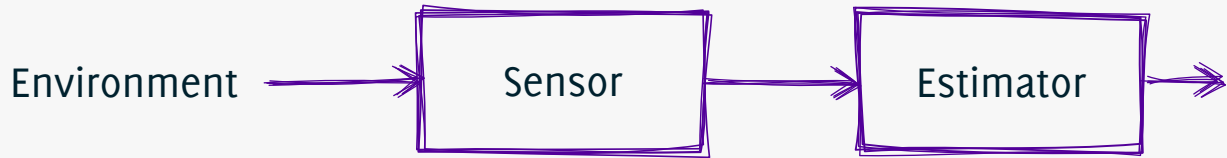


Estimation with active sensing

ADITYA MAHAJAN
MCGILL UNIVERSITY

February 11, 2011. ITA Workshop

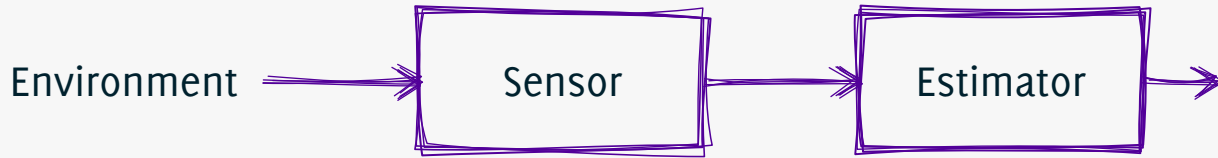
System Model



Sensor transmits packet to the estimator

- Each transmission incurs a constant cost
(energy required to switch on the radio and transmit a packet).
- Trade-off estimation quality with transmission energy
- The sensor decides when and how to transmit
(as opposed to the estimator scheduling the transmissions
... or the sensor communicating without coding)

System Model



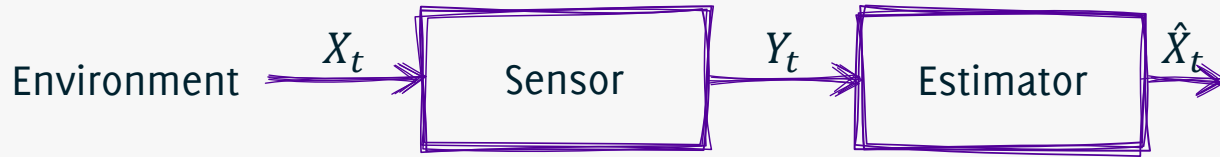
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When to communicate? And how to communicate?



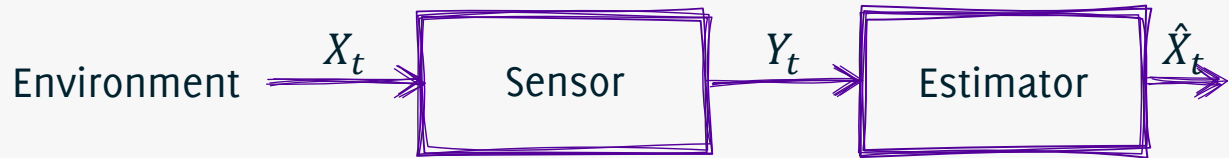
System Model



■ Environment

First order Markov process $\{X_t, t = 1, 2, \dots\}$, $X_t \in \mathbb{X}$

System Model



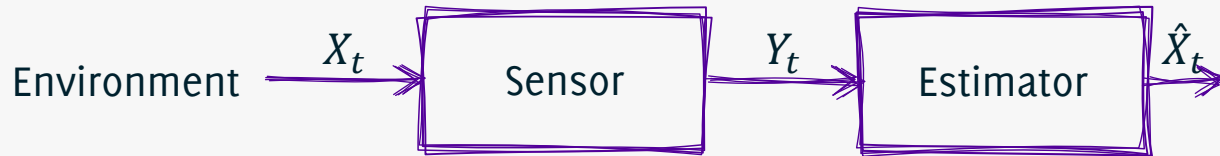
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$$Y_t = f_t(X_{1:t}, Y_{1:t-1}), \quad Y_t \in \mathbb{Y} := \mathbb{X} \cup \{\mathbf{b}\}$$

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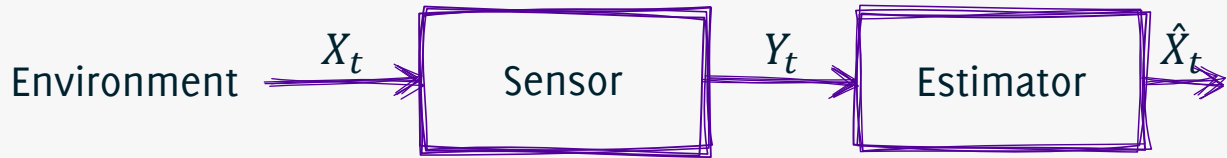
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$$\hat{X}_t = g_t(Y_{1:t}), \quad \hat{X}_t \in \mathbb{X}$$

■ Cost $c(Y_t) + d(X_t, \hat{X}_t)$

$$\triangleright c(y) = \begin{cases} 0, & y = \mathbb{b}, \\ c^*, & y \in \mathbb{X} \end{cases}$$

$\triangleright d$ is a metric on \mathbb{X}



Problem Formulation

- Transmission policy and estimation policies

$$\mathbf{f} = (f_1, f_2, \dots, f_T) \quad \mathbf{g} = (g_1, g_2, \dots, g_T)$$

- Performance of a policy

$$\mathcal{J}(\mathbf{f}, \mathbf{g}) := \mathbb{E}^{\mathbf{f}, \mathbf{g}} \left[\sum_{t=1}^T c(Y_t) + d(X_t, \hat{X}_t) \right]$$

- Causal Policies $(\mathcal{F}, \mathcal{G})$

$$\mathcal{F} = \prod_{t=1}^T F_t, \quad F_t = (\mathbb{X}^t \times \mathbb{Y}^{t-1} \mapsto \mathbb{Y})$$

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P1: Choose $(\mathbf{f}, \mathbf{g}) \in (\mathcal{F}, \mathcal{G})$ to minimize $\mathcal{J}(\mathbf{f}, \mathbf{g})$



Salient Features

- Sequential dynamic team

Sensor and estimator have access to different information at different times.



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Sensor and estimator have access to different information at different times.

- Causal real-time communication

Equivalent to a real-time communication system with input cost



Literature overview

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- **Implication:** Optimal policy is easy to implement . . .

. . . but the intuition relies of the Markov process being Gaussian.



For non-Gaussian processes, is the optimal policy easy to implement?

Solution approach

- Identify irrelevant information at sensor

Fix estimator. Set up the problem at the sensor as a MDP

- Find structure of optimal transmission and estimation policy:
Coding does not improve performance

A sequence of interchange arguments

- Identify a dynamic programming decomposition

Dynamic team problem. Use the coordinator approach proposed by [Mahajan, Nayyar, Teneketzis, 2008]

- Simplify DP using structural result

- Implementation issues



Removing irrelevant information

- Simple Causal policies $(\mathcal{F}', \mathcal{G})$

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■ Proposition

There is no loss of optimality in restricting attention to simple causal policies.



Proof Outline

Fix an arbitrary estimation policy. Look at the problem of choosing the best transmission policy.

Define $Z_t = (X_t, Y_{1:t-1})$. We can show that

■ $\mathbb{P}(Z_{t+1}|Z_{1:t}, Y_{1:t}) = \mathbb{P}(Z_{t+1}|Z_t, Y_t)$

■ $\mathbb{E}[c(Y_t) + d(X_t, \hat{X}_t)|Z_{1:t}, Y_{1:t}] = \mathbb{E}[c(Y_t) + d(X_t, \hat{X}_t)|Z_t, Y_t]$



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Hence, $\{Z_t, t = 1, 2, \dots\}$ is a controlled Markov process. Thus, there is no loss of optimality in restricting attention to

$$Y_t = f_t(Z_t) = f_t(X_t, Y_{1:t-1})$$



Structure of optimal policies

■ Structured Policies $(\mathcal{F}^*, \mathcal{G}^*)$

$$\mathcal{F}^* = \prod_{t=1}^T F_t^*,$$

$$F_t^* = \left\{ f_t \in \textcolor{red}{F}_t' : \forall x_t \in \mathbb{X}, \forall y_{1:t-1} \in \mathbb{Y}, f_t(x_t, y_{1:t-1}) \in \{x_t, \mathbb{b}\} \right\}$$



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Theorem: There is no loss of optimality in restricting attention to structured policies.



Implication: Real-time coding
does not improve performance

Proof Outline

- Define $(\mathcal{F}^{(s)}, \mathcal{G}^{(s)})$, $s = 0, 1, \dots, T$ as follows

$$(F_t^{(s)}, G_t^{(s)}) = \begin{cases} (F'_t, G_t) & t < s \\ (F_t^*, G_t^*) & t \geq s \end{cases}$$

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- Lemma \Rightarrow Theorem.

$$(\mathcal{F}^{(T+1)}, \mathcal{G}^{(T+1)}) = (\mathcal{F}', \mathcal{G}), \quad (\mathcal{F}^{(0)}, \mathcal{G}^{(0)}) = (\mathcal{F}^*, \mathcal{G}^*)$$



So how do we find an optimal policy?

Dynamic programming decomposition

- Non-classical information structure

$$Y_t = f_t(X_t, Y_{1:t-1})$$

Sensor

$$\hat{X}_t = g_t(Y_{1:t-1}, Y_t)$$

Estimator



Dynamic programming decomposition

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- Belongs to the class of tractable non-classical information structures identified in [Mahajan, Nayyar, Teneketzis, 2008]
- Look at the problem from the p.o.v. of a **coordinator** that observes the **common data** between the sensor and the estimator.
- The **local data** does not increase with time.



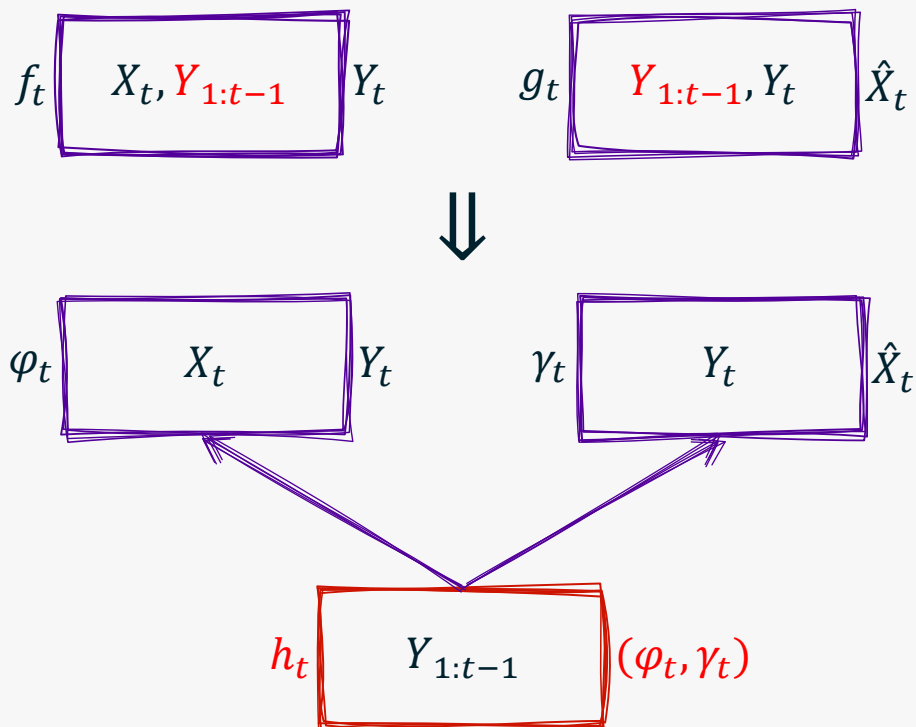
The coordinated system

$$f_t \left[X_t, Y_{1:t-1} \right] Y_t$$

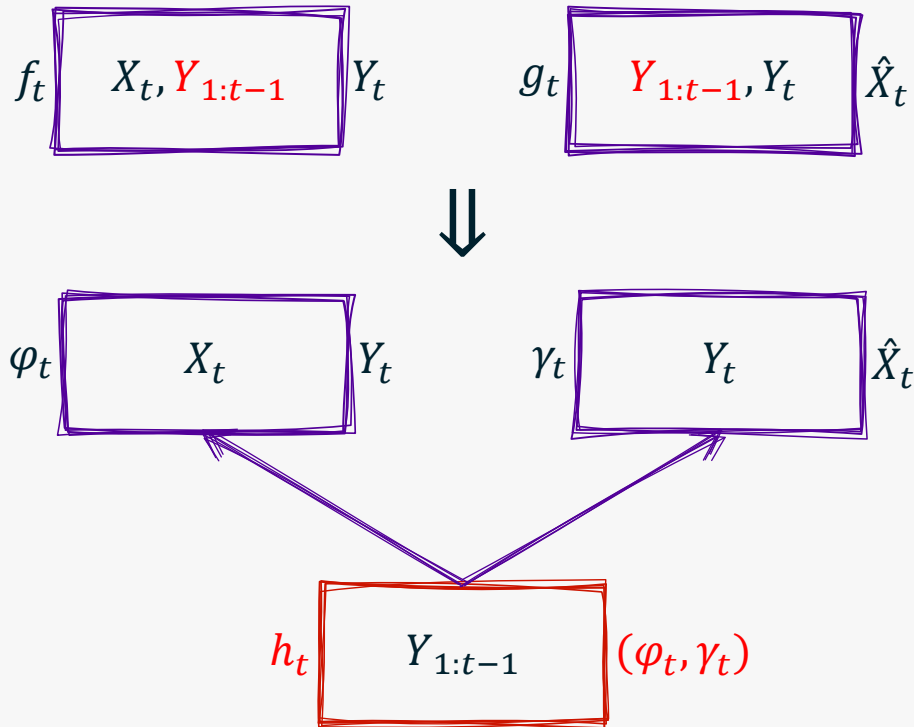
$$g_t \left[Y_{1:t-1}, Y_t \right] \hat{X}_t$$



The coordinated system



The coordinated system



Both systems are equivalent

The dynamic programming decomposition

The coordinated system is a centralized (single agent) POMDP



The dynamic programming decomposition

The coordinated system is a centralized (single agent) POMDP

■ Structure Result

Let $\pi_t = \mathbb{P}(X_t | Y_{1:t-1}, \varphi_{1:t-1}, \gamma_{1:t-1})$.

$$(\varphi_t, \gamma_t) = h_t(\pi_t)$$



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$$V_t(\pi_t) = \min_{\varphi_t, \gamma_t} \left\{ \mathbb{E} [c(X_t) + d(X_t, \hat{X}_t) + V_{t+1}(\text{update}(\pi_t, Y_t)) \mid \pi_t, \varphi_t, \gamma_t] \right\}$$



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■ Each step is a functional optimization problem



Refinement of DP

Restrict attention to structured policies



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■ $\forall f_t \in F_t^*, \quad f_t(x_t, y_{1:t-1}) \in \{x_t, \mathbb{b}\}.$

Thus, $\varphi_t(\cdot) \equiv \{x \in \mathbb{X} : \varphi_t(x_t) = x_t\} =: C_t$ (communication set)



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Thus, $\gamma_t(\cdot) \equiv \hat{x}_t^*$



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- Can be converted to countable state MDP.



Implementation issues

For a fixed transmission policy, $\pi_t \equiv (U_t, N_t)$ where

- ▶ U_t last non-blank transmission
- ▶ N time since last non-blank transmission



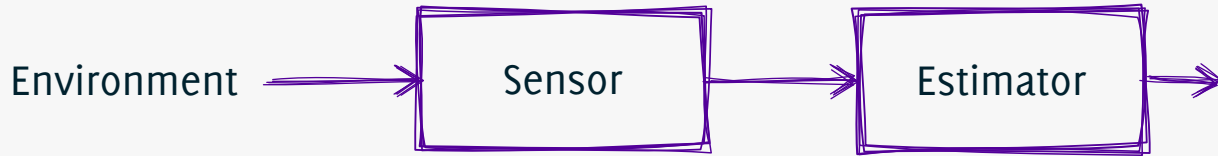
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Conclusion



- Use ideas from dynamic team theory and real-time communications
 - ▶ remove irrelevant data at sensor
 - ▶ coding does not improve performance
 - ▶ DP decomposition by investigating the coordinated system
 - ▶ Simplify DP using structure of optimal policy
- Transformed to denumerable MDP with set valued actions

$$V_t(\pi_t) = \min_{c_t, \hat{x}_t^*} \left\{ \mathbb{E} \left[c(X_t) + d(X_t, \hat{X}_t) + V_{t+1}(\text{update}(\pi_t, Y_t)) \mid \pi_t, \varphi_t, \gamma_t \right] \right\}$$



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Or, equivalently $\exists \text{index}_t : \mathbb{X} \mapsto \mathbb{R} \ni C_t = \{x \in \mathbb{X} : \text{index}_t(x) > \tau_t\}$



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- Why is this talk in the middle of two talks on bandits?

Consider “symmetric” Markov chains and a fixed decoding rule. Finding the optimal transmission policy is tangentially related to [Katehakis and Veinott, 1987] characterization of the Gittin’s index as a restart in i problem.

Could we reduce the problem to a 1-1/2 armed bandit problem?



Thank you.
Questions?

