

Fixed-finite delay decoding of i.i.d. sequences

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McGill University

Information Theory and Applications (ITA) Workshop

Motivation

© Block Markov coding schemes

Used to derive achievable schemes for multi-user info theory

- ▶ MAC with feedback
- ▶ Broadcast with feedback
- ▶ Relay

Share common features

- ▶ i.i.d. messages at each block
- ▶ decoding done with a fixed (block) delay

Motivation

© Proof techniques to derive structural results in teams

Usual way to prove structural properties:

- ▶ Person-by-person optimality
- ▶ Coordinator with common information

Both proceed by **backward induction**. Are there situations where backward induction does not work and one has to use a **forward induction** argument?

Motivation

© Proof techniques to derive structural results in teams

Our proof technique

- ▶ Genie-aided nodes that know additional information.
- ▶ Forward induction to **discard irrelevant variables**

What is the optimal structure (at block level) of block Markov coding schemes?

A new proof technique to prove structural results in teams.

Outline

⑤ Point-to-point communication

- ▶ Source coding
- ▶ Channel coding
- ▶ Channel coding with feedback

⑤ Two-transmitters one-receiver setup

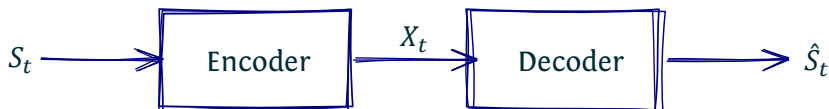
- ▶ ...

⑤ One-transmitter two-receivers

- ▶ ...

Point-to-point communication

Source Coding Model



i.i.d. source $\{S_1, S_2, \dots\}$

Sequential encoder $X_t = f_t(S_{1:t}, X_{1:t-1})$

Sequential decoder $\hat{S}_t = g_t(X_{1:t})$

Fixed-finite delay $\rho_t(S_{t-d}, \hat{S}_t)$

Objective: Choose $f = (f_1, \dots, f_T)$ and $g = (g_1, \dots, g_T)$ to minimize

$$J(f, g) = \mathbb{E} \left[\sum_{t=d+1}^T \rho_t(S_{t-d+1}, \hat{S}_t) \right]$$

Source Coding Model

© Literature Overview (delay = 0)

- ▶ i.i.d. source, delay = 0 (Ericson '79)

$$X_t = f_t(S_t)$$

- ▶ Markov source, delay = 0 (Witsenhausen '79, Walrand-Varaiya '82)

$$X_t = f_t(S_t, M_{t-1}) \quad \text{and} \quad X_t = f_t(S_t, \mathbb{P}(S_t | X_{1:t-1}))$$

Source Coding Model

© Literature Overview (delay > 0)

- ▶ Markov source, delay > 0 (Witsenhausen '79)

$$X_t = f_t(S_{t-d:t}, M_{t-1})$$

- ▶ i.i.d. source, delay > 0 (Asnani Wiessman '11)

$$X_t = f_t(S_{t-d:t}, \mathbb{P}(S_{t-d:t} \mid X_{1:t-1}))$$

Source Coding Model

© Literature Overview (delay > 0)

- ▶ Markov source, delay > 0 (Witsenhausen '79)

$$X_t = f_t(S_{t-d:t}, M_{t-1})$$

- ▶ i.i.d. source, delay > 0 (Asnani Wiessman '11)

$$X_t = f_t(S_{t-d:t}, \mathbb{P}(S_{t-d:t} \mid X_{1:t-1}))$$

© Our result

$$X_t = f_t(S_{t-d:t}, X_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t = g_t(X_{t-d:t})$$

A detour to MDPs

© Markov decision processes

$$X_{t+1} = f_t(X_t, U_t, W_t) \quad \text{and} \quad U_t = g_t(X_{1:t}, U_{1:t-1})$$

$$\min_g \mathbb{E}^g \left[\sum_{t=1}^T c_t(X_t, U_t) \right]$$

A detour to MDPs

© Markov decision processes

$$X_{t+1} = f_t(X_t, U_t, W_t) \quad \text{and} \quad U_t = g_t(X_{1:t}, U_{1:t-1})$$

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© Usual proof

A detour to MDPs

⊙ Markov decision processes

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$$\min_g \mathbb{E}^g \left[\sum_{t=1}^T c_t(X_t, U_t) \right]$$

⊙ Usual proof

Sequentially show that

$$U_T = g_T(X_T)$$

A detour to MDPs

⊙ Markov decision processes

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$$U_T = g_T(X_T)$$

$$U_{T-1} = g_{T-1}(X_{T-1})$$

A detour to MDPs

⊙ Markov decision processes

$$X_{t+1} = f_t(X_t, U_t, W_t) \quad \text{and} \quad U_t = g_t(X_{1:t}, U_{1:t-1})$$

$$\min_g \mathbb{E}^g \left[\sum_{t=1}^T c_t(X_t, U_t) \right]$$

⊙ Usual proof

Sequentially show that

$$U_T = g_T(X_T)$$

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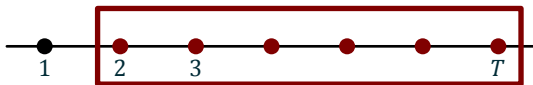
$$\dots = \dots$$

$$U_1 = g_1(X_1)$$

A detour to MDPs

© An alternative proof idea

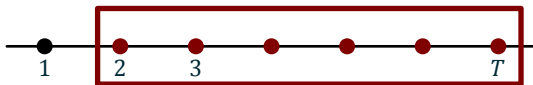
At time 2, consider a **coordinator** for all controllers from 2 up to T .



A detour to MDPs

⊙ An alternative proof idea

At time 2, consider a **coordinator** for all controllers from 2 up to T .



Common info (X_1, U_1, X_2)

Control action $\gamma^{(2)} = (\gamma_2^{(2)}, \gamma_3^{(2)}, \dots, \gamma_T^{(2)}) = h^{(2)}(X_1, U_1, X_2)$

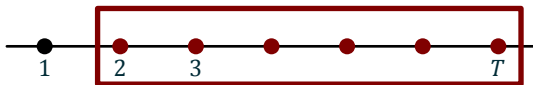
where $\gamma_t^{(2)}(\cdot) = g_t(X_{1:2}, U_1, \cdot) : X_{3:t} \mapsto U_t$

Cost at coordinator $\tilde{c}(X_2, \gamma^{(2)})$

A detour to MDPs

© An alternative proof idea

At time 2, consider a **coordinator** for all controllers from 2 up to T .



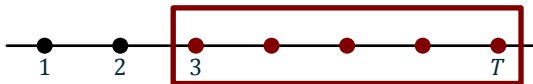
Without loss of optimality $\gamma^{(2)} = h^{(2)}(X_2)$, \Rightarrow

$$U_t = g_t(X_{2:t}, U_{2:t-1}), \quad \forall t > 2$$

A detour to MDPs

© An alternative proof idea

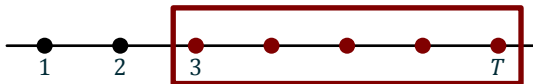
Repeat the argument for $t = 3$



A detour to MDPs

© An alternative proof idea

Repeat the argument for $t = 3$



Common info (X_2, U_2, X_3)

Control action $\gamma^{(3)} = (\gamma_3^{(3)}, \gamma_4^{(3)}, \dots, \gamma_T^{(3)}) = h^{(3)}(X_1, U_1, X_2)$

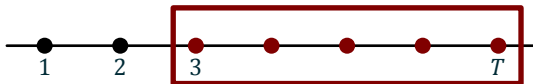
where $\gamma_t^{(3)}(\cdot) = g_t(X_{2:3}, U_2, \cdot) : X_{3:t}, U_{3:t-1} \mapsto U_3$

Cost at coordinator $\tilde{c}(X_3, \gamma^{(3)})$

A detour to MDPs

© An alternative proof idea

Repeat the argument for $t = 3$



Without loss of optimality $\gamma^{(3)} = h^{(3)}(X_3)$, \Rightarrow

$$U_t = g_t(X_{3:t}, U_{3:t-1}), \quad \forall t > 3$$

A detour to MDPs

© An alternative proof idea

Proceeding this way,

$$U_t = g_t(X_t), \quad \forall t$$

A detour to MDPs

© An alternative proof idea

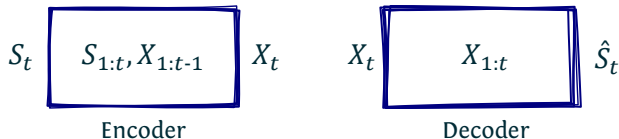
Proceeding this way,

$$U_t = g_t(X_t), \quad \forall t$$

Use this **forward induction** argument to prove the structural result for the source coding setup

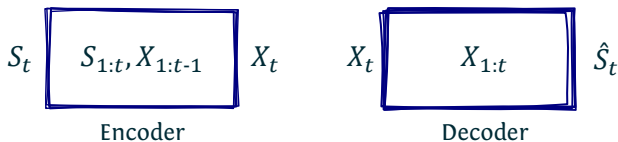
Structure of Optimal Encoder and Decoder

© Information structure of the system

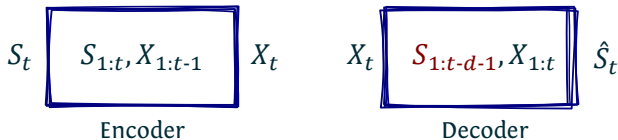


Structure of Optimal Encoder and Decoder

Information structure of the system

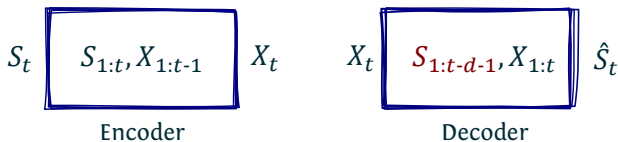


Genie aided decoder



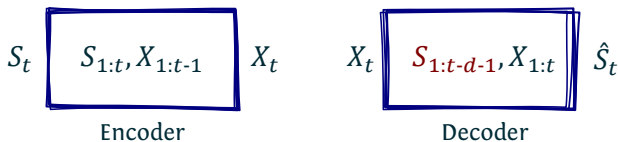
Structure of Optimal Encoder and Decoder

© Coordinator for $[d + 2, T]$



Structure of Optimal Encoder and Decoder

© Coordinator for $[d+2, T]$

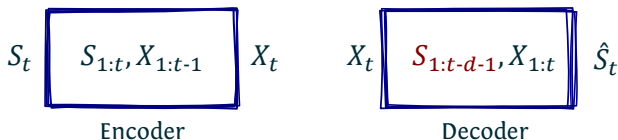


- ▶ Common info $(S_1, X_{1:d+1})$
- ▶ Control action: appropriate **partial functions** $\gamma^{(d+2)}$
- ▶ Cost at coordinator (depends on $\mathbb{P}(S_{2:T}, \hat{S}_{d+2:T})$)

$$\tilde{c}(X_{2:d+2}, \gamma^{(d+2)})$$

Structure of Optimal Encoder and Decoder

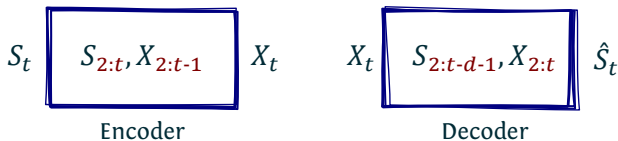
© Coordinator for $[d + 2, T]$



Without loss of optimality, $\gamma^{(d+2)} = h^{(d+2)}(X_{2:d+2})$.
All nodes in the coordinator may discard (S_1, X_1) .

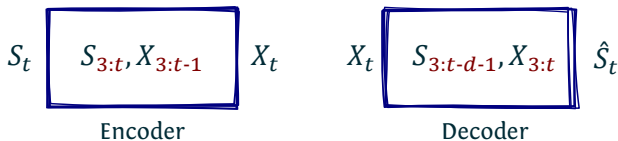
Structure of Optimal Encoder and Decoder

© Coordinator for $[d + 2, T]$



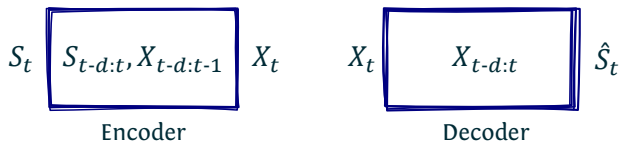
Structure of Optimal Encoder and Decoder

© Coordinator for $[d + 3, T]$



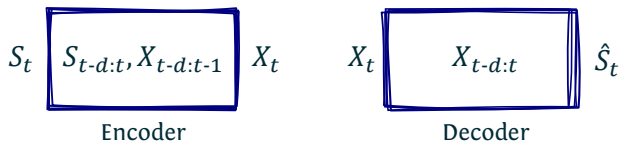
Structure of Optimal Encoder and Decoder

© Coordinator for T



Structure of Optimal Encoder and Decoder

© Coordinator for T



The decoder does not use the additional information $S_{1:t-d-1}$.

Channel coding model



i.i.d. source $\{S_1, S_2, \dots\}$

Sequential encoder $X_t = f_t(S_{1:t}, X_{1:t-1})$

Sequential decoder $\hat{S}_t = g_t(Y_{1:t})$

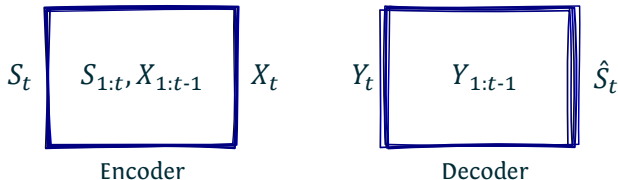
Fixed-finite delay $\rho_t(S_{t-d}, \hat{S}_t)$

Objective: Choose $f = (f_1, \dots, f_T)$ and $g = (g_1, \dots, g_T)$ to minimize

$$J(f, g) = \mathbb{E} \left[\sum_{t=d+1}^T \rho_t(S_{t-d+1}, \hat{S}_t) \right]$$

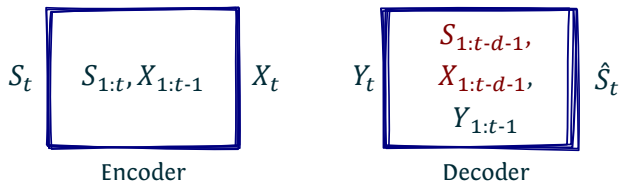
Structure of optimal encoder and decoder

© Information structure of the system



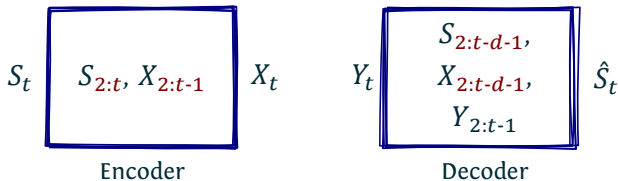
Structure of optimal encoder and decoder

© Genie aided decoder



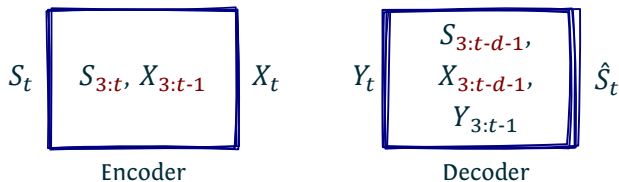
Structure of optimal encoder and decoder

© Coordinator for $[d + 2, T]$



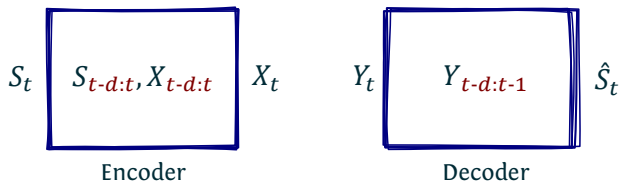
Structure of optimal encoder and decoder

© Coordinator for $[d + 3, T]$

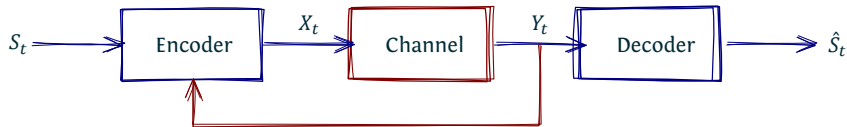


Structure of optimal encoder and decoder

© Coordinator for T



Channel coding with feedback



i.i.d. source $\{S_1, S_2, \dots\}$

Sequential encoder $X_t = f_t(S_{1:t}, X_{1:t-1}, Y_{1:t-1})$

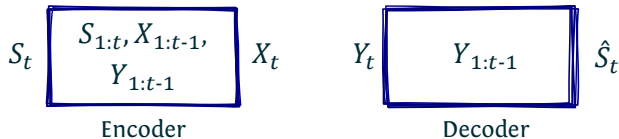
Sequential decoder $\hat{S}_t = g_t(Y_{1:t})$

Fixed-finite delay $\rho_t(S_{t-d}, \hat{S}_t)$

$$\min J(f, g) = \min \mathbb{E} \left[\sum_{t=d+1}^T \rho_t(S_{t-d+1}, \hat{S}_t) \right]$$

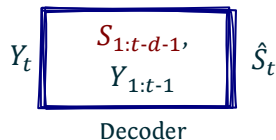
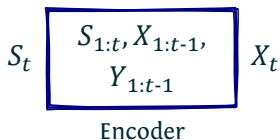
Structure of optimal encoder and decoder

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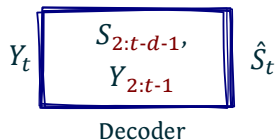
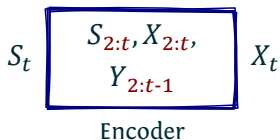
Structure of optimal encoder and decoder

© Genie aided decoder



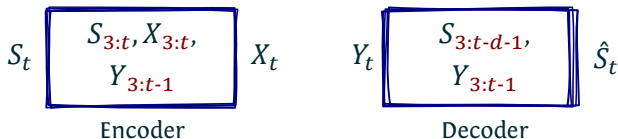
Structure of optimal encoder and decoder

© Coordinator for $[d + 2, T]$



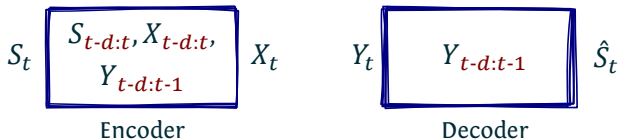
Structure of optimal encoder and decoder

⊙ Coordinator for $[d + 3, T]$



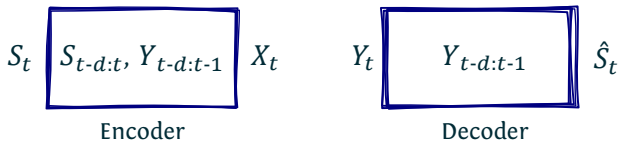
Structure of optimal encoder and decoder

© Coordinator for T



Structure of optimal encoder and decoder

© Fix decoder, consider optimal encoder



Given channel outputs and corresponding source outputs, channel inputs are irrelevant.

Summary of results for point-to-point setup

◎ Source coding

$$X_t = f_t(S_{t-d:t}, X_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t = g_t(X_{t-d:t})$$

◎ Channel coding

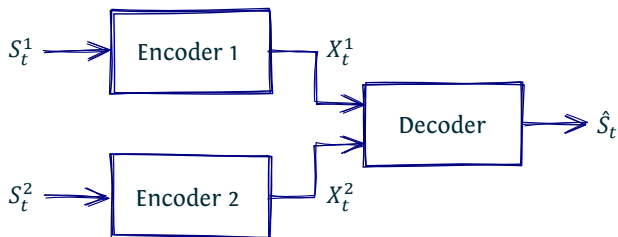
$$X_t = f_t(S_{t-d:t}, X_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t = g_t(Y_{t-d:t})$$

◎ Channel coding with feedback

$$X_t = f_t(S_{t-d:t}, Y_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t = g_t(Y_{t-d:t})$$

Multi-terminal communication
Two transmitters, one receiver

Distributed source coding



i.i.d. source $\{(S_1^1, S_1^2), (S_2^1, S_2^2), \dots\}$

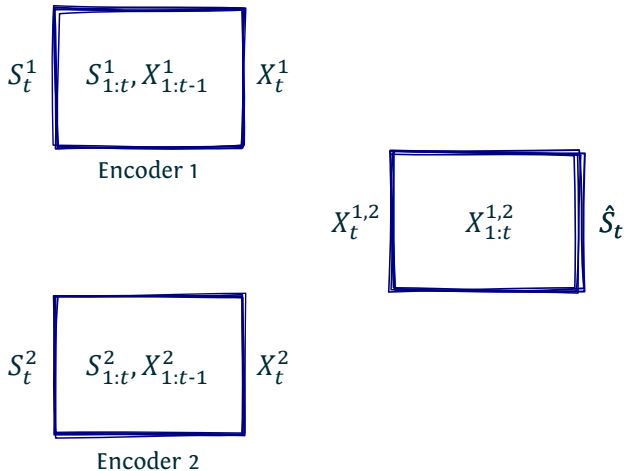
Sequential encoder $X_t^k = f_t(S_{1:t}^k, X_{1:t-1}^k)$

Sequential decoder $\hat{S}_t = g_t(X_{1:t}^1, X_{1:t}^2)$

Fixed-finite delay $\rho_t(S_{t-d}^1, S_{t-d}^2, \hat{S}_t)$

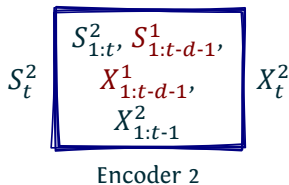
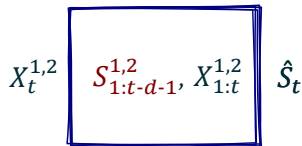
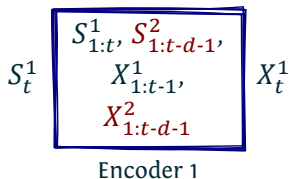
Structure of optimal encoders and decoder

© Original information structure



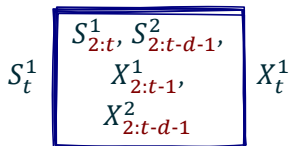
Structure of optimal encoders and decoder

Genie aided nodes

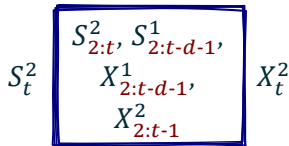
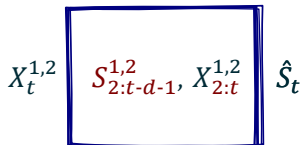


Structure of optimal encoders and decoder

⊙ Coordinator for $[d + 2, T]$



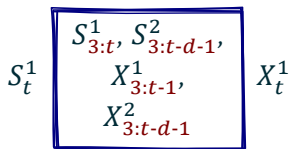
Encoder 1



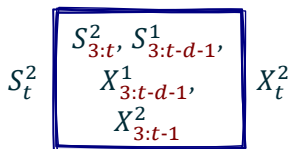
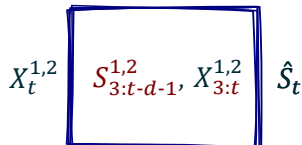
Encoder 2

Structure of optimal encoders and decoder

© Coordinator for $[d + 3, T]$



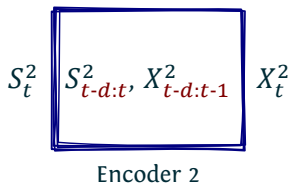
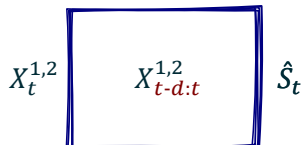
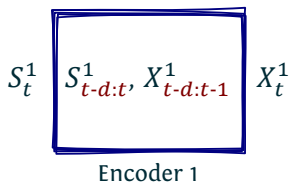
Encoder 1



Encoder 2

Structure of optimal encoders and decoder

Coordinator for T



Summary of results for 2-Tx 1-Rx

© Multi-terminal source coding

$$X_t^k = f_t^k(S_{t-d:t}^k, X_{t-d:t-1}^k) \quad \text{and} \quad \hat{S}_t = g_t(X_{t-d:t}^{1,2})$$

© Multiple access channel (MAC)

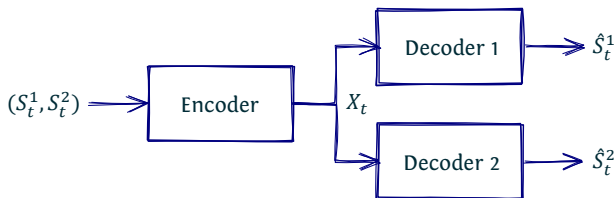
$$X_t^k = f_t^k(S_{t-d:t}^k, X_{t-d:t-1}^k) \quad \text{and} \quad \hat{S}_t = g_t(Y_{t-d:t})$$

© MAC with with feedback

$$X_t^k = f_t^k(S_{t-d:t}^k, X_{t-d:t-1}^k, Y_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t = g_t(Y_{t-d:t})$$

Multi-terminal communication
One transmitter, two receivers

Distributed reconstruction source coding



i.i.d. source $\{(S_1^1, S_1^2), (S_2^1, S_2^2), \dots\}$

Sequential encoder $X_t = f_t(S_{1:t}, X_{1:t-1})$

Sequential decoder $\hat{S}_t^k = g_t(X_{1:t})$

Fixed-finite delay $\rho_t(S_{t-d}^1, S_{t-d}^2, \hat{S}_t^1, \hat{S}_t^2)$

Summary of results for 1-Tx 2-Rx

⊙ Distributed reconstruction

$$X_t = f_t(S_{t-d:t}^{1,2}, X_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t^k = g_t^k(X_{t-d:t})$$

⊙ Broadcast channel

$$X_t = f_t(S_{t-d:t}^{1,2}, X_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t^k = g_t^k(Y_{t-d:t}^k)$$

⊙ Broadcast channel with feedback

$$X_t = f_t(S_{t-d:t}^{1,2}, Y_{t-d:t-1}) \quad \text{and} \quad \hat{S}_t^k = g_t^k(Y_{t-d:t}^k)$$

Concluding remarks

- © Optimal structure for fixed-finite delay reconstruction of i.i.d. sequences
 - ▶ Might be useful for block Markov coding schemes.
- © A (new?) proof technique to prove structural results
 - ▶ Genie-aided nodes with additional variables
 - ▶ Forward induction argument to remove irrelevant variables

Concluding remarks

© Dynamic programming decomposition

- ▶ All nodes have fixed memory. So we can use the approach of Sandell '74, M '08.

info state = $\mathbb{P}(\text{Memory of all nodes at } t)$

Time-varying optimal control laws!

- ▶ “State” observed after a delay. So we can also use the approach of Aicardi *et al.* '87, Nayyar M Teneketzis '12

info state = (local memory, *(partiallyevaluatedfunctions)*)

Time-invariant optimal control laws!

Thank you