

# *Design and Implementation of a SHA-256 Co-Processor*

Aditi Prakash

Department of Electrical Engineering  
Hochschule Darmstadt  
Schöfferstr 3, 64295 Darmstadt, Germany  
stadprak@stud.h-da.de

**Abstract**— Exponential rise in technological developments in the recent era has introduced a dire necessity of security for ensuring the authenticity of the information. Accordingly, the demand for cryptographic hash functions have increased. A hash function is a cryptographic algorithm without a key such as MD5, SHA-1 (Secure Hash Algorithm) and SHA-256. Different types of cryptographic algorithms have been developed in order to improve the performance of these information-protecting procedures. In this paper, software design and hardware design of SHA-256 is implemented and its performance is compared with the co-processor designed using NIOS II CPU. SHA-256 is a one-way hash function characterized by being highly secure and fast while having a high collision resistance. This paper gives an overview about cryptography and discusses in detail the design and implementation of SHA-256 on software and hardware, and as a co-processor. Further, the application of SHA-256 in blockchain and cryptocurrency mining is discussed.

## I. INTRODUCTION

Cryptography is the process of converting ordinary plain information into unintelligible text or vice-versa. It is a method of storing and transmitting data in a particular form so that only those for whom it is intended can read and process it. It not only protects data from theft or alteration but can also be used for user authentication. Data can be protected either by encryption or by hashing. While Encryption is a two-way function; Information can be retrieved if the key used to encrypt the data is known, hashing is one-way function; the input cannot be retrieved after hashing is performed.

Hashing is the practice of using an algorithm to map data of any size to a fixed length. This fixed length output is called hash-value. Such one-way functions have typical applications in Pseudo Random Generators, Message Authentication Codes, Public Key Cryptography, IFF (Identification, Friend or Foe Detection) etc...

Hash functions are used to protect data, verify a digital signature, or authenticate a wide range of online processes and transactions, using complex logical operations. The hash functions have unique properties making them highly secure. They are one-way functions, meaning for a given hash output  $h$ , it is computationally infeasible to retrieve the input message  $x$ , where  $H(x) = h$ . They also have a very high collision resistance, that is it is computationally infeasible to

have two inputs with the same hashed output ( $x \neq y$  and  $H(x) \neq H(y)$ ).

Besides, hash functions are highly sensitive. A slight change in the input message will completely change the output hash.

Although the input to a hash function can be of any length, the length of the output is finite. This means that at some point of time when huge amount of data is being hashed, the output combination ought to repeat for a different input data as well. Albeit, collisions exist and cannot be completely avoided.

This paper explores the design and implementation of SHA-256 algorithm. Initially a pure software approach of SHA-256 is explored by designing the algorithm in C language. Later, it is implemented with System Verilog with an appropriate testbench. Simulation results are analyzed with Model Sim Altera tool. Lastly, a wrapper is designed for the SHA-256 kernel to interface it with a NIOS II CPU and is modelled to behave as a co-processor/accelerator when connected with the hardware.

This paper is organized as follows: In section 2, the principle and working of SHA 256 is introduced; design of SHA-256 is discussed in section 3; and finally results of the design and application of SHA256 in currency mining are presented in section 4.

## II. PRINCIPLE OF SHA-256

### A. SHA-256 Functions and constants

SHA-256 algorithm uses a series of Rotate and Shift operations throughout the computation in order to calculate the hash function. These operations are combined together in different combinations to obtain six functions that are used in the algorithm at different stages.

Below is a short description of the functions:

$$\sigma_0(x) = \text{ROTR7}(x) \wedge \text{ROTR18}(x) \wedge \text{SHR3}(x) \quad (1)$$

$$\sigma_1(x) = \text{ROTR17}(x) \wedge \text{ROTR19}(x) \wedge \text{SHR10}(x) \quad (2)$$

$$\Sigma_0(x) = \text{ROTR2}(x) \wedge \text{ROTR13}(x) \wedge \text{ROTR22}(x) \quad (3)$$

$$\Sigma_1(x) = \text{ROTR6}(x) \wedge \text{ROTR11}(x) \wedge \text{ROTR25}(x) \quad (4)$$

Choice ( $x, y, z$ )  $\rightarrow$  Decides the second input based on the first input (e.g., if  $x = 1$ , choose  $y$  and if  $x = 0$  then choose  $z$ )

Majority ( $x, y, z$ )  $\rightarrow$  The result is the majority of the three bits

The constant  $K_t$  in SHA-256 is used to mix the message that is put into hash function. They are cube roots of prime numbers. This is done for the first 64 prime numbers, i.e.,  $0 \leq t \leq 63$  for  $K_t$ . The initial hash value for the first round of compression is stored into 8 registers  $a$  to  $h$ . They are initialized with the square root of first 8 prime numbers.

Two temporary hash values are required which are calculated as below:

$$T_i = \Sigma_1(e) + \text{Choice}(e, f, g) + h + K_i + W_i \quad (5)$$

$$T_i = \Sigma_0(a) + \text{Maj}(a, b, c) \quad (6)$$

### B. Message Padding

Before starting the actual computation of the hash value, data has to be padded since the hash function hashes the data in chunks of 512 bits at a time. Therefore, input message of length 'L' is padded with a '1' that indicates the termination of message input or acts as a separator, followed by 'k' zero bits where k is the smallest non-negative solution to  $L+1+k = 448 \pmod{512}$ , then the length of the message 'L' is expressed in binary in the last 64 bits.

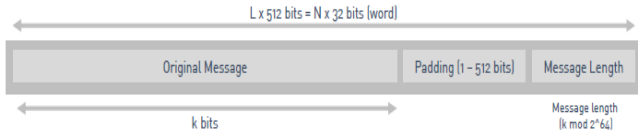


Fig 1: Structure of message after padding

If the message is longer than 512 bits then it will be padded to the next least multiple of 512 i.e., 1024. Later the padded message is divided into multiple blocks of 512 bits.

### C. Creation of Message Schedule

A message schedule,  $W_t$  needs to be created from each message block. It is required that the schedule is 64 words long.

For  $0 \leq t \leq 15$ , a message schedule is directly created from the input message while for  $16 \leq t \leq 63$ , a message schedule is calculated by using the following expression:

$$W_t = \sigma_1[W(t-2)] + W(t-7) + \sigma_0[W(t-15)] + W(t-16) \quad (7)$$

### D. Compression

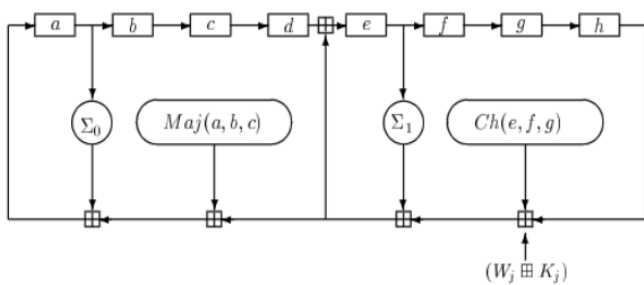


Fig 2: Compression in SHA-256

Compression uses four functions (Equation (3), (4), Choice and Majority) for its computation. The initial hash value is assigned to 8 variables and 64 iterative operations are performed along with constants ( $K_t$ ) and message schedule ( $W_t$ ) to obtain the final value of these 8 variables.

### E. Final Hash Creation

The computed hash values are added with the initial hash value and is fed as an input to the next block. The final block provides the hash value of 256 bits.

The below image provides an overall picture of the SHA-256 computation. C here in the image indicates the compression block

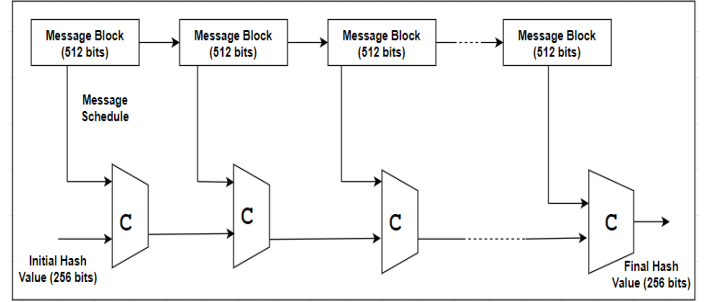


Fig 3: Block Diagram of SHA-256 operation

## III. DESIGN AND IMPLEMENTATION

### A. Software Design

The Software Design of SHA-256 is a simple C program that reads the input string from the user and provides the hash values as output.

The below image shows the flow of the C code.

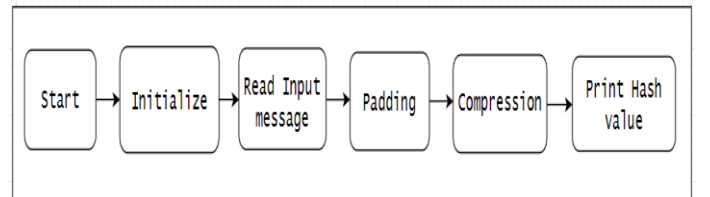


Fig 4: Software Block Diagram

### B. Hardware Design

Hardware Description Language System Verilog is used for the Hardware design. The computation of SHA-256 is designed using a State Machine. The SHA-256 kernel is clock synchronous with a clock of 50Mhz and performs computations on every positive edge. The processing of SHA-256 takes 176 cycles. The behaviour is verified by writing an appropriate test bench for the implementation.

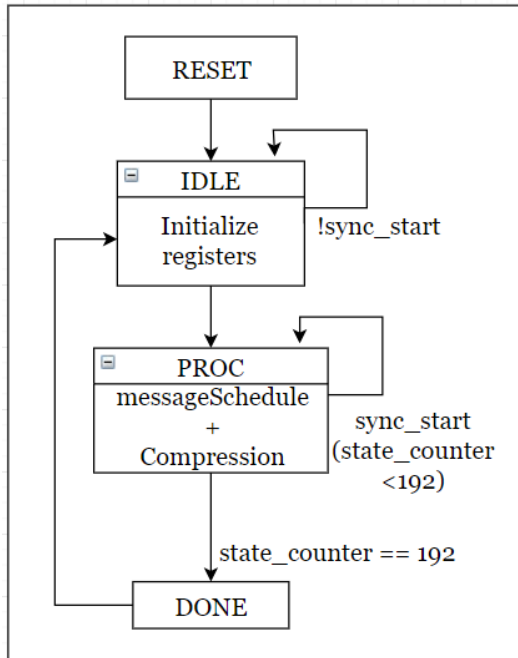


Fig 5: State machine for SHA-256

Below is a table that represents the number of clock cycles consumed by each stage in the processing state.

	No. of Clock cycles	Time (in ns)
Calculation of Message Schedule	48	960ns
Compression	128	2560
<b>Total</b>	<b>176</b>	<b>3520</b>

Table 1: Table depicting clock cycles

### C. Hardware Software Co-Design

The goal is to design an accelerator that offloads the CPU and performs the SHA-256 computation independent of the CPU. This reduces CPU load and speeds up the processor. For this, a NIOS II CPU is taken as a reference for the implementation.

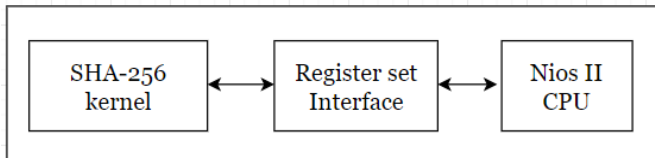


Fig 6: Block diagram of Accelerator interfaced with CPU

A testbench is designed in a such a way that it emulates the CPU behavior. A register module is used as an interface between the SHA-256 kernel and the CPU to provide the necessary interface.

Register set consists of control register, status register, input and output registers as shown below. A sum total of 26 registers, each of 32 bit is required for the design.

```

/* Register set
 * 1. CTRL REGISTER           : Control bits to start the sha256 processing
 * 32'd1 = start
 * 2. STATUS REGISTER        : Register to indicate completion of sha 256
 * 32'd1 = done
 * 3. DATA_IN               : 16x32 bit registers (512 bits) for input data
 * 4. DATA_OUT              : 8x32 bit registers (256 bits) for output data
 */
  
```

Fig 7: Register set used for the accelerator design

The testbench always selects address of the appropriate register before reading or writing the data.

Initially, the input message is stored into the registers. Upon receiving start signal from the testbench, SHA kernel starts the processing and stores the hash value in output registers after processing.

The testbench then reads the output hash value from the registers after a certain duration.

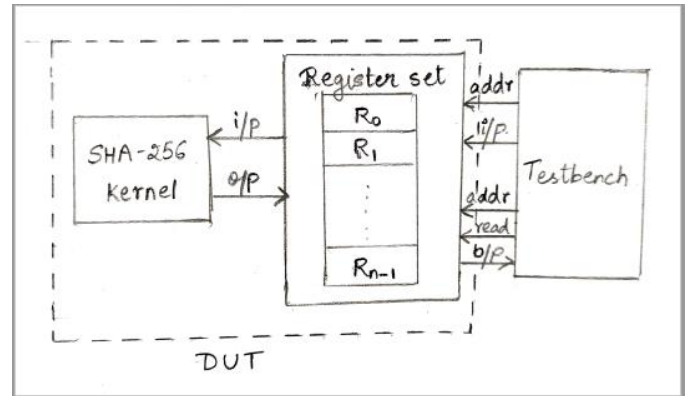


Fig 8: Block diagram of Hardware Software Co-Design

## IV. RESULTS AND DISCUSSION

This paper implements SHA-256 in software, hardware and as an accelerator to the NIOS II CPU. All the designs are tested with different inputs and each of them verified for their authenticity.

Below are the results of the tests performed for one of the input messages.

```

Console Problems Executables Debugger Console
<terminated> (exit value: 0) SHA_256.exe [C/C++ Application] C:\Assignment\SHA_256\SHA_256(Debug)\SHA_256.exe (08/03/21,
Enter the string
hello world
b9 4d 27 b9 93 4d 3e 8 a5 2e 52 d7 da 7d ab fa c4 84 ef e3 7a 53 80 ee 90 88 f7 ac e2 ef cd e9
  
```

Fig 9: Hash value output of C program

