Rational Agents

Paolo Turrini

Department of Computing, Imperial College London

Introduction to Artificial Intelligence 2nd Part

What you have seen

You have seen examples of computational problem-solving:

- Search
- Planning
- Pattern recognition via neural networks

- Able to reason about the world around
 - True facts (knowledge)
 - Plausible facts (beliefs)

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 - Imperfect and incomplete information
 - Quantifying uncertainty, attaching probabilities
 - Going for uncertain outcomes, calculating expected utility

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 - Going for uncertain outcomes, calculating expected utility
- Able to update their beliefs when confronted with new information (learning)

What is rationality?



Robert J. Aumann Nobel Prize Winner

'A person's behaviour is rational if it is in their best interests, given their information.'

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Agents (not only humans) can be rational!

- The agent and the world
 - Actions and knowledge
 - Inference
- Good decisions
 - Chance
 - Gains
- Good decisions in time
 - Chance and gains in time
 - Patience
 - Finding the best strategy
- Learning from experience
 - Finding a reasonable strategy



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- Week 2: Logic
- Week 3: Decision-making
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- Week 2: Logic
- Week 3: Decision-making [last year's exam]
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- Last day: Reinforcement Learning Lab

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The book

- Stuart Russell and Peter Norvig Artificial Intelligence: a modern approach. 3rd Edition.
 - You can get it for free. I'm not suggesting to download it.
 - Lots of useful exercises

How to contact me

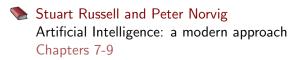
- Come and talk to me (After Thursday's class, Huxley 452)
- Send me an email (p.turrini@ic.ac.uk)
- Piazza

Knowledge Representation

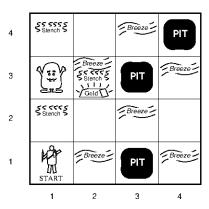
An uncertain world



The main reference



The Wumpus World



The Wumpus World

Sensors Breeze, Glitter, Smell Actuators Up, Down, Left, Right, Grab, Release, Shoot, Climb

Rewards 1000 escaping with gold, -1000 dying, -10 using arrow, -1 walking

Environment

- Squares adjacent to Wumpus are smelly
- Squares adjacent to pit are breezy
- Glitter iff gold is in the same square
- Shooting kills Wumpus if you are facing it
- Shooting uses up the only arrow
- Grabbing picks up gold if in same square
- Releasing drops the gold in same square

Knowledge base

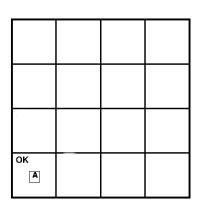
- A set of sentences representing what the agent thinks about the world.
 - 'I am in [2,1]'
 - 'I am out of arrows'
 - 'I smell Wumpus'
 - 'I'd better not go forward'
- We interpret it as what the agent knows,
 but it works just fine for what the agent believes.

Updating the knowledge base

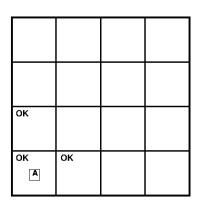
- What we TELL the knowledge base
- What we ASK the knowledge base

```
function KB-AGENT(percept) returns an action static: KB, a knowledge base t, a counter, initially 0, indicating time Tell(KB, Make-Percept-Sentence(percept, t)) action \leftarrow Ask(KB, Make-Action-Query(t)) Tell(KB, Make-Action-Sentence(action, t)) t \leftarrow t + 1 return action
```

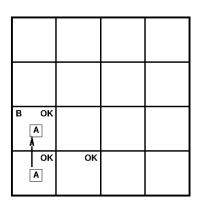
• The starting state...



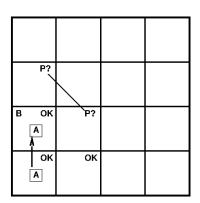
• and what we know.



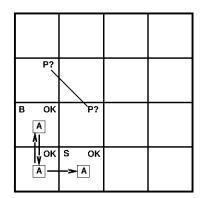
• B stands for Breeze



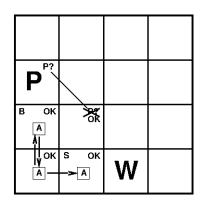
- Where is the pit?
- We are ruling out one square!



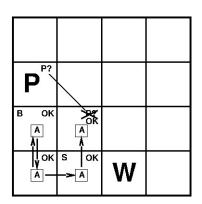
- S stands for smell
- What do we know?



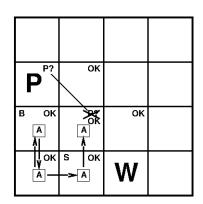
• Logic is the key!



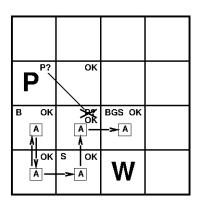
• The further we go the more we know



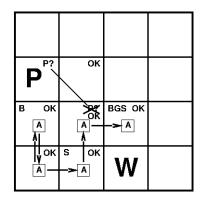
• The further we go the more we know



Gold!



- We know the way out
- Game over



Let $P_{i,j}$ be true if there is a pit in [i,j]. Let $B_{i,j}$ be true if there is a breeze in [i,j].

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\neg P_{1,1}
\neg B_{1,1}
B_{2,1}
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 $B_{2,1}$

"Pits cause breezes in adjacent squares"

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 $B_{2,1}$

"Pits cause breezes in adjacent squares"

$$B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$$

 $B_{2,1} \Leftrightarrow (P_{1,1} \vee P_{2,2} \vee P_{3,1})$

"A square is breezy if and only if there is an adjacent pit"

Expressivity: at what cost?

- OK if we were only dealing with finite objects
- But even then we would have to enumerate all the possibilities

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Propositional Logic lacks expressive power

First order logic

- Massive increase of expressivity
- But there are costs, e.g., decidability
- We will see how to exploit the gains while limiting the costs

KB with FOL

 We can encode the KB at each particular time point using FOL

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 Percept([Stench, Breeze, Glitter], 5) or
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 - Axioms from knowledge to knowledge, e.g., $\forall t \; AtGold(t) \land Action(Grab, t) \Rightarrow Holding(Gold, t + 1)$

Perception $\forall s, b, t \ Percept([s, b, Glitter], t) \Rightarrow AtGold(t)$

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Perception \forall s, b, t \; Percept([s, b, Glitter], t) \Rightarrow AtGold(t)
Location At(Agent, s, t)
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Decision-making \forall t \; AtGold(t) \Rightarrow Action(Grab, t)
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Location At(Agent, s, t)

Decision-making \forall t \; AtGold(t) \Rightarrow Action(Grab, t)

Internal reflection \forall t \; AtGold(t) \land \neg Holding(Gold, t) \Rightarrow Action(Grab, t), do we have gold already? (notice we cannot observe if we are holding gold, we need to track it)
```

Adjacent squares

$$\forall x, y, a, b \ Adjacent([x, y], [a, b]) \Leftrightarrow (x = a \land (y = b - 1 \lor y = b + 1) \lor (y = b \land (x = a - 1 \lor x = a + 1))$$

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"A square is breezy if and only if there is an adjacent pit"

$$\forall s , Breezy(s) \Leftrightarrow \exists r (Adjacent(r, s) \land Pit(r))$$

- We can go on and describe plans, causal rules, etc.
- But let's do some reasoning now

Facts and knowledge bases



'Joffrey Baratheon is a king'

Facts and knowledge bases



'Jon Snow is a person'

Facts and knowledge bases



'Jon Snow is a king'

Tell(*KB*, *King*(*Joffrey*))

```
Tell(KB, King(Joffrey))
Tell(KB, Person(Jon))
```

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Tell(KB, King(Joffrey))

Tell(KB, Person(Jon))

Tell(KB, \forall x \ King(x) \Rightarrow Person(x))
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Tell(KB, King(Joffrey))

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Tell(KB, \forall x \ King(x) \Rightarrow Person(x))

Ask(KB, \exists x Person(x)) is there a person?
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Askvar(KB, Person(x)) who is a person?
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Tell(KB, King(Joffrey))

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Askvar(KB, Person(x)) who is a person?

Askvar returns a list of substitutions: \{x/Joffrey\}, \{x/Jon\}
```

Definition

Given a sentence S and a substitution σ ,

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$$S\sigma = Smarter(Tyrion, Joffrey)$$

Definition

Given a sentence S and a substitution σ ,

 $S\sigma$ denotes the result of plugging σ into S; e.g.,

S = Smarter(x, y)

 $\sigma = \{x/\mathit{Tyrion}, y/\mathit{Joffrey}\}$

 $S\sigma = Smarter(Tyrion, Joffrey)$

Askvar(KB, S) returns some/all σ such that $KB \models S\sigma$

$$\forall x \; King(x) \land Greedy(x) \Rightarrow Evil(x)$$

$$King(Joffrey)$$

$$\forall y \; Greedy(y)$$

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We can get the inference immediately if we can find a substitution matching the premises of the implication to the known facts.

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$$\theta = \{x/Joffrey, y/Joffrey\}$$
 works



UNIFY(
$$\alpha, \beta$$
) returns θ if $\alpha \theta = \beta \theta$

| p | q | θ |
|-------------------|---------------------------|----------|
| Knows(Joffrey, x) | Knows(Joffrey, Sansa) | |
| Knows(Joffrey, x) | Knows(y, Sansa) | |
| Knows(Joffrey, x) | Knows(y, Mother(Joffrey)) | |
| Knows(Jon, x) | Knows(x, Mother(Jon)) | ' |

| p | q | $\mid 	heta$ |
|-------------------|---------------------------|---------------|
| Knows(Joffrey, x) | Knows(Joffrey, Sansa) | $\{x/Sansa\}$ |
| Knows(Joffrey, x) | Knows(y, Sansa) | |
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| Knows(Joffrey, x) | Knows(y, Sansa) | $\{x/Sansa, y/Joffrey\}$ |
| Knows(Joffrey, x) | Knows(y, Mother(Joffrey)) | $\{y/Joffrey, x/Mother(Joffrey)\}$ |
| Knows(Jon, x) | Knows(x, Mother(Jon)) | |

| p | q | $\mid 	heta \mid$ |
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| Knows(Joffrey, x) | Knows(y, Mother(Joffrey)) | $\{y/Joffrey, x/Mother(Joffrey)\}$ |
| Knows(Jon, x) | Knows(x, Mother(Jon)) | fail |

Standardising apart

Knows(Jon, x) & Knows(x, Mother(Jon)) fails

Standardising apart

$$Knows(Jon, x)$$
 & $Knows(x, Mother(Jon))$ fails

Standardising apart eliminates overlap of variables, e.g., $Knows(z_{17}, Mother(Jon))$

Definite clause:

disjunction of literals, exactly one of which positive

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$$\frac{p_1', p_2', \dots, p_n', (p_1 \wedge p_2 \wedge \dots \wedge p_n \Rightarrow q)}{q\theta} \quad \text{where } p_i'\theta = p_i\theta \text{ for all } i$$

Definite clause:

disjunction of literals, **exactly** one of which positive e.g., $(p_1 \land p_2 \land ... \land p_n \Rightarrow q)$

$$\frac{p_1', p_2', \dots, p_n', (p_1 \land p_2 \land \dots \land p_n \Rightarrow q)}{q\theta} \quad \text{where } p_i'\theta = p_i\theta \text{ for all } i$$

Assuming all variables are universally quantified...

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Assuming all variables are universally quantified...

$$p_1'$$
 is $King(Joffrey)$ p_1 is $King(x)$ p_2' is $Greedy(y)$ p_2 is $Greedy(x)$ θ is $\{x/Joffrey, y/Joffrey\}$ q is $Evil(x)$ $q\theta$ is $Evil(Joffrey)$

Soundness of GMP

Need to show that

$$p_1', \ldots, p_n', (p_1 \wedge \ldots \wedge p_n \Rightarrow q) \models q\theta$$

provided that $p_i'\theta = p_i\theta$ for all i

Lemma: If φ is definite clause, then $\varphi \models \varphi \theta$ by Universal Instantiation.

- $(p_1 \wedge \ldots \wedge p_n \Rightarrow q) \models (p_1 \wedge \ldots \wedge p_n \Rightarrow q)\theta = (p_1\theta \wedge \ldots \wedge p_n\theta \Rightarrow q\theta)$
- **3** From 1 and 2, $q\theta$ follows by ordinary Modus Ponens



- How to describe the world in logic
- Moving as a way to gather new facts
- Generalised modus ponens

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- How to describe the world in logic
- Moving as a way to gather new facts
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Coming next

- Making sound inferences
- Walking forward from the assumptions
- Walking backwards from the conclusion