## Implement and improve an efficient, layered tape with prefetching capabilities

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## Tape in Clad

The tape is a stack-like data structure that stores intermediate values in reverse mode AD during the forward pass for use during the backward (gradient) pass

# Previous Implementation of Tape

Clad used a monolithic memory buffer. If the tape was full while pushing then it reallocated double the capacity.

```
template <class InputIt, class NoThrowForwardIt>
CUDA HOST DEVICE void MoveData(InputIt first, InputIt last,
                                NoThrowForwardIt d first) {
  NoThrowForwardIt current = d first:
  for (; first != last; ++first, (void)++current) {
    ::new (const cast<void*>(
     static cast<const volatile void*>(clad addressof(*current))))
      T(std::move(*first));
 onstexpr static std::size t init capacity = 32;
CUDA HOST DEVICE void grow() {
  # (! capacity)
    capacity = init capacity;
  T* new data = AllocateRawStorage( capacity);
  if (Inew data) {
   destroy(begin(), end());
    printf("Allocation failure during tape resize! Aborting.\n");
    trap(EXIT FAILURE);
```

#### New Implementation of Tape

The new tape follows a slab-based structure with small buffer optimization. The slab size and buffer size are configurable template parameters with default values.

```
template <typename T. std::size t SBO SIZE = 64. std::size t SLAB SIZE = 1024>
         s(T) char raw data[SLAB SIZE * sizeof(T)]();
   Slab* next;
   CUDA HOST DEVICE Slab() : next(nullptr) ()
   CUDA HOST DEVICE T* elements() {
    cplusplus >= 201703L
        mich std::launder(reinterpret cast<T+>(raw data));
      return reinterpret cust<T*>(raw data):
       is(T) char m static buffer[SBO SIZE * suzunf(T)]{};
 bool m using sbo = true;
 Slab* m head = nullptr;
 std: size t m size = 0;
 CUDA HOST DEVICE T* sbo elements() {
    cplusplus >= 201703L
       orn std::launder(reloterpret cost<T*>(n static buffer));
   return reinterpret casts[*>(m static buffer):
 CUDA HOST DEVICE const T+ sho elements() const {
    colusplus >= 201703L
       urn std. launder(reinterpret cossistens) T*>(n static buffer));
       urn reinterpret cust<comst T'>(m static buffer);
```

```
template <typename... ArgsT>
CUDA HOST DEVICE void emplace back(ArgsT&&... args) {
   (m size < SBO SIZE) {
   // Store in SBO buffer
::new (const cast<void*>(static cast<const volatile void*>(
        sbo elements() + m size))) T(std::forward<ArgsT>(args)...);
    if (m using sbo)
     m using sbo = false:
    if ((m size - SBO SIZE) % SLAB SIZE == 0) {
      Slab* new slab = new Slab();
       (!m head) {
        m head = new slab:
        Slab* last = m head;
         hile (last->next)
          last = last->next;
        last->next = new slab;
    Slab* slab = m head;
    std::size t idx = (m size - SBO SIZE) / SLAB SIZE:
     mile (idx--)
      slab = slab->next;
   ::new (const cast<void*>(static cast<const volatile void*>(
        slab->elements() + ((m size - SBO SIZE) % SLAB SIZE))))
        T(std::forward<ArgsT>(args)...);
 m size++;
```

#### Enhanced Benchmarks

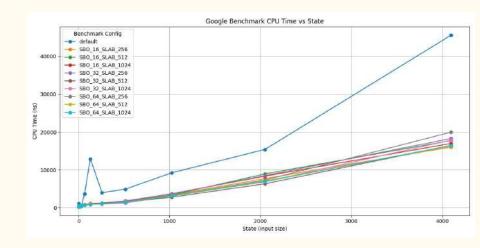
Configurable tape memory benchmarks added to test different configurations and find optimal slab and buffer size

```
template <std::size t SBO SIZE, std::size t SLAB SIZE>
static void BM TapeMemory Templated(benchmark::State& state) {
  int block = state.range(0);
  AddBMCounterRAII MemCounters(*mm.get(), state);
  clad::tape<double, SBO SIZE, SLAB SIZE> t;
  for (auto : state)
      func<double, SBO SIZE, SLAB SIZE>(t, 1, block * 2 + 1);
}

#define REGISTER TAPE BENCHMARK(sbo, slab)
      BENCHMARK TEMPLATE(BM TapeMemory Templated, sbo, slab)
      ->RangeMultiplier(2)
      ->Range(0, 4096)
      ->Iterations(1)
      ->Name("BM TapeMemory/SBO_" #sbo " SLAB_" #slab)

REGISTER TAPE BENCHMARK(64, 1024);
REGISTER TAPE BENCHMARK(32, 512);
```

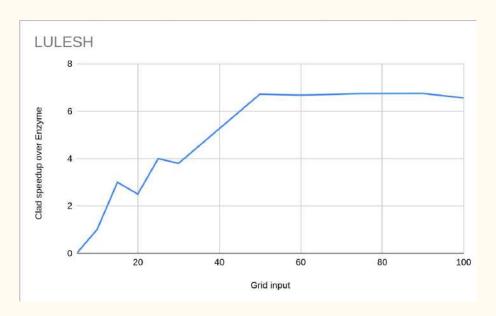
Slab size 64 and buffer size 1024 performed the best consistently over many runs



#### Lulesh Benchmarks

grid size = 50: clad is 6x faster than enzyme

grid size = 100: clad is 8.6x faster than enzyme



### Current Progress

- Benchmark script to compare two revisions (PR #1394)
- Modified tape structure to slab-based (PR #1404)
- Added small buffer optimization (PR #1404)
- Enhanced benchmarks (PR #1404)

#### Future Work

- Add thread safety to tape
- Add offloading mechanism
- Implement CPU-GPU transfer (Stretch goal)
- Implement Checkpointing (Stretch goal)

## Thank You!