

$$1 \text{ (a)} \quad V \cap W = \{\{1\}\}$$

$$V \cup W = \{\{1\}, \{2\}, 1, 2\}$$

$$P W = \{\emptyset, \{\{1\}\}, \{1\}, \{2\}, \{\{1\}, 1\}, \{\{1\}, 2\}, \{1, 2\}, \{\{1\}, 1, 2\}\}$$

$$V \cap P W = \{\{1\}, \{2\}\}$$

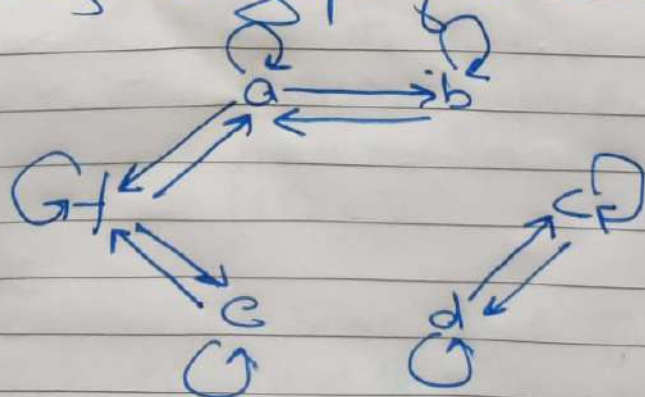
$$V \times W = \{\langle \{1\}, \{1\} \rangle, \langle \{1\}, 1 \rangle, \langle \{1\}, 2 \rangle, \langle \{2\}, \{1\} \rangle, \langle \{2\}, 1 \rangle, \langle \{2\}, 2 \rangle\}$$

$$V \Delta W = \{\{2\}, 1, 2\}$$

b Since  $R$  is reflexive,  $\langle a, a \rangle, \langle b, b \rangle, \langle c, c \rangle, \langle d, d \rangle, \langle e, e \rangle, \langle f, f \rangle$  must belong to  $R$

Since it is symmetric,  $\langle b, a \rangle, \langle f, a \rangle, \langle c, d \rangle, \langle f, e \rangle$  must belong to  $R$ .

Drawing the graph of the elements we know  $\in R$



So to make  $R$  transitive, we must have  $\langle e, a \rangle, \langle e, b \rangle, \langle f, b \rangle, \langle a, e \rangle, \langle b, f \rangle, \langle b, e \rangle$

Notice that symmetry is still maintained

Hence, we need to add the following pairs:

$\langle a, a \rangle, \langle b, b \rangle, \langle c, c \rangle, \langle d, d \rangle, \langle e, e \rangle, \langle f, f \rangle, \langle b, a \rangle, \langle f, a \rangle, \langle c, d \rangle, \langle f, e \rangle, \langle e, a \rangle, \langle e, b \rangle, \langle f, b \rangle, \langle a, e \rangle, \langle b, f \rangle, \langle b, e \rangle$



Page \_\_\_\_\_

(c) Take an arbitrary  $x \in (A \cup B) \cap (C \cup D)$ .  
Then by def. of intersection, we have ' $x \in (A \cup B)$ '  
and ' $x \in (C \cup D)$ '. So, we have both ' $x \in (A \cup B)$ '  
and ' $x \in (C \cup D)$ '. By def. of union, we have  
' $x \in A$  or  $x \in B$ ' and ' $x \in C$  or  $x \in D$ '. We have 4 cases.

(' $x \in A$  and  $x \in C$ ): Using 'and', we get ' $x \in A$ '. Then,  
' $x \in A$  or  $x \in C$ '. By definition of union, ' $x \in (A \cup C)$ '.  
Again, by def. of union, ' $x \in (A \cup C) \cup (B \cap D)$ '.

(' $x \in A$  and  $x \in D$ ): Using 'and', we get ' $x \in A$ '. Proceeding  
similarly as above, we get ' $x \in (A \cup C) \cup (B \cap D)$ '.

(' $x \in B$  and  $x \in C$ ): Using 'and', we get ' $x \in C$ '. Then,  
' $x \in A$  or  $x \in C$ '. By def. of union, ' $x \in (A \cup C)$ '. Again,  
by def. of union, ' $x \in (A \cup C) \cup (B \cap D)$ '.

(' $x \in B$  and  $x \in D$ ): By def. of intersection we get  
' $x \in B \cap D$ '. Then ' $x \in (A \cup C)$  or  $x \in (B \cap D)$ '. Hence, we  
have ' $x \in (A \cup C) \cup (B \cap D)$ '.

In all the four cases we have ' $x \in (A \cup C) \cup (B \cap D)$ '.  
Hence, we conclude that ' $x \in (A \cup C) \cup (B \cap D)$ '. Further  
as  $x$  was an arbitrary element of  $(A \cup B) \cap (C \cup D)$ ,  
we have ' $(A \cup B) \cap (C \cup D) \subseteq (A \cup C) \cup (B \cap D)$ '.



(d) To show  $(\overline{R \cup S})^{-1} \subseteq \overline{R^{-1}} \cap \overline{S^{-1}}$

Take arbitrary  $x, y$  such that ' $x(\overline{R \cup S})^{-1}y$ '.  
Then by def. of inverse, ' $y(\overline{R \cup S})x$ '. So by  
def. of complement, ' $\langle y, x \rangle \notin (R \cup S)$ '. By De  
Morgan's law, we get ' $\langle y, x \rangle \notin R$  and  
' $\langle y, x \rangle \notin S$ '. By def. of complement, we get

Suppose ' $\langle x, y \rangle \in R^{-1}$ '. Then ' $\langle y, x \rangle \in R$ ' by def. of  
inverse, which is a contradiction. So we conclude  
' $\langle x, y \rangle \notin R^{-1}$ '. Similarly, ' $\langle x, y \rangle \notin S^{-1}$ '. Then by  
def. of complement, we have ' $x \notin R^{-1}$  and  $x \notin S^{-1}$ '.  
So, by def. of intersection, we have ' $x \notin \overline{R^{-1}} \cap \overline{S^{-1}}$ '.

As  $x$  was an arbitrary element of  $(\overline{R \cup S})^{-1}$ , we conclude,  
 $(\overline{R \cup S})^{-1} \subseteq \overline{R^{-1}} \cap \overline{S^{-1}}$

To show  $\overline{R^{-1}} \cap \overline{S^{-1}} \subseteq (\overline{R \cup S})^{-1}$

Since De Morgan's law is stated in both  
directions, this follows from the above proof,  
reversing the steps in reverse order.

• To show  $(\overline{R \cup S})^{-1} = \overline{R^{-1}} \cap \overline{S^{-1}}$

We have ' $(\overline{R \cup S})^{-1} \subseteq \overline{R^{-1}} \cap \overline{S^{-1}}$ ' and ' $\overline{R^{-1}} \cap \overline{S^{-1}} \subseteq (\overline{R \cup S})^{-1}$ '. Hence, by definition of set equality,  
we have  $(\overline{R \cup S})^{-1} = \overline{R^{-1}} \cap \overline{S^{-1}}$

• To show  $\overline{R^{-1}} \cap \overline{S^{-1}} = (\overline{R \cup S})^{-1}$

We have ' $\overline{R^{-1}} \cap \overline{S^{-1}} \subseteq (\overline{R \cup S})^{-1}$ ' and ' $(\overline{R \cup S})^{-1} \subseteq \overline{R^{-1}} \cap \overline{S^{-1}}$ '. Hence, by def. of set equality,  
we have  $\overline{R^{-1}} \cap \overline{S^{-1}} = (\overline{R \cup S})^{-1}$



c To show  $R$  is reflexive:

Take any  $p, q \in \mathbb{N}$ . We have ' $p \times 1 = p$ ' and ' $q \times 1 = q$ '. But since  $1 \in \mathbb{N}$ , 'there exists  $n, m \in \mathbb{N}$  such that  $p \times n = p$  and  $q \times m = q$ '. Hence, we have ' $\langle p, q \rangle R \langle p, q \rangle$ '

To show  $R$  is transitive:

Assume that ' $\langle a, b \rangle R \langle c, d \rangle$ ' and ' $\langle c, d \rangle R \langle e, f \rangle$ '. Then by def. of  $R$ , there exists  $p, q, r, s \in \mathbb{N}$  such that ' $a \times p = c$ ' and ' $b \times q = d$ ' and ' $c \times r = e$ ' and ' $d \times s = f$ '. Using ' $a \times p = c$ ' and ' $c \times r = e$ ', we have ' $a \times (p \times r) = e$ '. Using ' $b \times q = d$ ' in ' $d \times s = f$ ', we have ' $b \times (q \times s) = f$ '. As we have both ' $a \times (p \times r) = e$ ' and ' $b \times (q \times s) = f$ ', and ' $p \times r \in \mathbb{N}$ ' and ' $q \times s \in \mathbb{N}$ ', by def. of  $R$ , we get ' $\langle a, b \rangle R \langle e, f \rangle$ '

To show  $R$  is anti symmetric

Assume that ' $\langle a, b \rangle R \langle c, d \rangle$ ' and ' $\langle c, d \rangle R \langle a, b \rangle$ ' both hold. Then there exists  $p, q, r, s \in \mathbb{N}$  such that ' $a \times (p \times r) = a$ ' and ' $b \times (q \times s) = b$ '. (Proof is same as case for transitive). As we have  $p, q, r, s \in \mathbb{N}$ , we conclude  $p = q = r = s = 1$ . Hence  $e = a \times p = a$  and  $d = b \times q = b$ . So ' $\langle a, b \rangle = \langle c, d \rangle$ '.

As  $R$  is reflexive, transitive and anti symmetric, it follows that  $R$  is a partial order.



2a i) Not a function as  $\langle 1, 1 \rangle$  and  $\langle 1, -1 \rangle$  both belong to the function, so it has 1 has 2 images.

ii) This is a function. (It is a total function). It is injective and surjective.  $f^{-1}(y) = \left(\frac{1}{y} + 8\right)^{1/3}$

iii) This is a total function. Not injective as  $\langle -1, 1 \rangle$  and  $\langle 1, 1 \rangle$  both belong to it. Not surjective as no  $x$  gets mapped to a negative number. No inverse.

iv) It is a total function. It is injective. It is not surjective as  $\frac{1}{2}$  is not part of the range. No inverse.

b) Given  $|A| = m$  and  $|B| = n$ .  
 $|P(A)| = 2^m$  and  $|P(B)| = 2^n$

Considering total functions from  $P(A)$  to  $P(B)$ , each element in  $P(A)$  can choose among  $|P(B)|$  different values for its image. Moreover, there are  $|P(A)|$  elements in  $P(A)$ , so we make this choice  $|P(A)|$  times.

$$\text{So no. of total functions} = |P(B)|^{|P(A)|} = (2^n)^{(2^m)} = 2^{n \cdot 2^m}$$

When considering partial functions, each element in  $P(A)$  has an additional choice to remain undefined. So each element in  $P(A)$  has  $(1 + |P(B)|)$  choices.

$$\text{So no. of partial functions} = (|P(B)| + 1)^{|P(A)|} = (2^n + 1)^{(2^m)}$$



(c) Take any  $a \in \mathbb{N}$ . Then  $\{a\} \in F$ .  $\{a\}$  is a minimal element of  $(F, R)$ . Take any  $V \in F$  such that  $V \not\subseteq \{a\}$ . Then ' $V = \emptyset$ ' or ' $V = \{a\}$ '. Since these are the only 2 subsets of  $\{a\}$ . Moreover, by def. of  $R$ , we have ' $V \neq \emptyset$ '. So ' $V = \emptyset$ ' cannot hold and we must have ' $V = \{a\}$ '.

So we have shown that (since  $V$  was arbitrary)  
 $\forall V \in F. (V R \{a\} \Rightarrow V = \{a\})$

Note that minimal element is not unique.  
 $\{1\}, \{2\}, \{3\}, \dots$  are all minimal

No least element exists. Take any minimal element  $\{a\}$ . Then take ' $b = a + 1$ ' and ' $c = a + 2$ ', then we have ' $\neg (\{a\} R \{b, c\})$ '. Hence, ' $\forall b \in F (\{a\} R b)$ ' does not hold for any  $a$ .

No maximal element exists as well. Suppose towards a contradiction that ' $V = \{a_1, a_2, \dots, a_n\}$ ' is a maximal element (note that  $V$  is finite). Define  $W = V \cup \{(\sum_{i=1}^n a_i) + 1\}$ . Then  $V \neq W$ . Further,  $V \subseteq W$  and  $V \neq \emptyset$ . So we have  $V R W$ . By def. of  $V$  being maximal, we get  $V = W$ , contradiction to ' $V \neq W$ '. Hence no maximal element can exist.

By similar reasoning, no greatest element exists as well.



(d) Given  $A_1 \approx A_2$  and  $B_1 \approx B_2$ . So there exists bijective functions  $f: A_1 \rightarrow A_2$  and  $g: B_1 \rightarrow B_2$

Define  $h: (A_1 \times B_1) \rightarrow (B_2 \times A_2)$  as follows  
 $h(a, b) = \langle g(b), f(a) \rangle$

This function is well defined as  $g(b) \in B_2$  and  $f(a) \in A_2$ , so  $\langle g(b), f(a) \rangle \in B_2 \times A_2$ . We will show  $h$  is bijective.

( $h$  is injective): Assume that ' $h(a_1, b_1) = h(a_2, b_2)$ '. Then, by def. of  $h$ , we have ' $\langle g(b_1), f(a_1) \rangle = \langle g(b_2), f(a_2) \rangle$ '. Then, we have ' $g(b_1) = g(b_2)$ ' and ' $f(a_1) = f(a_2)$ '. As  $f, g$  are both bijective and hence injective, we get ' $b_1 = b_2$ ' and ' $a_1 = a_2$ '. So we have ' $\langle a_1, b_1 \rangle = \langle a_2, b_2 \rangle$ '. So,  $h$  is injective.

( $h$  is surjective): Take any ' $\langle b_2, a_2 \rangle \in B_2 \times A_2$ '. Then  $b_2 \in B_2$  and  $a_2 \in A_2$ . As  $f, g$  are bijective and hence surjective, there exists  $a_1 \in A_1$  such that ' $f(a_1) = a_2$ ', and there exists  $b_1 \in B_1$  such that ' $g(b_1) = b_2$ '. But then ' $h(a_1, b_1) = \langle g(b_1), f(a_1) \rangle = \langle b_2, a_2 \rangle$ '. So there exists ' $\langle a, b \rangle \in (A_1 \times B_1)$ ' such that ' $h(a, b) = \langle a_2, b_2 \rangle$ '. Hence  $h$  is surjective.

Hence  $h$  is bijective. But then 'there exists a ~~function~~ bijective function  $h: (A_1 \times B_1) \rightarrow (B_2 \times A_2)$ '. So by def. we have  $A_1 \times B_1 \approx B_2 \times A_2$

Hence  $A_1 \times B_1 \approx B_2 \times A_2$



(e) i) Let  $f: \mathbb{N} \rightarrow \{1\}$  be a function. Then, for all  $x \in \mathbb{N}$ , we have  $f(x) = 1$ . So, there can only be 1 such function.

So the set of all functions from  $\mathbb{N}$  to  $\{1\}$  is finite and hence countable by definition.

ii) When considering partial functions, each element can be undefined as well. So we can write any such function as  $f: \mathbb{N} \rightarrow \{0, 1\}$ , and the set of such functions as  $\{0, 1\}^{\mathbb{N}}$ .

Using exercise 65, we know that  $\{0, 1\}^{\mathbb{N}} \approx p(\mathbb{N})$ .

Suppose towards a contradiction that  $\mathbb{N} \approx \{0, 1\}^{\mathbb{N}}$ . Then by transitivity of  $\approx$ , we have  $\mathbb{N} \approx p(\mathbb{N})$ , which by exercise 66, we know is a contradiction. Hence we have  $\mathbb{N} \not\approx \{0, 1\}^{\mathbb{N}}$ .

So, the set of all partial functions from  $\mathbb{N}$  to  $\{1\}$  is not countable.