Question 1. Complete the following table giving values of expressions in column 1 for the first 5 cycles. The values of integer flows for the first 5 cycles are given in rows, labelled P and Q.

Expression\Cycle	0	1	2	3	4	
Р	2	7	5	1	3	
Q	0	13	6	5	2	
pre(P)	nil	2	7	5	1	
Q+1	1	14	7	6	3	
P -> Q	2	13	6	5	2	
P -> (pre(Q)+1)	2	1	14	7	6	
P fby Q	2	0	13	6	5	

Question 2. Consider the following Lustre nodes

```
node foo(P: int) returns (Q, R: int)
var W, Y, Z: int;
let
        Q = W* Z;
        W = 0 -> (pre(Z) +1);
        Z = W + Y;
        Y = sqr(W);
        R = Q + W;
tel

node sqr(A :int) returns (B: int)
let
        B = A*A;
tel
```

Is this program causally correct? What is the order in which values of variables Q,R,W,Y,Z are calculated in each cycle?

Ans: **Yes**, the program is **casually correct**. The order in which values of variables Q,R,W,Y,Z are calculated in each cycle is as follows (the order is from **top to bottom**, i.e., W first and R last (or 5th here)):

```
W (1st)
Y (2nd)
Z (3rd)
Q (4th)
```

(5th)

R

Question 3.

(a) Define a Lustre node returning the following sequence of values. (Code required) 1, 4, 13, 40, 121, ...

(**Hint**: Consider how the difference of two successive terms grows.)

Note: In lustre, zero input nodes are not allowed. If your node is properly simulated in heptagon you can submit it. But keep the extension as .lus (Update Feb 12, 19.09 PM)

Ans: (In lustre, zero input nodes are not allowed. So I took a dummy value of type bool as input. It has no significance in the given program. Can be ignored. Compile & Run the node sequence to get the output).

Code:

```
node sequence(dummy: bool) returns (term: int)
var m : int;
let
  m = 1 -> (pre(m) * 3);
  term = 1 -> (pre(term) + m);
tel
```

(b) Complete the definition of following Lustre node (Code required)

```
node gen(req:bool) returns (ack: bool)
```

such that ack is true in the current cycle if and only if req has been true for the last 3 cycles (including the current cycle).

(Hint: First count for how many previous cycles req has been true continuously.)

Ans:

Code:

```
node gen(req:bool) returns (ack: bool)
var count: int;
let
   count = 0 -> if pre(req) then (pre(count) + 1) else 0;
   ack = if count >= 2 and req then 1 else 0;
tel
```

Question 4.

(a) Please describe in English the output produced by the following **Heptagon** node.

```
node t(x: bool^5) returns (y: bool);
let
    y = fold<<5>>(or)(x, false)
tel
```

Ans: The above Heptagon node t takes a boolean array x of 5 elements as input and produces a single boolean value y as output. y is false (or 0) if all the elements in x are false, otherwise, y is true or 1 (i.e. if at least 1 element of x is true then y is also true (or 1)). Here, actually:

```
y = false or x[0] or x[1] or x[2] or x[3] or x[4]
```

(b) What happens if you change the equation to y = fold <<5>> (or) (x, true)? Describe the output.

```
Ans: If we change the equation to y = fold<<5>>(or)(x, true), then the output y will always be true (or 1). Here, actually:
```

```
y = true \text{ or } x[0] \text{ or } x[1] \text{ or } x[2] \text{ or } x[3] \text{ or } x[4] = true \text{ (or 1)}
```

(c) Complete the definition of the following **Heptagon** node with parameter **n**. (Code required)

```
node mutex<<n>>(ack: bool^n) returns (ok: bool)
```

The node should check for mutual exclusion of ack[i] and ack[j]. That is **ok is true provided expression ack[i] and ack[j] is false for all i, j pairs with i not= j.**

Ans: (We will compile and run a node similar to display() to run this mutex node)

```
node f(ack: bool; initial: int) returns (count: int)
let
  count = if ack then initial + 1 else initial;
tel
node mutex<<n: int>>(ack: bool^n) returns (ok: bool)
var count: int;
let
  count = fold << n >> f(ack, 0);
 ok = if count < 2 then true else false;
tel
-- Code 1 used for Testing
node display() returns(ok: bool)
  ok = mutex<< 5 >>([false,false,false,false,true]);
tel
-- Code 2 used for Testing
--node display() returns(ok: bool)
-- ok = mutex<< 5 >>([false,false,true,false,true]);
--tel
```

Question 5. Study the following **Heptagon** code for ripple adder.

```
node mxor(x, y: bool) returns (c: bool)
let
    c = (x and not y) or ((not x) and y);
tel

node fa(x, y, cin: bool) returns (z, cout: bool)
let
    z = mxor(mxor(cin, x), y);
    cout = if cin then (x or y) else (x and y);
tel

node rippleadd<<n:int>>(a: bool^n; b: bool^n) returns (c: bool^n; over: bool)
let
    (c, over) = mapfold<<n>>fa(a, b, false);
tel
```

A sample code file with the above code simulating a 4-bit adder is given: <u>Download File</u>.

(a) Manually compute the output at 0th cycle of rippleadd for constants a=[0,1,1,0,1,1,1,0] and b=[1,1,0,1,0,1,1,1] given as input to rippleadd<<8>>(a,b).

```
Ans: 1
0
0
0
0
1
(This 1 is value of over (i.e. overflow) defined in rippleadd node)
```

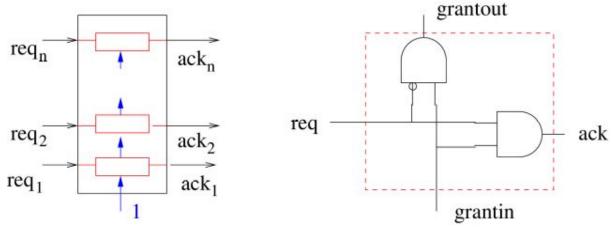
(b) Using this node rippleadd, define a node **counter** which counts in binary modulo 64. (It should **output unsigned 8 bit binary numbers corresponding to the decimal numbers** 0, 1, 2, ..., 63, 0, 1, 2,.... in successive cycles). **(Code required)**

Ans: (We will compile and run the node **counter()**. The node counter() also gives **overflow** as output which is **ignored**. The mod 64 values are represented by c7,c6,...,c0).

```
node mxor(x, y: bool) returns (c: bool)
let
    c = (x and not y) or ((not x) and y);
tel
```

```
node fa(x, y, cin: bool) returns (z, cout: bool)
let
  z = mxor(mxor(cin, x), y);
  cout = if cin then (x \text{ or } y) else (x \text{ and } y);
tel
node rippleadd<<n:int>>(a: bool^n; b: bool^n) returns (c: bool^n; over: bool)
  (c,over) = mapfold << n >> fa(a, b, false);
tel
node counter() returns (c0, c1, c2, c3, c4, c5, c6, c7: bool; overflow: bool)
var a, b, c: bool^6;
let
  a = [false, false, false, false, false];
  b = [true, false, false, false, false];
  (c,overflow) = (a, false) -> rippleadd<<6>>(pre(c), b);
  c0 = c[0]; c1 = c[1]; c2 = c[2]; c3 = c[3]; c4 = c[4]; c5 = c[5]; c6 = false; c7 = c[6]
false;
tel
```

Question 6. (Synchronous Bus Arbiter) An arbiter arbitrates between multiple requests coming at each cycle and gives acknowledgement to at most one of them. Consider the arbiter circuit below. **(Code required)**



Each **red box** (denoting a cell) in the left hand side figure is expanded to the **cell** circuit given in the right hand side figure.

```
Using Heptagon, model each cell as a
node cell(req, grantin: bool) returns (ack, grantout: bool)
```

Model a 5 cell arbiter as an assembly of 5 cells as shown in the figure. Call this node **arbiter**(req:bool^5) returns (ack:bool^5)

(**Hint**: See the implementation of rippleadd given in the previous question. Extra credit will be given if you can program this as an **n cell arbiter** with **parameter n**.)

Define a suitable **display node** to show the output. Simulate the arbiter using the tool **Heptagon** and check its functioning. Submit a screenshot of sample output using the sim2chro display.

Which of the following properties does this arbiter have?

- (a) Mutual exclusion of ack[i], ack[j] for i not= j.
- (b) No spurious ack, i.e. ack[i] => req[i]
- (c) No lost cycles, i.e. in any cycle, if there is at least 1 true request, the arbiter should have at least one true ack.

Ans : (The code given below is the implementation of **n bit arbiter.** For 5 bit arbiter, uncomment the 2 functions given in comments below and comment the last 2 functions. **Compile and run** the node **display** to get the **outputs**.)

This arbiter has all the 3 properties:

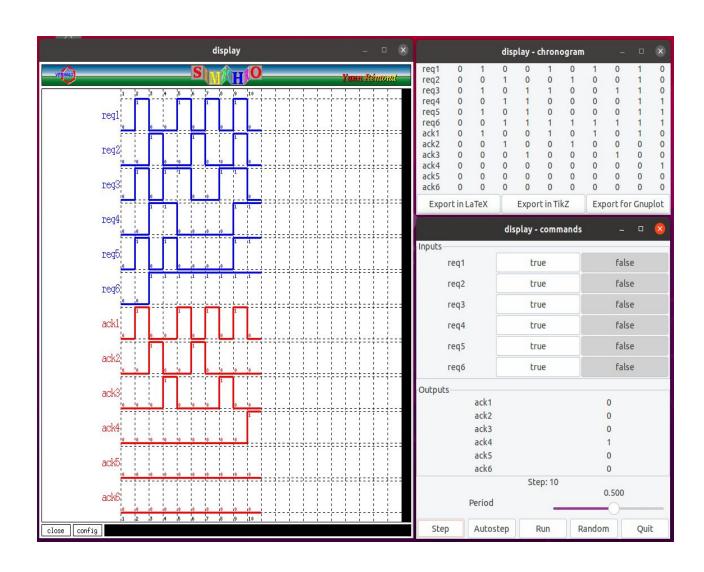
- a) Mutual Exclusion
- b) No spurious ack
- c) No lost cycles

```
node cell(req, grantin: bool) returns (ack, grantout: bool)
  ack = req and grantin;
  grantout = (not req) and grantin;
tel
--node arbiter(reg:bool^5) returns (ack:bool^5)
--var out: bool;
--let
-- (ack, out) = mapfold<<5>>cell(req, true);
--tel
--node display(req1, req2, req3, req4, req5: bool)
--returns (ack1, ack2, ack3, ack4, ack5: bool)
--var req, ack: bool^5;
--let
-- req = [req1, req2, req3, req4, req5];
-- ack = arbiter(req);
-- ack1 = ack[0]; ack2 = ack[1]; ack3 = ack[2]; ack4 = ack[3]; ack5 = ack[4];
--tel
```

```
node arbiter<<n:int>>(req:bool^n) returns (ack:bool^n)
var out: bool;
let
    (ack, out) = mapfold<<n>>cell(req, true);
tel

node display(req1, req2, req3, req4, req5, req6: bool)
returns (ack1, ack2, ack3, ack4, ack5, ack6: bool)
var req, ack: bool^6;
let
    req = [req1, req2, req3, req4, req5, req6];
    ack = arbiter<<6>>(req);
    ack1 = ack[0]; ack2 = ack[1]; ack3 = ack[2]; ack4 = ack[3]; ack5 = ack[4]; ack6 = ack[5];
tel
```

Screenshot of output:



Question 7. A monitor node for a property S takes as input a set of flows to observe (e.g. p, q: bool). It outputs a single boolean flow ok. The idea is that at every clock cycle, ok is true if the property S holds for the past sequence of inputs (including the current cycle). For example:

Property S: "p is continuously true in the past" has the monitor node

```
node smonitor(p: bool) returns (ok: bool)
let
    ok = p -> (pre(ok) and p);
tel
```

Answer the following

(a) Give the output of the above **smonitor** node for the input flow p = true true false true false true true

Ans: ok = true true true false false false false false (or 1 1 1 0 0 0 0 0)

(b) Give a monitor for the following property: "p is continuously true in the past AND q has occured at least once in the past." (Code required)

Ans:

Code:

```
node monitor(p, q: bool) returns (ok: bool)
var s, t: bool;
let
    s = q -> (pre(s) or q);
    t = p -> (pre(t) and p);
    ok = s and t;
tel
```

(c) Give a monitor node for the following property: Assume that a, b, c are boolean flows. "Everytime a occurs, c will remain continuously true from then on until a b occurs". Specify additional assumptions that you make in your design (E.g. what happens if a and b occur simultaneously). (Code required)

Ans:

Assumption: if a and b occurs simultaneously, i have given priority to b (i.e., c will be false in this case)

```
node monitor(a, b: bool) returns (c: bool)
let
    c = (a and (not b)) -> ((pre(c) or a) and (not b));
tel
```

(d) Give a monitor node for the following property: Assume that req and ack are boolean flows. "If req has been true for the last 3 cycles (including the current cycle) then ack must be true in the current cycle." (Code required)

Clarification: req and ack are both input boolean flows. ok is the output flow. ok is true if the property is observed/followed. (Update Feb 12, 18.00 PM)

Ans:

```
node monitor(req, ack:bool) returns (ok: bool)
var count: int;
let
    count = 0 -> if pre(req) then (pre(count) + 1) else 0;
    ok = (not (if count >= 2 and req then 1 else 0)) or ack;
tel
```