

**Department of Computer Science and Engineering**

**FACULTY OF ENGINEERING AND TECHNOLOGY  
UNIVERSITY OF LUCKNOW  
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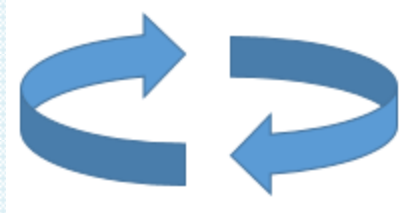


**CS-501**

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# SWAPPING AND MEMORY ALLOCATION

# SWAPPING



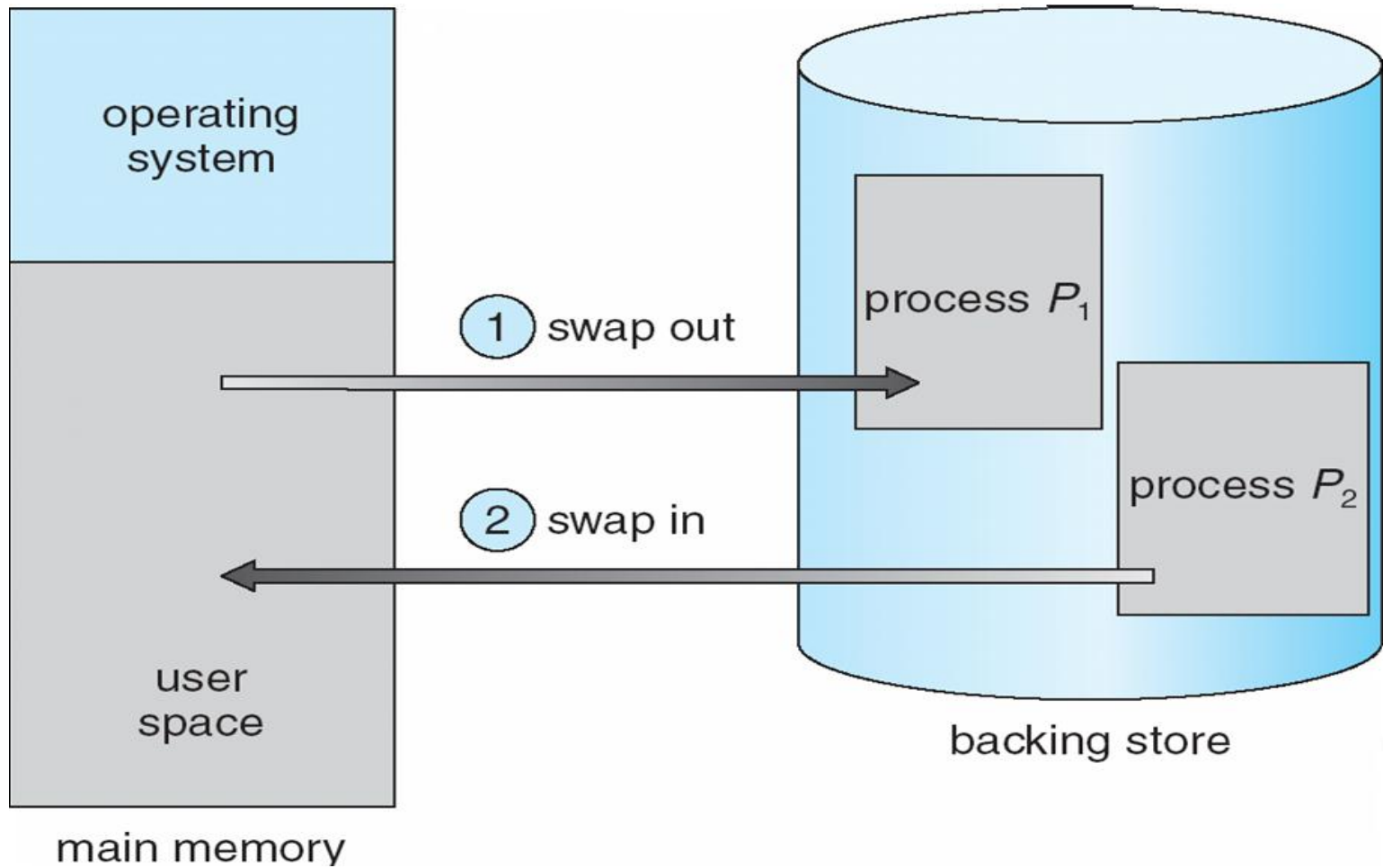
# Swapping

- A *process* can be swapped temporarily out of memory to a backing store, and then brought back into memory for continued execution.
- **Backing store**
  - Disk large enough to *accommodate* copies of all memory images for all users.
  - Must provide direct access to these memory images.
- **Roll out, roll in**
  - *swapping* variant used for priority-based scheduling algorithms.
  - Lower-priority process is swapped out so *higher-priority* process can be loaded and executed.

# Swapping

- Major part of swap time is *transfer time*.
- Total transfer time is directly proportional to the amount of *memory* swapped.
- Modified versions of *swapping* are found on many systems (i.e., UNIX, Linux, and Windows).
- System maintains a *ready queue* of ready-to-run processes which have memory images on disk.

# Swapping



# MEMORY ALLOCATION

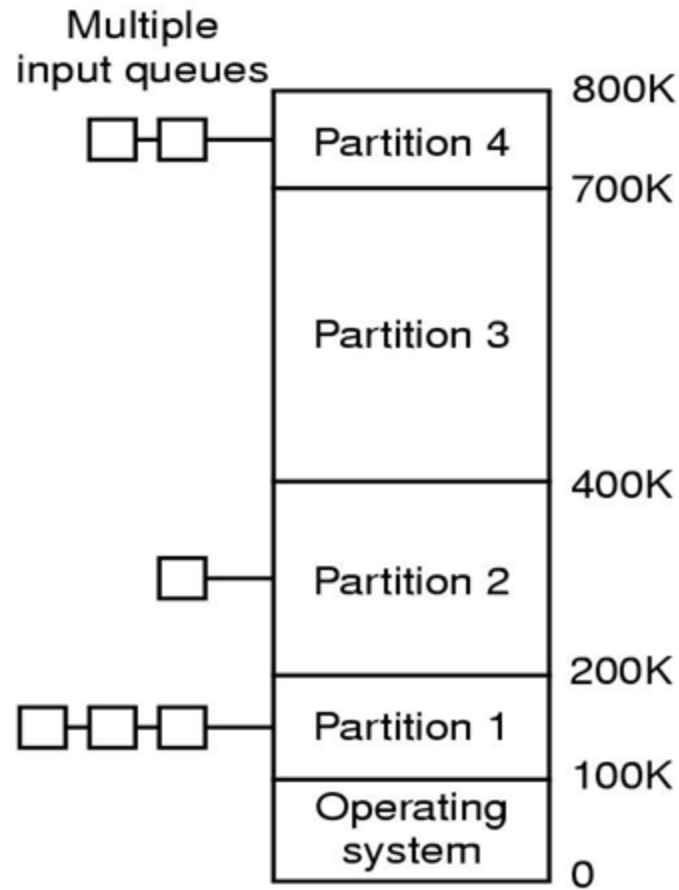
# Memory Allocation

- **Problem:**
  - How to allocate memory for multiple *processes* (in a multi-programming environment).
- **Solutions:**
- **Contiguous allocation**
  - Fixed partitions
  - Dynamic partitions
- **Paging**

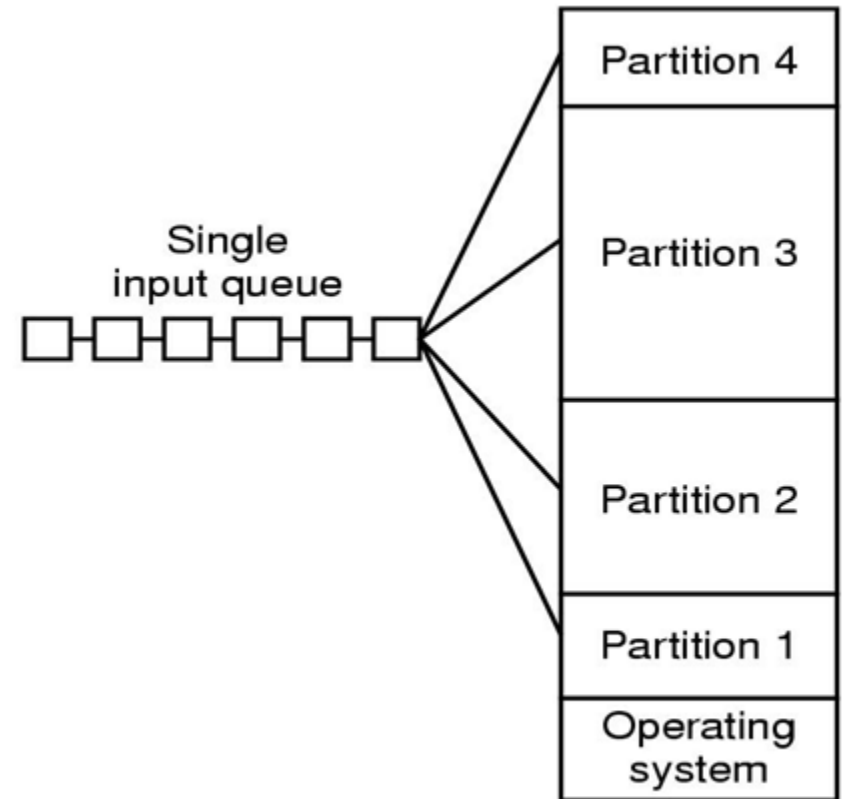


# Contiguous Allocation (Fixed Partitions)

# Contiguous Allocation (Fixed Partitions)



(a)



(b)

# Contiguous Allocation (Dynamic Partitions)

# Contiguous Allocation (Dynamic Partitions):

- *Main memory* usually have two partitions:
  - Resident operating system, usually held in *low memory* with interrupt vector.
  - User processes, held in *high memory*.

# Contiguous Allocation (Dynamic Partitions):

- Multiple-partition allocation

- **Hole**

- Block of available *memory*.
    - Holes of various size are *scattered* throughout memory.

- When a *process* arrives, it is allocated memory from a large enough hole.

- Operating system maintains information about:

- a) *Allocated* partitions

- b) *Free* partitions (hole).

# Dynamic Storage-Allocation Problems

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- **Problem**

- How to satisfy a request of size  $n$  from a list of free holes.

# Dynamic Storage-Allocation Problems

- **Solution:**

- *First-fit:* Allocate the first hole that is big enough.

- *Best-fit:* Allocate the smallest hole that is big enough; must search entire list, unless ordered by size

- Produces the smallest leftover hole

- *Worst-fit:* Allocate the largest hole; must also search entire list

- Produces the largest leftover hole

- **Note:** First-fit and best-fit better than worst-fit in terms of *speed and storage utilization*.



# Question-1

- Given five memory partitions of 100 KB, 500 KB, 200 KB, 300 KB, and 600 KB (in order), how would the first-fit, best-fit, and worst-fit algorithms place processes of size 212 KB, 417 KB, 112 KB, and 426 KB (in order)? Rank the algorithms in terms of how efficiently they use memory.

# Question-2

- Given six memory partitions of 300 KB, 600 KB, 350 KB, 200 KB, 750 KB, and 125 KB (in order), how would the first-fit, best-fit, and worst-fit algorithms place processes of size 115 KB, 500 KB, 358 KB, 200 KB, and 375 KB (in order)? Rank the algorithms in terms of how efficiently they use memory.

# References

1. Silberschatz, Galvin and Gagne, “Operating Systems Concepts”, Wiley.
2. William Stallings, “Operating Systems: Internals and Design Principles”, 6<sup>th</sup> Edition, Pearson Education.
3. D M Dhamdhere, “Operating Systems: A Concept based Approach”, 2<sup>nd</sup> Edition, TMH.

**Thank You.**

