August 27th Lectur 2 Model of Stochastic dynamical Systems (SDS) (A) General Model Xtx = & (Xt, Ut, Wt) -System State Signence 2) $Y_t = h_t(x_t, v_t) - Observatur$ Sequence (=0,1,2,... 7/4 (whole numbers) Ut - Central action Princtin random $Y_0, W_0, W_1, \ldots, V_o, V_1, \ldots$ variables (Their nature is not under our Central.) - Queveing enample - arrival sequence of job A_{0}, A_{1}, \ldots unit amount of work. Ai 6 80,13 W_0', W_1', W_2, \ldots Service Sequence at Servy 1 W_0^2 , W_1^2 , W_2^2 , ... System State - number in W; E {0,13

$$X_{tH}^{1} = max \left(X_{t}^{1} + 1_{S} u_{t} = Q_{1}, A_{t} = 1_{S}^{2} - W_{t}^{1}, 0 \right)$$

$$X_{t+2}^{2} = max \left(X_{t}^{2} + 1_{S} u_{t} = Q_{2}, A_{t} = 1_{S}^{2} - W_{t}^{2}, 0 \right)$$

$$Y_{t}^{1} = X_{t}^{1}, Y_{t}^{2} = X_{t}^{2} \quad (Perfect observations)$$

$$- Shace Craft example
$$X_{tH} = X_{t} + U_{t} + W_{t}$$

$$Y_{t} = CX_{t} + V_{t}$$$$

B) Classify Control Strategies / laws / Policies

1) OPEN-LOOP CONTROL

Fin Some Uo, U, ,... { Upg } 20 Consider system evolution

$$X_1 = f_0(X_0, V_0, W_0)$$

$$X_2 = \int_1 (X_1, U_1, \omega_1)$$

(for all)

$$= b_{1,0} (x_0, u_0, u_1, w_0, w_1)$$

 $Yt \quad Xt = \int_{E_{1,0}} (X_{0}, U_{0}, U_{1}, ..., U_{E_{1}}, W_{0}, W_{1}, ..., W_{E_{1}})$

Furthermon, Starting from time m & f, we have + f>m $X_{t} = b_{t-1, m} (X_{m}, U_{m_{1}, ..., U_{t-1}}, W_{m_{1}, ..., U_{t-1}})$ Similarly, Yo = ho (xorvo) $Y_1 = h_1(x_1, v_1)$ ν, (η, (κο, μο, ν,),ν,) = h1,0 (X0, W0, W0, V1) Different. Yt = ht,0 (Xo, Uo,..., Ut, wo,..., WH, Vt) What about the Corresponding deterministic system? (a) No & Wt3+70 & {Vt3+30 (6) or Set to constante w.p. 1. $x_{t+1} = \phi_t(x_t, u_t) \qquad (Y_t = h_t(x_t))$ of Eury + > 0 is the Central Sequence Xo is given Fining {Ur3+>,0, x+= b+1,0(x0,40,...,4+1) Feedbach policies $g = (g_0, g_1, g_2, \dots) = g_1 + g_1$

=>
$$U_{t} = g_{t}(x_{0:t}, u_{0:t-1})$$

 x_{0} is given
 $x_{t} = \int_{t-1}^{\infty} (x_{0}, g_{0}, \dots, g_{t-1})$

Claim: For deterministic Systems (as sherified about) feedbach Centrol Cases and open- Coop centrol Cases are equivalent.

Proof: (=>) Consider an open loop Control Caw \hat{U}_0 , \hat{U}_1 , ...

Define $g = (g_0, g_1, \dots)$ such that (s.f.) $g_+(x_{0:t}, u_{0:t}) = \hat{u}_t$ f(s.f.)(Constant)

= $\chi^{9}_{t} = b_{t-1,0} (\chi_{0}, g_{0}, \dots, g_{t-1})$

 $= b_{t-1,0} \left(x_0, \hat{u}_0, \dots, \hat{u}_{t-1} \right) = x_t^{OL}.$

(E) Consider om arbitrary feedback control bus $\hat{\mathcal{G}} = (\hat{\mathcal{G}}_0, \dots, \hat{\mathcal{G}}_{\ell}, \dots).$

 $\hat{\mathcal{U}}_0 = \hat{\mathcal{Y}}_0(x_0), \quad \hat{\mathcal{U}}_1 = \hat{\mathcal{Y}}_1(x_0, x_1, \hat{\mathcal{U}}_0), \dots$

$$\hat{\mathcal{U}}_{t} = \hat{g}_{t}(x_{0}, \dots, x_{t}, \hat{\mathcal{U}}_{0}, \dots, \hat{\mathcal{U}}_{t-1}), \dots$$
an known at $t=0$! (Can be computed)

Why?

$$\hat{\mathcal{U}}_{0} = \hat{g}_{0}(x_{0})$$

$$\Rightarrow x_{1} = b_{1}(x_{0}, \hat{\mathcal{U}}_{0}) = b_{0,0}(x_{0})$$

$$\hat{\mathcal{U}}_{1} = \hat{g}_{1}(x_{0}, \hat{\mathcal{U}}_{0})$$

$$= \hat{g}_{1,0}(x_{0})$$

$$\vdots$$
Same logic holds for the whole Seguence.

Choose Open-Coop (ential (aw $\hat{\mathcal{U}}_{0}, \hat{\mathcal{U}}_{1}, \dots$

Thun, X DL = X g + + !

(5) SDS - Open-wood versus feedback central Cours. Feedback control laws can gin a richer Sit of System trajectnies.

linear 1-D with perfect observations Enample Xtt = Xt + Ut +Wt Consider 9 2 (90,9,...) S.t.

Then
$$X_{t+}^9 = 9_t(x_0;t, u_0;t+1)$$

$$= 9_t(x_t) = -x_t$$
Then $X_{t+}^9 = x_t + u_t + w_t$

$$= x_t + 9_t(x_t) + w_t$$

$$= w_t + t$$

$$= w_t +$$

mutually independent. $\mathcal{N}(0, Q+tR)$ $X_{t}^{2} = W_{t-1} \sim \mathcal{N}(0, R)$ $x^{9} = x_{0} \sim N(0, \alpha)$ We cannot match! SDS are different from deterministic System. 6 Other key difference Remember, +t>, m $X_t = b_{t-1, m} (X_m, U_m, ..., U_{t-1})$ Knowing the State at time in Suffices to determine fection evolution. Shewfically, XmH = 6 m (xm, um) Knowledge of Xo, ...; Xm-1, Uo, ..., Ums

gins no additional advantage. No true for a SDS. Question: (LHS) 1P(x9+A) x0,..., xt anditional enfectation uo, ..., ur) (RHS) A ennt set. LHS IP (X ST FA / X+, U+) = IP(b+ (x+, u+, w+) & A | x+, u+) = IP($\S \omega: b_t(x_t, u_t, \omega) \in A_s^2 | u_t, u_t$) RHS IP (XIM GA / No:+, Uo:+) = IP (bt (x't, ut, wt) EA | No,..., Nt, uo,..., ut) Xt is a function of $Y_0, \omega_0, \ldots, \omega_{t-1}, v_0, \ldots, v_{t-1}$

 U_{+}^{9} is a function of X_{6} , W_{0} , ..., W_{E-1} , V_{0} , ..., V_{E-1} , V_{+}

If there is defendence of We on any of these vandom variables, Chin LHS may not equal KHS. However, it wis independent of κο, Wo, ..., Wt-1, Vo, ..., Ut-1, Vt, thun LHS= RHS!

lore assumption

Al. Ko, {W+3+>,0, {V+3+>,0 an mutually in dependent

Under A1

1P(xot CA | xo:t, uo:t) = IP (xoff EA \ Xf, Uf) = IP(of (xt, ut, wt) FA (xt, ut) Nota function of 3!

Policy independence!

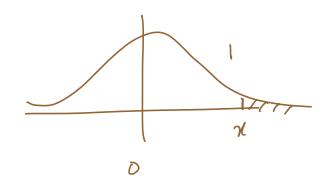
PE (XI'n -, ME) = X++ n++M+

 $W_t \sim \mathcal{N}(0,R)$

A = 2 x > 03 IP(b+(x+, u+, w+) GA | x+, u+) and Al was true

Wt~N(OLK) Q(x)= 1P(N(011)>, n)

$$= Q \left(- \frac{\chi_{F} - \chi_{F}}{\sqrt{E}} \right)$$



Revise Protability

Basic Probability Space (12, 7, 1P)

1 - Sample Shace

J - krunt Shace 6-algebra

Collection of Subsite of 1.

1. Non-empty - either 1 or \$ GI

2. Closed under complements

AGA, then $A^{c} = \Lambda \setminus AGJ$.

3. closed under Countable unions

 $A_1, A_2, \dots \in \Lambda$ VA; EN

lin U Ai = U Ai

$$\Lambda = \begin{cases} 1,2,\dots \end{cases} = IN \qquad J = 2^{\Lambda} = P(n)$$

$$\Lambda := \begin{cases} i \end{cases} \quad i = 1,2,\dots \qquad \text{(all subsite of } \Lambda)$$

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2. (Countable additivity)

$$A_{1,\ldots,A_{n,\ldots}} \in \mathcal{F}$$

f multially disjoint $i \neq j$ $A_i \cap A_j = \phi$ (multially exclusion)