Unique Paths Problem - Solutions Notes

1. Recursive Solution

- Idea: Start from (0,0) and move only right or down until reaching (m-1,n-1). - Base Case: If out of bounds \rightarrow 0, if at destination \rightarrow 1. - Time Complexity: O(2^(m+n)), very high (exponential). - Space Complexity: O(m+n) recursion stack.

```
class Solution {
public:
    int countPaths(int i, int j, int m, int n) {
        if(i >= m || j >= n) return 0;
        if(i == m-1 && j == n-1) return 1;
        return countPaths(i+1, j, m, n) + countPaths(i, j+1, m, n);
    }
    int uniquePaths(int m, int n) {
        return countPaths(0, 0, m, n);
    }
};
```

2. Recursive + Memoization

- Optimization: Store results of subproblems in DP table. - Avoids recomputation of overlapping subproblems. - Time Complexity: O(m*n). - Space Complexity: O(m*n) + recursion stack.

```
class Solution {
public:
    int solve(int i, int j, int m, int n, vector<vector<int>>& dp) {
        if(i >= m || j >= n) return 0;
        if(i == m-1 && j == n-1) return 1;
        if(dp[i][j] != -1) return dp[i][j];
        return dp[i][j] = solve(i+1, j, m, n, dp) + solve(i, j+1, m, n, dp);
    }
    int uniquePaths(int m, int n) {
        vector<vector<int>> dp(m, vector<int>(n, -1));
        return solve(0, 0, m, n, dp);
    }
};
```

3. Tabulation (Bottom-Up DP)

- Build DP table iteratively. - dp[i][j] = dp[i-1][j] + dp[i][j-1]. - Base Case: dp[0][0] = 1. - Time Complexity: $O(m^*n)$. - Space Complexity: $O(m^*n)$.

4. Space Optimization

- Use only two rows (previous + current). - Reduces space from $O(m^*n) \to O(n)$. - Time Complexity: $O(m^*n)$. - Space Complexity: O(n).