

CO202 – Software Engineering – Algorithms

Dynamic Programming - Solutions

Exercise 1: Compute LCS

	j	0	1	2	3	4	5	6
i		y_j	B	D	C	A	B	A
0	x_i							
1	A							
2	B							
3	C							
4	B							
5	D							
6	A							
7	B							

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	j	0	1	2	3	4	5	6
i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0						
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

```

8: for j = 1 to n
9:   for i = 1 to m
10:    if  $x_i == y_j$ 
11:       $c[i,j] = c[i-1,j-1] + 1$ 
12:       $b[i,j] = \nwarrow$ 
13:    elseif  $c[i-1,j] \geq c[i,j-1]$ 
14:       $c[i,j] = c[i-1,j]$ 
15:       $b[i,j] = \uparrow$ 
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i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	0					
2	B	0						
3	C	0						
4	B	0						
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i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0					
2	B	0						
3	C	0						
4	B	0						
5	D	0						
6	A	0						
7	B	0						

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0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0					
2	B	0	1					
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4	B	0						
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0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0					
2	B	0	↖ 1					
3	C	0						
4	B	0						
5	D	0						
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i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑					
2	B	0	↖					
3	C	0	1					
4	B	0						
5	D	0						
6	A	0						
7	B	0						

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9:   for i = 1 to m
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i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑					
2	B	0	↖					
3	C	0	↑					
4	B	0						
5	D	0						
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i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0	↑ 0	↑ 0	↖ 1	← 1	↖ 1
2	B	0	↖ 1	← 1	← 1	↑ 1	↖ 2	← 2
3	C	0	↑ 1	↑ 1	↖ 2	← 2	↑ 2	↑ 2
4	B	0	↖ 1	↑ 1	↑ 2	↑ 2	↖ 3	← 3
5	D	0	↑ 1	↖ 2	↑ 2	↑ 2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3	↖ 4
7	B	0	↖ 1	↑ 2	↑ 2	↑ 3	↖ 4	↑ 4

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	j	0	1	2	3	4	5	6
i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0	↑ 0	↑ 0	↖ 1	← 1	↖ 1
2	B	0	↖ 1	← 1	← 1	↑ 1	↖ 2	← 2
3	C	0	↑ 1	↑ 1	↖ 2	← 2	↑ 2	↑ 2
4	B	0	↖ 1	↑ 1	↑ 2	↑ 2	↖ 3	← 3
5	D	0	↑ 1	↖ 2	↑ 2	↑ 2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3	↖ 4
7	B	0	↖ 1	↑ 2	↑ 2	↑ 3	↖ 4	↑ 4

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	j	0	1	2	3	4	5	6
i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0	↑ 0	↑ 0	↖ 1	← 1	↖ 1
2	B	0	↖ 1	← 1	← 1	↑ 1	↖ 2	← 2
3	C	0	↑ 1	↑ 1	↖ 2	← 2	↑ 2	↑ 2
4	B	0	↖ 1	↑ 1	↑ 2	↑ 2	↖ 3	← 3
5	D	0	↑ 1	↖ 2	↑ 2	↑ 2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3	↖ 4
7	B	0	↖ 1	↑ 2	↑ 2	↑ 3	↖ 4	↑ 4

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	j	0	1	2	3	4	5	6
i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0	↑ 0	↑ 0	↖ 1	← 1	↖ 1
2	B	0	↖ 1	← 1	← 1	↑ 1	↖ 2	← 2
3	C	0	↑ 1	↑ 1	↖ 2	← 2	↑ 2	↑ 2
4	B	0	↖ 1	↑ 1	↑ 2	↑ 2	↖ 3	← 3
5	D	0	↑ 1	↖ 2	↑ 2	↑ 2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3	↖ 4
7	B	0	↖ 1	↑ 2	↑ 2	↑ 3	↖ 4	↑ 4

Exercise 1: Compute LCS

	j	0	1	2					
i		y_j	B	D					
0	x_i	0	0	0					
1	A	0	↑	↑					
2	B	0	↖	←					
3	C	0	↑	↑					
4	B	0	↖	↑	↑	↑	↖		
5	D	0	↑	↖	←	↑	↑	↑	↑
6	A	0	↑	↑	↑	↖	↑	↖	
7	B	0	↖	↑	↑	↑	↖	↑	←

```

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	j	0	1	2	3	4	5	6
i		y_j	B	D	C	A	B	A
0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0	↑ 0	↑ 0	↖ 1	← 1	↖ 1
2	B	0	↖ 1	← 1	← 1	↑ 1	↖ 2	← 2
3	C	0	↑ 1	↑ 1	↖ 2	← 2	↑ 2	↑ 2
4	B	0	↖ 1	↑ 1	↑ 2	↑ 2	↖ 3	← 3
5	D	0	↑ 1	↖ 2	← 2	↑ 2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3	↖ 4
7	B	0	↖ 1	↑ 2	↑ 2	↑ 3	↖ 4	← 4

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	j	0	1	2	3	4	5	6
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0	x_i	0	0	0	0	0	0	0
1	A	0	↑ 0	↑ 0	↑ 0	↖ 1	← 1	↖ 1
2	B	0	↖ 1	← 1	← 1	↑ 1	↖ 2	← 2
3	C	0	↑ 1	↑ 1	↖ 2	← 2	↑ 2	↑ 2
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5	D	0	↑ 1	↖ 2	← 2	↑ 2	↑ 3	↑ 3
6	A	0	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3	↖ 4
7	B	0	↖ 1	↑ 2	↑ 2	↑ 3	↖ 4	← 4

Exercise 2: Compute Levenshtein Distance

	j	0	1	2	3	4	5	6	7	8	9
i		y_j	E	M	P	I	R	I	C	A	L
0	x_i										
1	I										
2	M										
3	P										
4	E										
5	R										
6	I										
7	A										
8	L										

Exercise 2: Compute Levenshtein Distance

	j	0	1	2	3	4	5	6	7	8	9
i		y_j	E	M	P	I	R	I	C	A	L
0	x_i	0	1	2	3	4	5	6	7	8	9
1	I	1									
2	M	2									
3	P	3									
4	E	4									
5	R	5									
6	I	6									
7											
8											
9											

```
8: for j = 1 to n
9:   for i = 1 to m
10:    c =  $x_i == y_j$  ? 0 : 1
11:    d[i,j] = min(d[i-1,j] + 1, d[i,j-1] + 1, d[i-1,j-1] + c)
```

Exercise 2: Compute Levenshtein Distance

	j	0	1	2	3	4	5	6	7	8	9
i		y_j	E	M	P	I	R	I	C	A	L
0	x_i	0	1	2	3	4	5	6	7	8	9
1	I	1	R	1							
2	M	2									
3	P	3									
4	E	4									
5	R	5									
6	I	6									
7	C										
8	A										
9	L										

```

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Exercise 2: Compute Levenshtein Distance

	j	0	1	2	3	4	5	6	7	8	9
i		y_j	E	M	P	I	R	I	C	A	L
0	x_i	0	1	2	3	4	5	6	7	8	9
1	I	1	R1								
2	M	2	R2	D2							
3	P	3									
4	E	4									
5	R	5									
6	I	6									
7	C	7									
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9	L	9									

```

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i		y_j	E	M	P	I	R	I	C	A	L
0	x_i	0	1	2	3	4	5	6	7	8	9
1	I	1	R 1								
2	M	2	R D 2								
3	P	3	R D 3								
4	E	4									
5	R	5									
6	I	6									
7	C	7									
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Exercise 2: Compute Levenshtein Distance

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i		y_j	E	M	P	I	R	I	C	A	L
0	x_i	0	1	2	3	4	5	6	7	8	9
1	I	1	R 1								
2	M	2	R D 2								
3	P	3	R D 3								
4	E	4	K 3								
5	R	5									
6	I	6									
7	C	7									
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9	L	9									

```

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```

Exercise 2: Compute Levenshtein Distance

[illegible]

```

8:  for j = 1 to n
9:      for i = 1 to m
10:         c = xi == yj ? 0 : 1
11:         d[i,j] = min(d[i-1,j] + 1, d[i,j-1] + 1, d[i-1,j-1] + c)

```

Exercise 2: Compute Levenshtein Distance

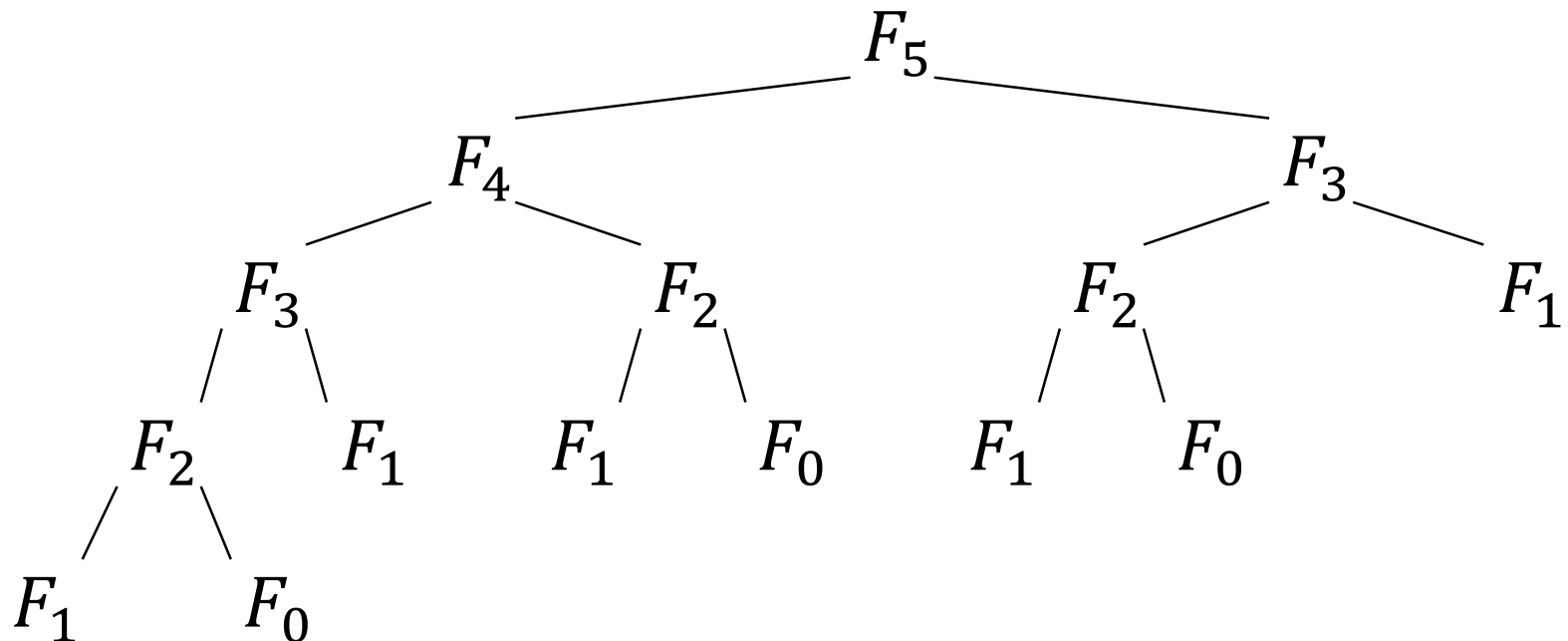
I	M	P	E	R	I	-	A	L
E	M	P	I	R	I	C	A	L
REPLACE	KEEP	KEEP	REPLACE	KEEP	KEEP	INSERT	KEEP	KEEP

	j	0	1	2	3	4	5	6	7	8	9		
i		y_j	E	M	P	I	R	I	C	A	L		
0	x_i	0	I 1	I 2	I 3	I 4	I 5	I 6	I 7	I 8	I 9		
1	I	1	D 1	R 1	I 2	I 3	K 3	I 4	I 5	I 6	I 7	I 8	
2	M	2	D 2	R 2	D 2	K 1	I 2	I 3	R 4	R 5	R 6	R 7	
3	P	3	D 3	R 3	D 3	D 2	K 1	I 2	I 3	I 4	I 5	I 6	I 7
4	E	4	D 4	K 3	D 3	D 2	R 2	R 3	R 4	R 5	R 6	R 7	
5	R	5	D 5	D 4	R 4	D 3	R 3	D 2	K 2	I 3	I 4	I 5	I 6
6	I	6	D 6	D 5	R 5	D 4	K 3	D 3	D 2	K 2	I 3	I 4	I 5
7	A	7	D 7	D 6	R 6	D 5	D 4	R 4	D 3	D 2	R 3	K 3	I 4
8	L	8	D 8	D 7	R 7	D 6	D 5	R 5	D 4	D 3	R 3	D 2	K 3

Exercise 3: Fibonacci Sequence

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, ...

$$F(n) = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F(n-1) + F(n-2), & \text{otherwise} \end{cases}$$



Exercise 3: Fibonacci Sequence

NAÏVE-FIBONACCI(*n*)

```
1: if n == 0
2:     return 0
3: if n == 1
4:     return 1
5: return NAÏVE-FIBONACCI(n-1) + NAÏVE-FIBONACCI(n-2)
```

Running time of NAÏVE-FIBONACCI:

$$T(n) = O(2^{0.694n})$$

Exercise 3: Fibonacci Sequence

BOTTOM-UP-FIBONACCI(*n*)

1: **if** *n* == 0

2: **return** 0

3: let *f*[0..*n*] be a new array

4: *f*[0] = 0

5: *f*[1] = 1

6: ?

7: ?

8: ?

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, ...

$$F(n) = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F(n-1) + F(n-2), & \text{otherwise} \end{cases}$$

What is the running time of BOTTOM-UP-FIBONACCI?

Exercise 3: Fibonacci Sequence

BOTTOM-UP-FIBONACCI(n)

1: **if** $n == 0$

2: **return** 0

3: let $f[0..n]$ be a new array

4: $f[0] = 0$

5: $f[1] = 1$

6: **for** $i = 2$ **to** n

7: $f[i] = f[i-1] + f[i-2]$

8: **return** $f[n]$

0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, ...

$$F(n) = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F(n-1) + F(n-2), & \text{otherwise} \end{cases}$$

What is the running time of BOTTOM-UP-FIBONACCI?

$T(n) = O(n)$...or?

Exercise 4: Coin Change Problem

Coin change is the problem of finding the least number of coins for a given amount of money.

For example, the UK coin set contains the following coins:

- 1p, 2p, 5p, 10p, 20p, 50p, £1, £2, and £5 (very uncommon).
- For £2.82, the optimal change is £2, 50p, 20p, 10p, 2p, so 5 coins.

1. Write a mathematical recurrence equation that determines the least number of coins.
2. Devise a pseudo-code, bottom-up dynamic programming algorithm `coin_change(n, coins)`.

Exercise 4: Coin Change Problem

1. Write a mathematical recurrence equation that determines the least number of coins.

$$counts[n] = \begin{cases} 0, & \text{if } n = 0 \\ \min_{coin} (counts[n - coin] + 1), & \text{otherwise} \end{cases}$$

Exercise 4: Coin Change Problem

2. Devise a pseudo-code, bottom-up dynamic programming algorithm `coin_change(n, coins)`.

```
01: def coin_change(n, coins):
02:     if n == 0:
03:         return 0
04:     counts = [0]*(n+1)
05:     change = [0]*(n+1)
06:     for i in range(1, n+1):
07:         counts[i] = n+1
08:         for coin in coins:
09:             if coin <= i:
10:                 count = 1 + counts[i-coin]
11:                 if count < counts[i]:
12:                     counts[i] = count
13:                     change[i] = coin
14:     return counts, change
```

Exercise 5: Fibonacci Challenge

D&C Fibonacci Revisited

Naïve:

$$F(n) = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F(n-1) + F(n-2), & \text{otherwise} \end{cases}$$

Let's rewrite Fibonacci

$$F(n) = F(2k) = F(k)^2 + 2F(k)F(k-1) \quad \text{for even } n$$

$$F(n) = F(2k-1) = F(k)^2 + F(k-1)^2 \quad \text{for odd } n$$

$$F(n) = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F(\lceil n/2 \rceil)^2 + 2F(\lceil n/2 \rceil)F(\lceil n/2 \rceil - 1), & \text{if } n \geq 2 \text{ and } n \text{ is even} \\ F(\lceil n/2 \rceil)^2 + F(\lceil n/2 \rceil - 1)^2, & \text{if } n \geq 2 \text{ and } n \text{ is odd} \end{cases}$$

Exercise 5: Fibonacci Challenge

D&C Fibonacci Revisited

$$F(n) = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F(\lceil n/2 \rceil)^2 + 2F(\lceil n/2 \rceil)F(\lceil n/2 \rceil - 1), & \text{if } n \geq 2 \text{ and } n \text{ is even} \\ F(\lceil n/2 \rceil)^2 + F(\lceil n/2 \rceil - 1)^2, & \text{if } n \geq 2 \text{ and } n \text{ is odd} \end{cases}$$

DC-FIBONACCI(*n*)

```
1: if n == 0 || n == 1
2:     return n
3: else
4:     a = DC-FIBONACCI((n+1)/2)
5:     b = DC-FIBONACCI((n+1)/2-1)
6:     if n is even
7:         return a * (a + 2 * b)
8:     else
9:         return a * a + b * b
```

$$T(n) = 2T(n/2) + \Theta(1)$$

Case 1: $d < \log_b a$, then $T(n) = O(n^{\log_b a})$

Exercise 5: Fibonacci Challenge

Dynamic Programming vs. Divide-and-Conquer

BOTTOM-UP-FIBONACCI(n)

```
1: if n == 0
2:     return 0
3: let f[0..n] be a new array
4: f[0] = 0
5: f[1] = 1
6: for i = 2 to n
7:     f[i] = f[i-1] + f[i-2]
8: return f[n]
```

DC-FIBONACCI(n)

```
1: if n == 0 || n == 1
2:     return n
3: else
4:     a = DC-FIBONACCI((n+1)/2)
5:     b = DC-FIBONACCI((n+1)/2-1)
6:     if n is even
7:         return a * (a + 2 * b)
8:     else
9:         return a * a + b * b
```

Asymptotic running time: $O(n)$
But, who wins in practice?