Homework - 2

Q.1.

a. Code attached to submission

1.) (b) fitness = 1 + individuals
Goal: calculate the probability that a given alternate allele present at frequency 1/2N in generation 0 becomes extinct in generation 1.
There is exactly 1 copy of the alternate allele in a population of size N. [Diploid >> 2N total alleles]
The probability of the allele not being selected is 1 - 1 2N
So, it is the probability of it not being selected to be passed on to generation I in any of the 2N opportunities.
$P(\text{extinction}) = \left(1 - \frac{1}{2N}\right)^{2N}$
for N->00 this can be approximated as the exponential limit
$=> P(extinction) \approx \frac{1}{e}$
≈ 0.36787944

c. Output:

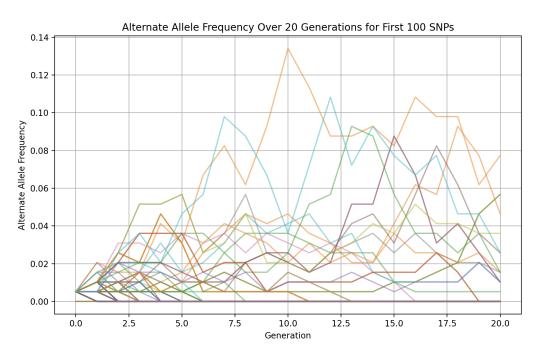
Estimated extinction probability: 0.3514

Theoretical extinction probability: 0.36787944117144233

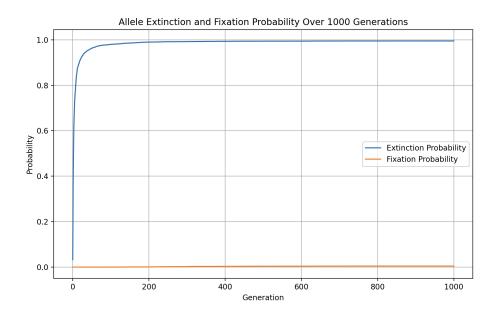
Extinction count: 3514, Total SNPs: 10000

The output derived approximately matches our expected value of extinction probability.

d.



e.



P(child does not inherit alt alleles)

=
$$\frac{1.5}{(N+0.5)} * 0.5 + (\frac{1.5}{N-0.5}) * 0.5$$

+ $\frac{1.5}{(N+0.5)} - \frac{1.5}{N-0.5}$

= $1 - (\frac{1.5}{N+0.5}) * 0.5 - \frac{1.5}{N-0.5}$

= $\frac{1.5}{(N+0.5)} * 0.5 - \frac{1.5}{(N-0.5)} * 0.5$

= $\frac{1.5}{(N-0.5)} * 0.5$

= $\frac{1.5}{(N-0$

f. Output obtained:

Estimated extinction probability at SNP42 after 1 generation: 0.237

g. Output obtained:

Estimated extinction probability at SNP42 after 100 generations: 0.015 Estimated fixation probability at SNP42 after 100 generations: 0.985

h. Output obtained:

Estimated extinction probability at SNP42 (deleterious) after 100 generations: 1.0 Estimated fixation probability at SNP42 (deleterious) after 100 generations: 0.0

Q.2.

- a. Code attached to submission.
- b. Output from code:

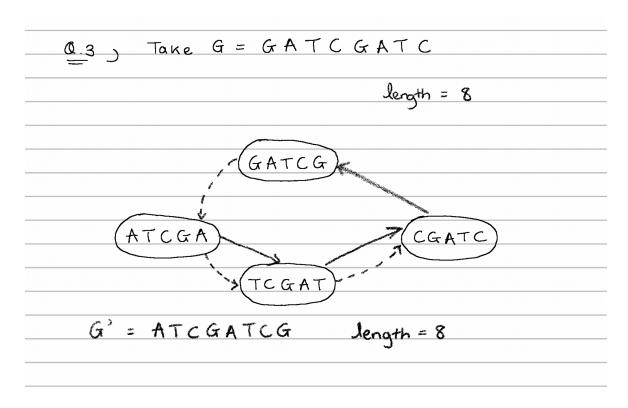
Number of SNPs with p-value < 0.05: 458 Expected number of SNPs with p-value < 0.05 by chance: 500.0

c. Output from code:

SNP	Uncorrected p_value	Corrected p-value	Diseases Odds ratio for Heterozygous individuals	Disease Odds ratio for homozygous alternate individuals
1000	1.732482e-16	1.732482e-12	2.000930	5.471660
2000	9.230119e-07	9.230119e-03	2.111610	7.867450
3000	5.629213e-10	5.629213e-06	2.046087	3.334257
4000	4.770372e-07	4.770372e-03	1.641808	4.201993

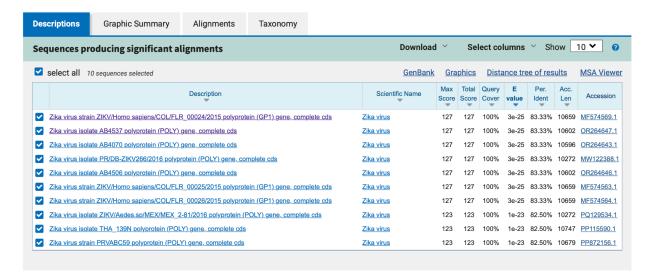
1101	a				II Z III CONCINGUICY_COD	
	SNP	uncorrected_p_value	corrected_p_value	odds_ratio_het	odds_ratio_homo_alt	
1000	SNP1000	1.732482e-16	1.732482e-12	2.000930	5.471660	
2000	SNP2000	9.230119e-07	9.230119e-03	2.111610	7.867450	
3000	SNP3000	5.629213e-10	5.629213e-06	2.046087	3.334257	
	SNP4000	4.770372e-07	4.770372e-03	1.641808	4.201993	
)(myenv) aditipotnis@Aditis−Air ~ % ■						

d. Even if two SNPs have similar disease odds ratios, the chi-squared value will depend on the actual counts in each cell of the contingency table, the expected counts under the null hypothesis, and how evenly distributed the genotypes are among diseased and non-diseased individuals. If one SNP has a better representation in each genotype category, the observed data may deviate less from the expected, resulting in a smaller chi-squared value and thus a higher p-value.



Q.4.

a. According to the BLAST results it is clear that the infection is a Zika Virus.



100% Query Cover means that your entire query sequence was aligned to the Zika virus genome. This indicates that the given infected section of the sequence matches the full length of these Zika virus sequences. The **E-value** is **3e-25** for most of the matches, which is extremely low. This indicates that the probability of this alignment occurring by chance is negligible. In other words, the match is highly significant and very unlikely to be random.

b. **Position 6–64**: The sequence comparisons highlight where the strains have identical nucleotides and where they differ. The given patient's strain (MF574569.1) appears to be **most similar to strain 4** for a large portion of the sequence. It is also quite similar to strain 5 but with more differences.

Tree Rendering results

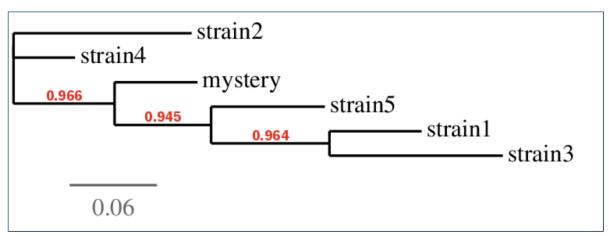


Figure 1: Phylogenetic tree (the branch length is proportional to the number of substitutions per site).

```
6 64
 MF574569.1 GCCTACCTTG AGAAGCAATC AGACACTCAA TATGTCTGCA AAAGATCCGC
 strain4
              GCCTACATTG AGAAACTATC ACATACAAAA TCTGTCAGCA AAAGATCTGC
 strain2 GCCTACATTG AAAAGCTATC AGATACAGAA CTTGTCAGCA AAAGAACTGC
 strain3 GCCTACCTTG AAAAGCTATT ACGCACAGGA TTTGTCAGCA GAAGGTCTGA
 strain5
              GCCTACTTTG AGAAGCTATC ATACACAGGA TTTGTCAGCA AAAGATCCGG
 strain1
              GCCTACCTTG AGAAGCTATT ATACACAGGA TTTGTCAGCA AAAGATCTGA
              GCTATCTCCC CAGG
              ACTATATGGC TAAG
              ACTATATGGC AAAG
              ACTATATGCC CAAG
              ACTATATCCC TAAG
              ACTATATGCC TAAG
Q.5.
   a.
def small parsimony tree(T, D, k, M):
  // given that:
  // T: A tree with n leaves and m internal nodes
  // D: Matrix (n x k) where each row is the k-dimensional trait vector for a species
  // M: Maximum possible value for any trait
  n = len(D) # number of leaves (species)
  // Step 1: Initialize dynamic programming table
  DP = \{\} # DP[u][i] will store minimum cost at node u for trait i
  Traceback = {} # Traceback[u][i] will store the best trait value for internal node u, trait i
  // Step 2: Assign leaf nodes with given trait vectors
  for leaf in T.leaves():
    DP[leaf] = \{\}
    for i in range(k):
      DP[leaf][i] = \{D[leaf][i]: 0\}  # No cost at leaves, trait values are fixed
  // Step 3: Bottom-up dynamic programming
  for node in T.postorder traversal():
    if node not in T.leaves():
      DP[node] = \{\}
      Traceback[node] = {}
      for i in range(k): # Iterate over each trait
         DP[node][i] = \{\}
        min cost = float('inf')
```

```
for val in range(M + 1): # Iterate over all possible trait values for this node
  cost = 0
  for child in T.children(node):
    # Find minimum cost of assigning val to current node's trait i
    min_child_cost = float('inf')
    best_val = -1
    for child_val in DP[child][i]:
        current_cost = DP[child][i][child_val] + abs(val - child_val)
        if current_cost < min_child_cost:
            min_child_cost = current_cost
            best_val = child_val
        cost += min_child_cost
        Traceback[node][i] = best_val</pre>
```

DP[node][i][val] = cost # Store the minimum cost for node's trait i having value 'val'

```
// Step 4: Traceback to find the best trait assignment for internal nodes
```

```
Du = {} # Du will store the best trait vector for each internal node
root = T.root()
Du[root] = [0] * k

# Initialize the root's trait vector based on DP values
for i in range(k):
    best_val = min(DP[root][i], key=DP[root][i].get) # Best trait value for trait i
    Du[root][i] = best_val

# Traceback to assign values to other internal nodes
for node in T.preorder_traversal():
    if node != root:
        parent = T.parent(node)
        Du[node] = [0] * k
        for i in range(k):
            Du[node][i] = Traceback[parent][i]
```

// Step 5: Return the trait vector assignment for each internal node return Du, DP[root]

- b. Time complexity can be derived by analyzing the 2 steps:
 - For the tree traversal the runtime will be O(n), as the tree is traversed twice (during DP and traceback) and the tree has n nodes (n species).
 - There are k traits.. For each node and each trait, you have O(M^2) operations (forward and backward passes through possible trait values).

With these considerations the time complexity is $O(n * k * M^2)$ assuming M is significant enough to consider else O(n * k).