

IMPERIAL COLLEGE OF SCIENCE, TECHNOLOGY AND MEDICINE

EXAMINATIONS 2016

BEng Honours Degree in Computing Part I
MEng Honours Degrees in Computing Part I
BEng Honours Degree in Mathematics and Computer Science Part I
MEng Honours Degree in Mathematics and Computer Science Part I
for Internal Students of the Imperial College of Science, Technology and Medicine

*This paper is also taken for the relevant examinations for the
Associateship of the City and Guilds of London Institute*

PAPER C140=MC140

LOGIC

Tuesday 3 May 2016, 14:00
Duration: 80 minutes

Answer ALL TWO questions

Paper contains 2 questions
Calculators not required

- 1 a Draw the formation tree of the formula

$$\forall x P(x) \vee \neg Q(x) \wedge \exists y \neg R(x, y) \rightarrow x = y,$$

and indicate all free occurrences of variables in your tree.

- b Use propositional equivalences to show that the two formulas

$$p \vee q \quad \text{and} \quad (p \rightarrow q) \rightarrow q$$

are logically equivalent. In each step of your argument, state the equivalence you use (e.g., $\neg\neg A \equiv A$). Use only one equivalence in each step.

- c In natural deduction, the alternative \vee -elimination rule *Alt* $\vee E$ is as follows:

| | | |
|---|------------|------------------------------|
| 1 | $A \vee B$ | got this somehow |
| 2 | A | ass |
| 3 | \vdots | hard work |
| 4 | B | proved somehow |
| 5 | B | <i>Alt</i> $\vee E(1, 2, 4)$ |

Let \vdash denote the standard system of natural deduction. Let \vdash^a denote the system of natural deduction obtained from \vdash by adding the rule *Alt* $\vee E$ and removing the standard rule $\vee E$.

- Show that *Alt* $\vee E$ is a derived rule of \vdash .
- Show that $A \vee B, A \rightarrow C, B \rightarrow C \vdash^a C$.
(You may assume that PC is a derived rule of \vdash^a .)
- Explain why $\vee E$ is a derived rule of \vdash^a , using part c(ii) or otherwise.
- What does it mean to say that \vdash^a is *sound and complete*?
- Is \vdash^a sound and complete? Justify your answer.

The three parts carry, respectively, 10%, 25%, and 65% of the marks.

2a Prove by natural deduction that

$$\exists x P(x), \quad \exists x \forall y (y = x) \quad \vdash \quad \forall x P(x).$$

In the rest of this question, L is the 2-sorted signature with sorts Nat and $[\text{Nat}]$, constants $0, 1, 2, \dots : \text{Nat}$ and $[] : [\text{Nat}]$, function symbols $+, -, \times, :, ++, !!, \#$, and relation symbols $<, \leq$ and merge , of the appropriate sorts (as in lectures).

Variables i, j, k, m, n , etc., have sort Nat , and xs, ys, zs, us , etc., have sort $[\text{Nat}]$.

The intended semantics is an L -structure M whose domain consists of the natural numbers $0, 1, 2, \dots$ (sort Nat) and all lists of natural numbers (sort $[\text{Nat}]$). The symbols of L are interpreted in M as in lectures. For example, $M \models \text{merge}(ys, zs, xs)$ if and only if xs is a permutation of $ys++zs$ and the relative order of entries in ys and in zs is retained in xs .

You are given an L -formula $\text{in}(n, xs)$ expressing that n is an entry in xs , and an L -formula $\text{count}(n, xs, k)$ expressing that n occurs exactly k times in xs .

b Write down L -formulas expressing the following properties of xs and ys :

- i) Some entry in xs is an even number.
- ii) There are no repeated entries in xs .
- iii) ys is the result of sorting xs in ascending order.
- iv) xs has more even entries than odd entries.

c The binary relation $\text{entries}([\text{Nat}], [\text{Nat}])$ is specified informally by:
 $M \models \text{entries}(xs, ys)$ if and only if ys lists the entries in xs without repetitions.

Example: $M \models \text{entries}([1, 3, 1, 3, 3], ys)$ if and only if ys is $[1, 3]$ or $[3, 1]$.

Below are three *incorrect* attempts to express entries by an L -formula:

A1. $\forall n (\text{in}(n, xs) \leftrightarrow \text{in}(n, ys))$

A2. $\forall n (\text{in}(n, xs) \leftrightarrow \text{count}(n, ys, 1))$

A3. $\exists zs \text{merge}(ys, zs, xs) \wedge \forall n (\text{in}(n, xs) \rightarrow \text{count}(n, ys, 1))$

- i) For each of the formulas A1–A3, write down lists xs and ys for which the formula is true but $\text{entries}(xs, ys)$ is false, or vice versa.
- ii) Write down an L -formula $B(xs, ys)$ that does express entries (that is, $M \models \text{entries}(xs, ys) \leftrightarrow B(xs, ys)$ for all lists xs, ys). You do not need to justify your answer. You may use in and count freely.

The three parts carry, respectively, 35%, 40%, and 25% of the marks.