

Assignment 2

Job Partitioning on two machines using Greedy Algorithm

Team members: NAME - EMAIL - UID

- Chris Davis Jaldi – jaldi.2@wright.edu – U01099335
- Aditya Mallakula – mallakula.2@wright.edu – U01093160
- Vanaja Uppala – uppala.19@wright.edu – U01080568

Greedy Heuristic Used:

This scheduling algorithm's greedy heuristic prioritizes the **earliest job finish time**:

- Jobs are scheduled by selecting the one that completes the quickest and can be placed on an available machine (prioritizing machine M1 over M2).
- By sorting the jobs by end time, the algorithm ensures that each job selected leaves the most amount of time for subsequent jobs to be scheduled.

Greedy Algorithm Pseudocode:

1. Initialization:

- We start by creating two lists, “*M1_jobs*” and “*M2_jobs*”, which contain jobs assigned to machines M1 and M2, respectively.
- To keep track of the end times of the last job scheduled on each machine, we define two variables, “*M1_end*” and “*M2_end*”, which are both initialized to zero.
- We also define a variable called “*total_jobs*” to keep track of how many jobs are scheduled.

2. Core loop:

- The algorithm starts by sorting the list of jobs by end time in ascending order.
- For each job (*job_id*, *start_time*, *end_time*) on the sorted job list:
 - If the job's “*start_time*” is greater than or equal to “*M1_end*”, assign it to machine M1 and add the job ID to “*M1_jobs*”.
 - Update “*M1_end*” to “*end_time*”.
 - Increase the “*total_jobs*” count.
- Otherwise, if the job's “*start_time*” is greater than or equal to “*M2_end*”, assign it to machine M2.
 - add the “*job_id*” to “*M2_jobs*”.
 - Set “*M2_end*” to “*end_time*”.
 - Increase the “*total_jobs*” count.

3. Output:

- The loop repeats until all jobs have been processed.
- After all jobs are scheduled, the algorithm writes the following information into an output file.
 - Total number of jobs scheduled (total_jobs).
 - The list of jobs assigned to M1 (M1_jobs).
 - The list of jobs assigned to M2 (M2_jobs).

Complexity Analysis:

- **Time Complexity:**
 - Job Loading Complexity: Reading n jobs from the input file takes $O(n)$ time.
 - Sorting Jobs: Sorting the jobs by their end times takes $O(n \log n)$.
 - Scheduling Jobs: Iterating through the sorted jobs and scheduling them takes $O(n)$ time.
 - Overall Complexity: $O(n \log n)$
 - The dominant step is the sorting of the jobs, so the overall time complexity of the algorithm is $O(n \log n)$.