

## ■ Subarray XOR Equal to K

Given an array of integers and an integer K, count the number of subarrays such that the XOR of all elements in the subarray is equal to K.

■ Intuition: - XOR behaves like addition in prefix-sum problems but with different properties. - If we know the XOR of prefix  $[0..i]$  and prefix  $[0..j]$ , then XOR of subarray  $[i+1..j]$  can be computed as:  $\text{prefixXOR}[j] \oplus \text{prefixXOR}[i]$  - We maintain a running XOR ( $\text{xr}$ ) as we traverse the array. - If  $\text{xr} == k$ , then the subarray from 0 to current index has XOR = k. - Otherwise, we check if there exists a prefix with XOR =  $(\text{xr} \oplus k)$ . If yes, then some subarray between that prefix and the current index has XOR = k. - We store the frequency of each prefix XOR in a hashmap (dictionary).

■ Algorithm: 1. Initialize count = 0,  $\text{xr} = 0$ , and a hashmap to store frequencies of prefix XORs. 2. Traverse the array: - Update  $\text{xr}$  with current element. - If  $\text{xr} == k$ , increment count. - Compute  $\text{rem} = \text{xr} \oplus k$ . - If  $\text{rem}$  exists in hashmap, add its frequency to count. - Store current  $\text{xr}$  in hashmap (increase its frequency). 3. Return count at the end.

### Python Code:

```
from collections import defaultdict

class Solution:
    def subarrayXor(self, arr, k):
        n = len(arr)
        count = 0
        xr = 0
        mpp = defaultdict(int)

        for i in range(n):
            xr ^= arr[i]

            if xr == k:
                count += 1

            rem = xr ^ k
            count += mpp.get(rem, 0)

            mpp[xr] += 1

        return count
```