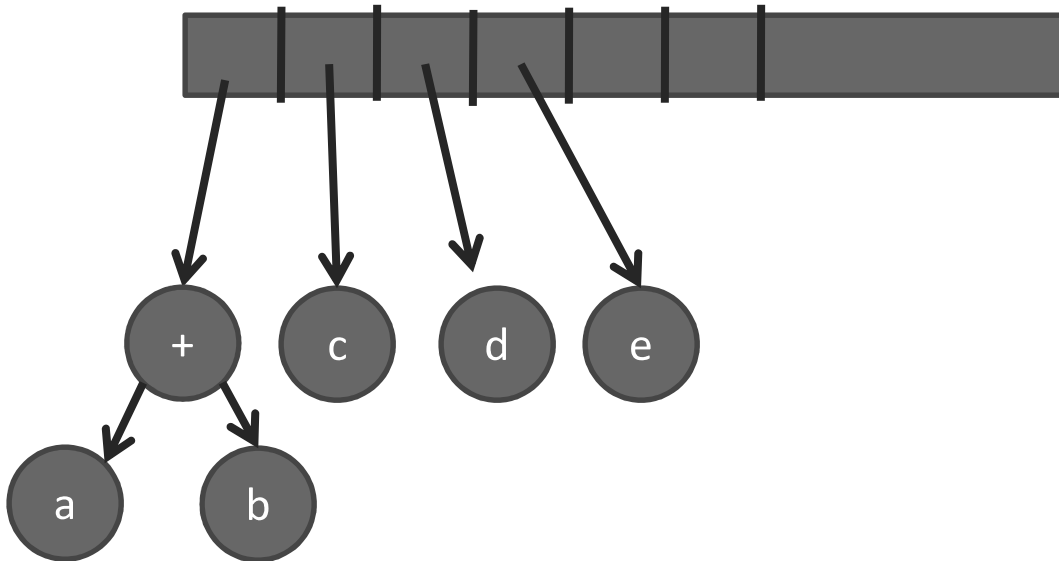
Constructing an expression Tree

- Convert postfix expression to expression tree
- $a\ b\ +\ c\ d\ e\ +\ * \ *$



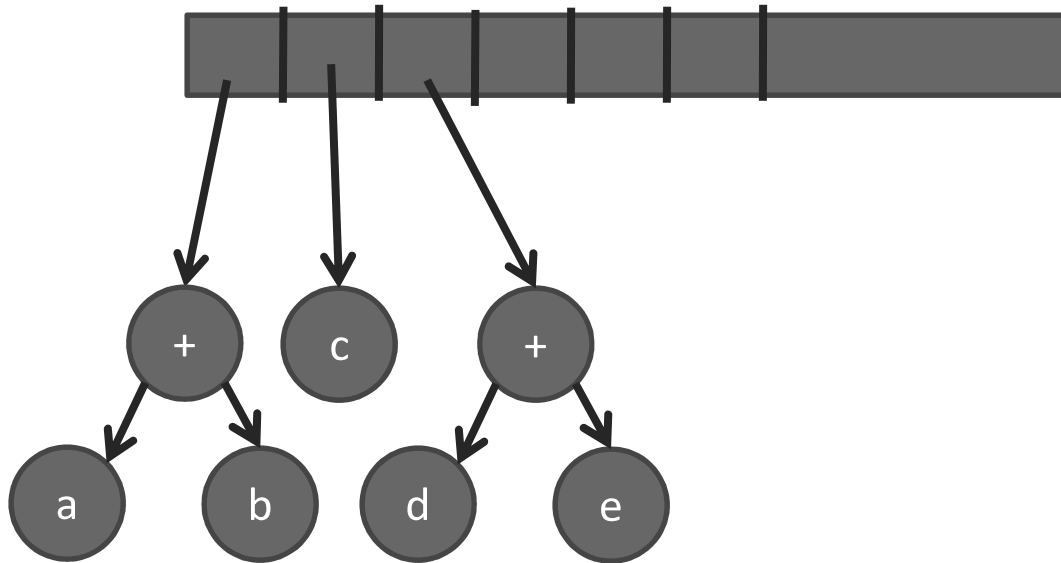
Constructing an expression Tree

- Convert postfix expression to expression tree
- $a\ b\ +\ c\ d\ e\ +\ * \ *$



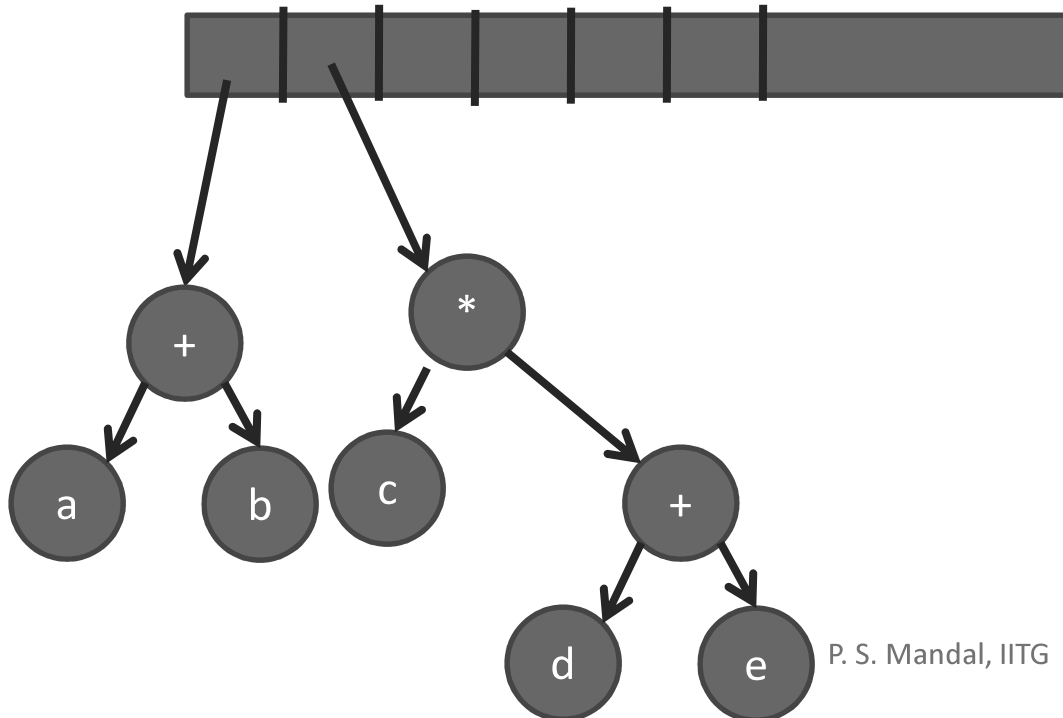
Constructing an expression Tree

- Convert postfix expression to expression tree
- $a\ b\ +\ c\ d\ e\ +\ * \ *$



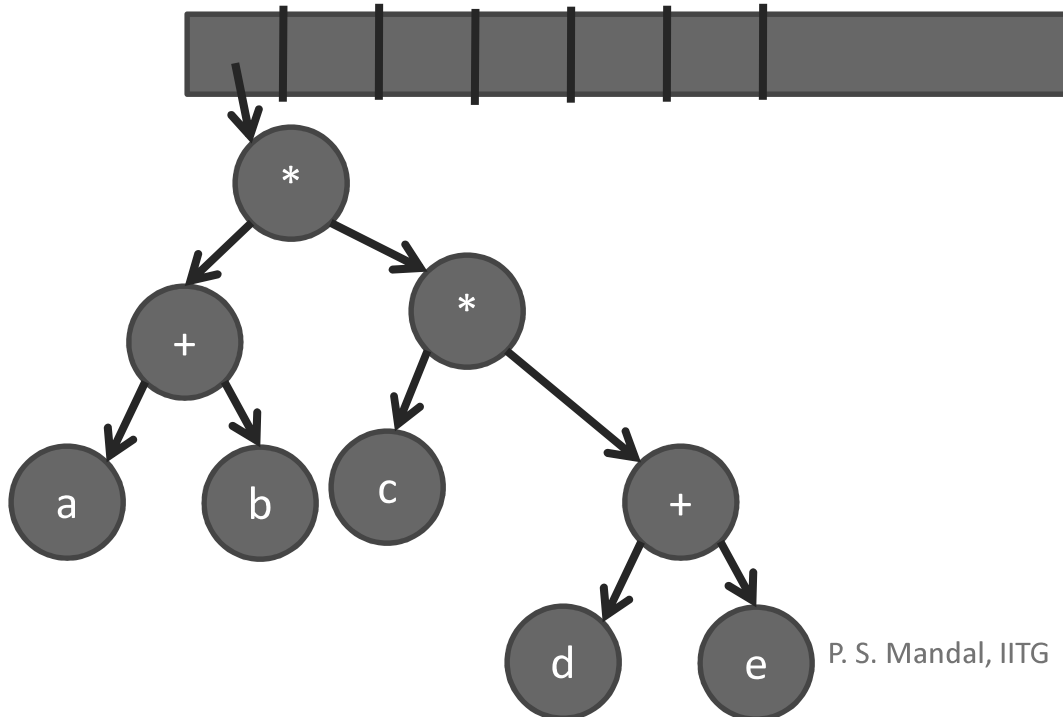
Constructing an expression Tree

- Convert postfix expression to expression tree
- $a\ b\ +\ c\ d\ e\ +\ *\ *$



Constructing an expression Tree

- Convert postfix expression to expression tree
- $a\ b\ +\ c\ d\ e\ +\ * \ *$

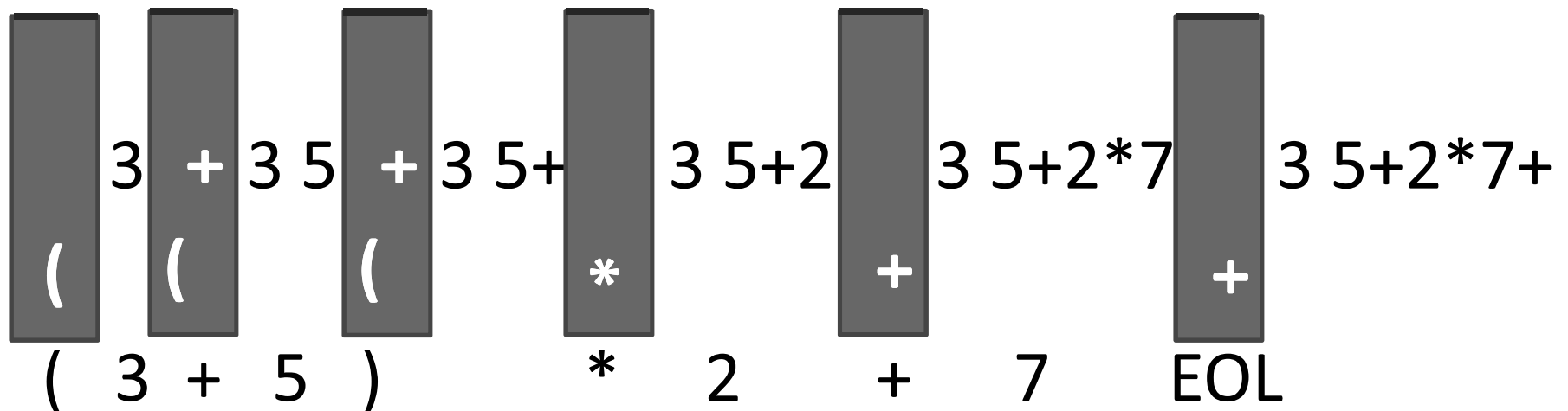


Infix to Postfix

- **Note:** input of an expression tree is a postfix expression.
- How we can convert infix expression to postfix expression ?
 - We can use stacks to do so

Infix to Postfix Example

- Infix expression: $(3 + 5) * 2 + 7$
- The operator stack holds just the operators. Operands are sent to the output directly.



- Postfix expression: $3\ 5\ +\ 2\ *\ 7\ +$