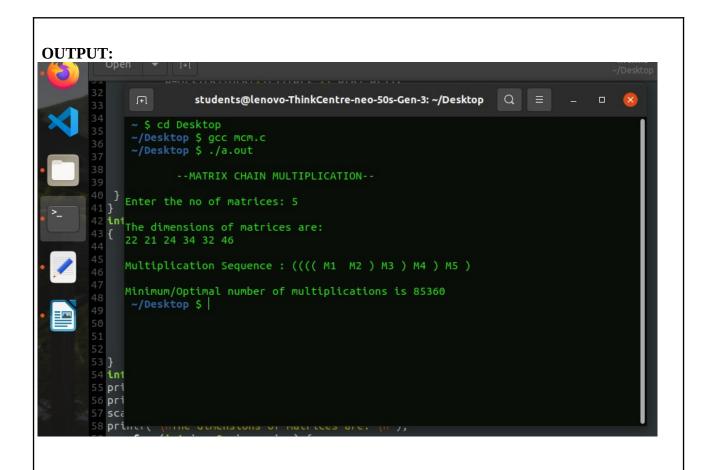
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Experiment No.	5

AIM:	Implement matrix chain multiplication using dynamic programming	
PROGRAM		
THEORY:	<ul> <li>matrices of size m*n and n*p when multiplied, they generate a matrix of size m*p and the number of multiplications performed are m*n*p.</li> <li>Now, for a given chain of N matrices, the first partition can be done in N-1 ways. For example, sequence of matrices A, B, C and D can be grouped as (A)(BCD), (AB)(CD) or (ABC)(D) in these 3 ways.</li> <li>So a range [i, j] can be broken into two groups like {[i, i+1], [i+1, j]}, {[i, i+2], [i+2, j]},, {[i, j-1], [j-1, j]}.</li> <li>Each of the groups can be further partitioned into smaller groups and we can find the total required multiplications by solving for each of the groups.</li> <li>The minimum number of multiplications among all the first partitions is the required answer.</li> </ul>	
INPUT:	Input nuumber of matrices and generate random dimensions	

```
MatrixChainMultiplication(mat[], low , high):
ALGORITHM:
                        // 0 steps are required if low equals high
                        // case of only 1 sized sub-problem.
                        If(low=high):
                          return 0
                        // Initialize minCost to a very big number.
                        minCost = Infinity
                        // Iterate from k = low to k = high-1
                        For(k=low to high-1):
                           Cost = Cost of Multiplying chain on left side +
                               Cost of Multiplying chain on right side +
                               Cost of Multiplying matrix obtained from left
                               and right side.
                           */
                          cost=MatrixChainMultiplication(mat, low, k)+
                             MatrixChainMultiplication(mat, low+1, high)+
                             mat[low-1]*mat[k]*mat[high]
                          // Update the minCost if cost<minCost.
                          If(cost<minCost):</pre>
                             minCost=cost
                        return minCost
```

```
}
}
void multiply(){
int q,k;
for(int i=n;i>0;i--)
  for(int j=i;j \le n;j++)
   if(i==j)
    m[i][j]=0;
   else
     for(int k=i;k \le j;k++)
      q=m[i][k]+m[k+1][j]+p[i-1]*p[k]*p[j];
      if(q{<}m[i][j])
        m[i][j]=q;
        s[i][j]=k;
int chain(int p[], int i, int j)
  if(i == j)
     return 0;
   int k,min=INT_MAX,count=0;
  for (k = i; k < j; k++) {
     count = chain(p, i, k) + chain(p, k + 1, j) + p[i - 1] * p[k] * p[j];
     if (count < min)
        min = count;
   }
   return min;
int main(){
printf("\n\t--MATRIX CHAIN MULTIPLICATION--\n")
```

```
printf("\nEnter the no of matrices: ");
scanf("%d",&n);
printf("The dimensions of matrices are: \n");
  for (int i = 0; i \le n; i++) {
     p[i] = (rand()\%(46 - 15 + 1)) + 15;
     printf("%d ", p[i]);
for(int i=1;i<=n;i++)
for(int j=i+1;j<=n;j++)
m[i][i]=0;
m[i][j]=INT\_MAX;
s[i][j]=0;
multiply();
printf("\nMultiplication Sequence : ");
print(1,n);
printf("\nMinimum/Optimal\ number\ of\ multiplications\ is\ %d\n",chain(p,\ 1,\ n));
return 0;
```



## RESULT ANALYSIS:

**Time Complexity** - We are using three nested for loops, each of which is iterating roughly O(n) times. Hence, the overall time complexity is O(n3).

**Space Complexity -** We are using an auxiliary dp array of dimensions,  $(n-1)\times(n-1)$  hence space complexity is O(n2)

In the matrix chain multiplication problem, the minimum number of multiplication steps required to multiply a chain of matrices has been calculated.

• Determining the minimum number of steps required can highly speed up the

	<ul> <li>multiplication process.</li> <li>It takes O(n3) time and O(n2) auxiliarly space to calculate the minimum number of steps required to multiply a chain of matrices using the dynamic programming method.</li> <li>It is very important to decide the order of multiplication of matrices to perform the task efficiently.</li> <li>So the minimum number of multiplication steps required to multiply a chain matrix of length n has been computed.</li> <li>Trying out all the possibilities recursively is very time consuming hence the method of dynamic programming is used to find the same.</li> <li>The minimum number of steps required to multiply a chain of matrices can be found in O(n3) time complexity using dynamic programming.</li> </ul>
CONCLUSION:	In this experiment I understood about how to find the minimalcost of matrix chain multiplication using dynamic programming