

Aditya Mahajan Yale

Acknowledgements: Demos Teneketzis (UMich), Ashutosh Nayyar (UMich), Sekhar Tatikonda (Yale), Ramji Venkataramanan (Stanford) April 13, 2010, UC Berkeley

Stochastic control \Rightarrow Information theory

Converses in info theory

- ► Capacity of FSM channels Goldsmith-Varaiya-96
- ➤ Capacity of FSM channels with f/b Vishwanathan-99, Chen-Berger-05, Yang-Kavcic-Tatikonda-05, Tatikonda-Mitter-09, Kavcic-Mandic-Huang-Ma-09
- ► Error exponents of FSM channels with f/b Como-Yüksel-Tatikonda-09
- ► FS-MAC with partial state info Como-Yüksel-09

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Coding schemes

➤ Optimal coding schemes for channels with feedback
Hornstein-63, Schalkwijk-Kailath-66,
Shayevitz-Feder-09, Coleman-09,
Bae-Anastasopoulos-10, Gorantla-Coleman-10

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Coding schemes

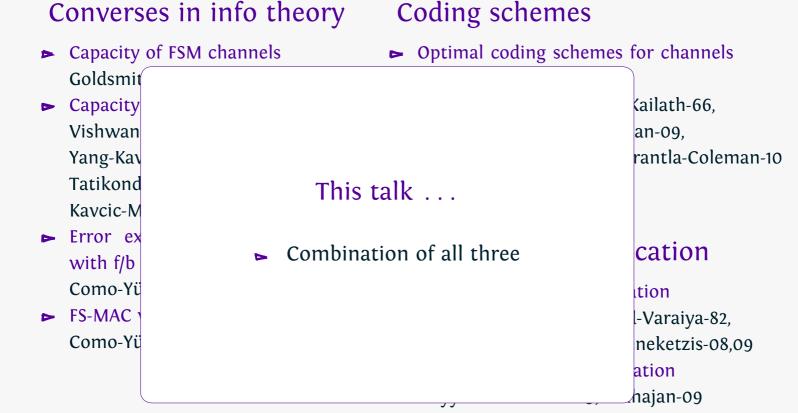
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Real-time communication

- ➤ Point-to-point communication
 Witsenhausen-79, Walrand-Varaiya-82,
 Teneketzis-06, Mahajan-Teneketzis-08,09
- ► Multi-terminal communication Nayyar-Teneketzis-09, Mahajan-09

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Stochastic control ⇒ **Information theory**



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This talk

Setup

- Discrete memoryless multiple access channel (MAC) with feedback
- ► Inner bounds on capacity region (achievable schemes) using block Markov superposition coding

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- Inner bounds on capacity region (achievable schemes) using block
 Markov superposition coding

Approach

- Setup block Markov superposition codes as decentralized control problem with delayed sharing of information
- ► A systematic study of auxiliary random variables

This talk ...

Setup

- Discrete memoryless multiple access channel (MAC) with feedback
- Inner bounds on capacity region (achievable schemes) using block
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- Setup block Markov superposition codes as decentralized control problem with delayed sharing of information
- ► A systematic study of auxiliary random variables

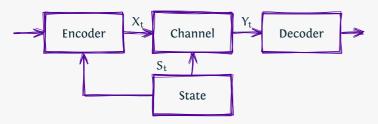
Main point

- ► Block Markov superposition coding schemes give rise to a specific information structure
- ➤ Stochastic control can provide insights into the design of coding schemes for a specific information structures

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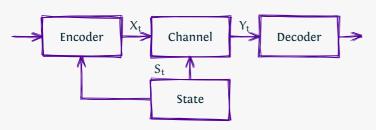
Are auxiliary random variables in info theory related to information states in stochastic control?

Channel coding with causal side information



[Shannon-58]

Channel coding with causal side information

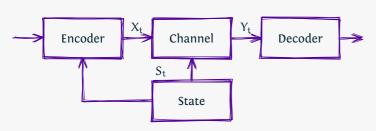


[Shannon-58]

Code functions

$$C = \max_{\substack{P_F \\ F: S \to \mathcal{X}}} I(F \land Y)$$

Channel coding with causal side information



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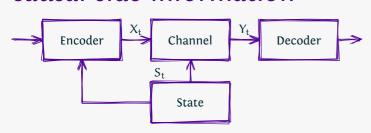
[Shannon-58]

Auxiliary random variables

$$C = \max_{P_U, f: \mathcal{U} \times \mathcal{S} \to \mathcal{X}} I(U \wedge Y)$$

where
$$|\mathcal{U}| \leq |\mathcal{Y}|$$

Channel coding with causal side information



Code functions

$$C = \max_{\substack{P_F \\ F: S \to \mathcal{X}}} I(F \land Y)$$

[Shannon-58]

Auxiliary random variables

We can think of the auxiliary random variable indexing the set of all functions from S to X.

$$C = \max_{P_U, f: \mathcal{U} \times \mathcal{S} \to \mathcal{X}} I(U \wedge Y)$$

where $|\mathcal{U}| \leq |\mathcal{Y}|$

[Keshet, Steinberg, Merhav, 2007]

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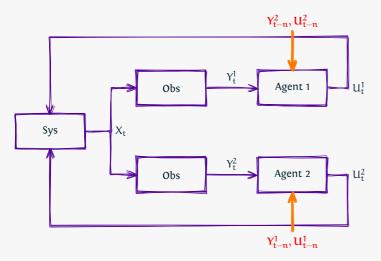
What are information states?

The (information) state should be a summary ('compression') of some data (the 'past') known to someone (an observer or a controller) and sufficient for some purposes (input-output map, optimization, dynamic programming).

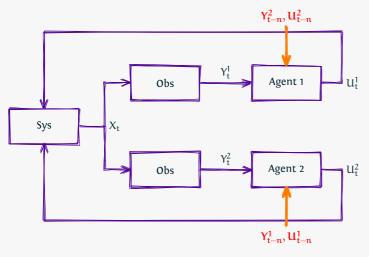
[Witsenhausen-1976]



Delayed sharing information structure



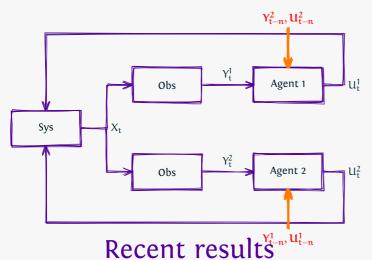
Delayed sharing information structure



History

- A bridge between classical (n = 0) and non-classical $(n = \infty)$ info structures
- ➤ Witsenhausen, 1971 proposed the n-DSIS and asserted a structure of optimal control policies
- ▶ Varaiya and Walrand, 1979 proved that Witsenhausen's assertion is true for n = 1 but false of n > 1

Delayed sharing information structure



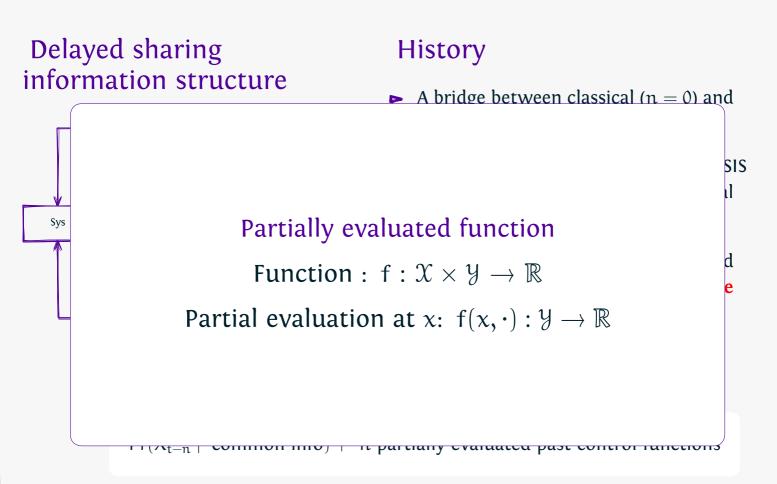
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Information state: [Nayyar-Mahajan-Teneketzis-10]

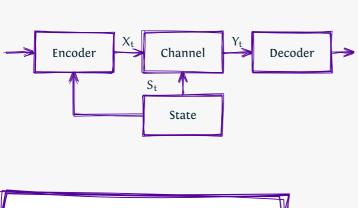
 $Pr(X_{t-n} \mid common info) + n-partially evaluated past control functions$





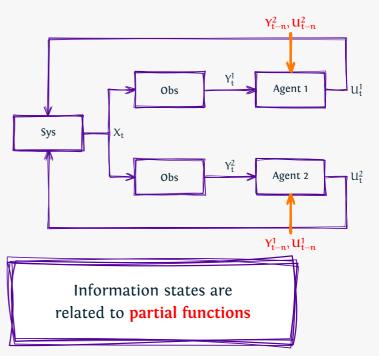
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Channel coding with causal side information



Auxiliary random variable is related to functions

Delayed sharing information structure

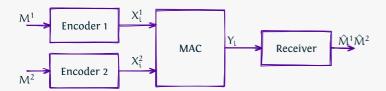




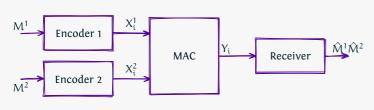
Outline

- 1. Overview of MAC with feedback
 - ▶ Cover-Leung scheme
 - Bross-Lapidoth scheme
 - Venkataramanan-Pradhan scheme
- 2. Formulation as a decentralized control problem with delayed sharing
 - ▶ delay = 1
 - ► delay = 2
- 3. Conclusion

Multiple access channel



Multiple access channel



Encoder

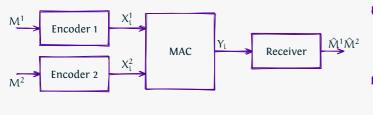
$$X_{[1:n]}^1 = f^1(M^1), \quad X_{[1:n]}^2 = f^2(M^2)$$

Channel

$$\mathbb{P}(Y_i \mid X_{[1:i]}^1 X_{[1:i]}^2) = \mathbb{P}(Y_i \mid X_i^1 X_i^2)$$

- ▶ Decoder: $\hat{M}^1\hat{M}^2 = d(Y_{[1:n]})$
- $ightharpoonup Error: \mathbb{P}(\hat{M}^1\hat{M}^2 \neq M^1M^2)$

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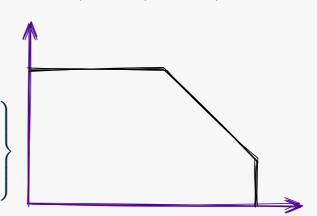
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Capacity

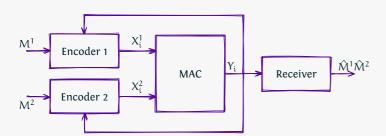
$$C = \bigcup_{P_{X^1}P_{X^2}} \left\{ (R_1, R_2) : \frac{R_1 \leqslant I(X^1 \wedge Y \mid X^2)}{R_2 \leqslant I(X^2 \wedge Y \mid X^1)} \right\}$$



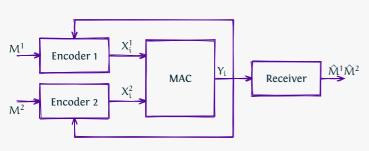
Alshwede-71, Liao-72

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MAC with feedback

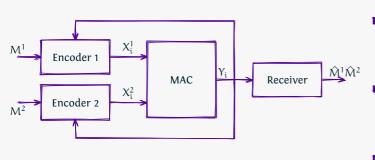


MAC with feedback



- ► Encoder $X_i^1 = f^1(M^1, Y_{[1:i]}), \quad X_i^2 = f^2(M^2, Y_{[1:i]})$
- Channel $\mathbb{P}(Y_i \mid X_{[1:i]}^1 X_{[1:i]}^2) = \mathbb{P}(Y_i \mid X_i^1 X_i^2)$
- ► Decoder: $\hat{M}^1\hat{M}^2 = d(Y_{[1:n]})$
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MAC with feedback



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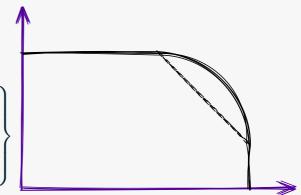
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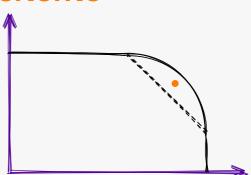
Capacity

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[Kramer-03]

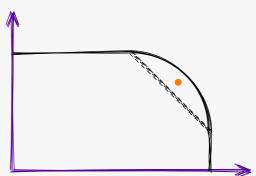
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Main idea

Suppose we want to transmit at a rate point outside the no-feedback capacity region. Communicate in two phases

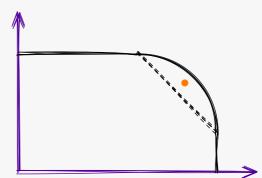
[Gaarder-Wolf-75]

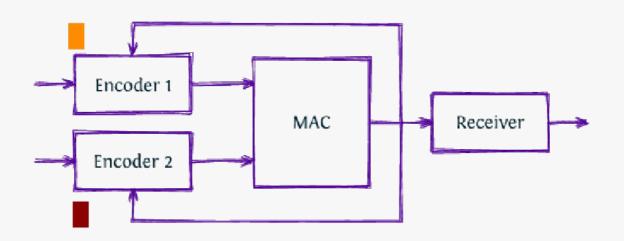


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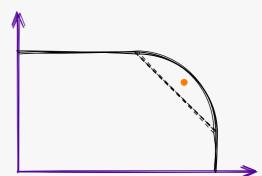


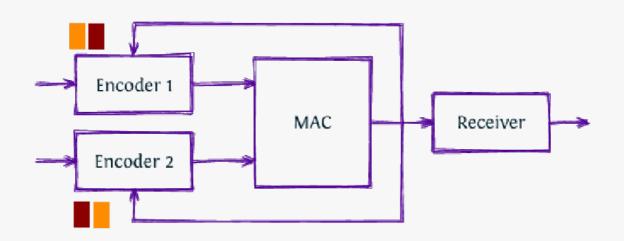


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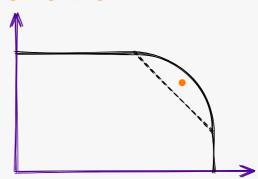


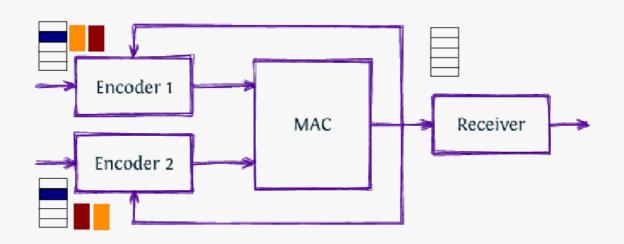


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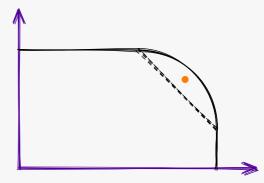


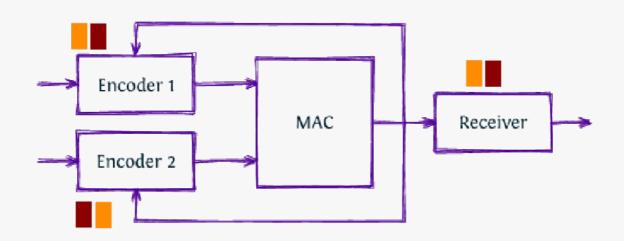


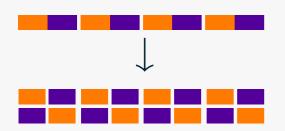
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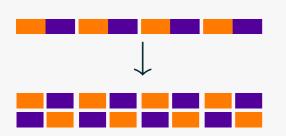
[Gaarder-Wolf-75]

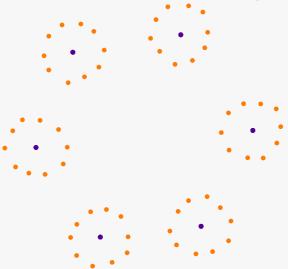




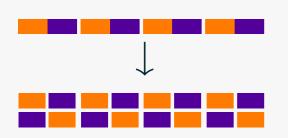


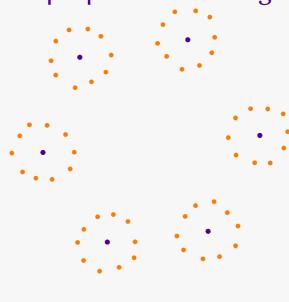
Superposition coding





Superposition coding



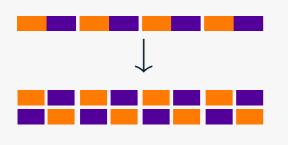


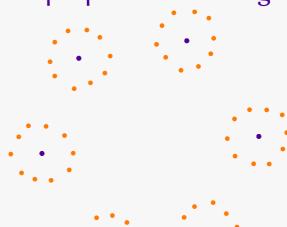
Binning



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Superposition coding





Cover-Leung region

$$R_{CL} = \bigcup_{P_{U}P_{X^{1}|U}P_{X^{2}|U}} \left\{ (R_{1},R_{2}) : \begin{matrix} R_{1} \leqslant I(X^{1} \wedge Y \mid UX^{2}) \\ R_{2} \leqslant I(X^{2} \wedge Y \mid UX^{1}) \\ R_{1} + R_{2} \leqslant I(X^{1}X^{2} \wedge Y) \end{matrix} \right\}$$

[Cover-Leung-81]

HH HH HH HH HH

Binning



Converse

Cover-Leung region

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[Cover-Leung-81]

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[Cover-Leung-81]

CL region is tight for . . .

- \rightarrow H(X¹|YX²) = 0 [Willems-82]
- ► Partial feedback [Willems-83]
- ► Binary erasure MAC [Willems-84]



Converse

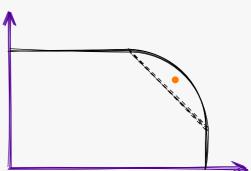
Cover-Leung region

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[Cover-Leung-81]

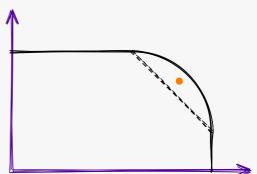
- CL region is tight for . . . CL region is **not** tight for . . .
- ► $H(X^1|YX^2) = 0$ [Willems-82]
- ► Gaussian MAC [Ozarow-84]
- Partial feedback [Willems-83]
- ► Poisson MAC [Bross-Lapidoth-05]
- Binary erasure MAC [Willems-84]
 - ▶ Others . . . [Kramer-02, Bross-Lapidoth-05, Venkataramanan-Pradhan-091

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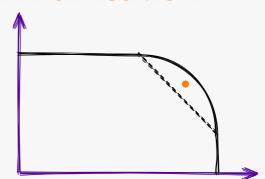
Main idea

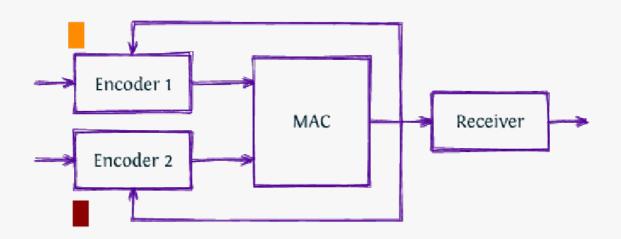
Each block now consists of three phases [Bross-Lapidoth-05]



Main idea

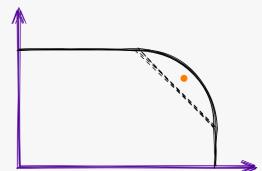
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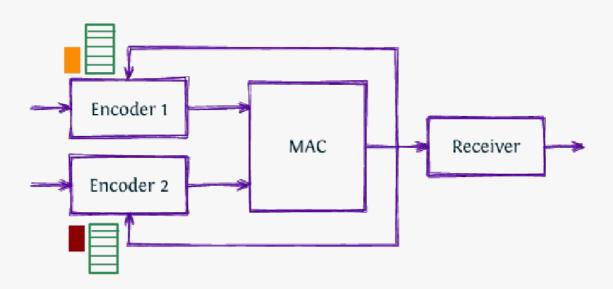




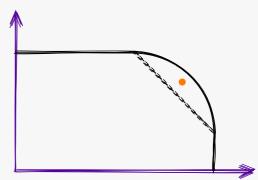
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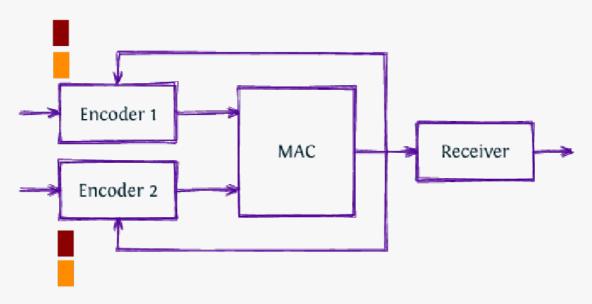
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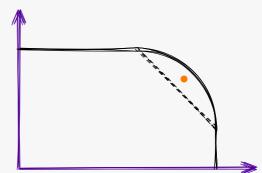
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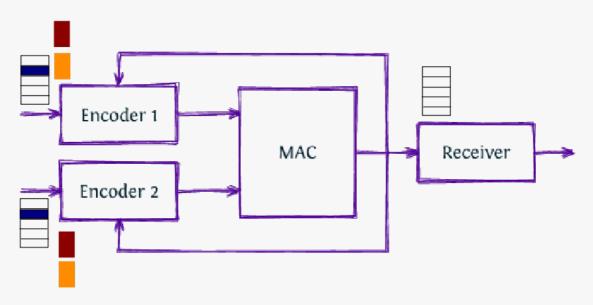




Main idea

Each block now consists of three phases
[Bross-Lapidoth-05]

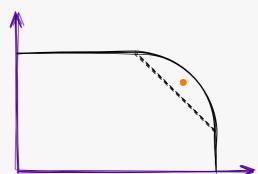


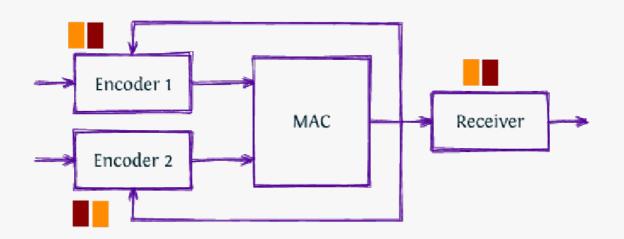


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Main idea

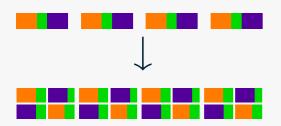
Each block now consists of three phases [Bross-Lapidoth-05]



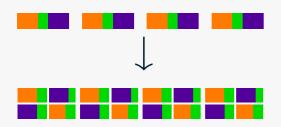




Bross-Lapidoth Region



Bross-Lapidoth Region



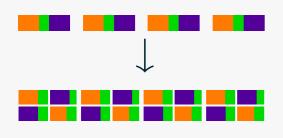
Capacity

$$R_{BL} = \bigcup_{\substack{P_{U}P_{X^{1}|U}P_{X^{2}|U}\\g_{1}:\mathcal{X}^{1}\times\mathcal{Y}\rightarrow\mathcal{Y}^{1},g_{2}:\mathcal{X}^{2}\times\mathcal{Y}\rightarrow\mathcal{Y}^{2}\\stuff}} \begin{cases} R_{1} \leqslant (1+\eta)^{-1}I(X^{1}\wedge Y \mid UX^{2}V^{1})\\ R_{2} \leqslant (1+\eta)^{-1}I(X^{2}\wedge Y \mid UX^{1}V^{2})\\ R_{1} + R_{2} \leqslant (1+\eta)^{-1}I(X^{1}X^{2}\wedge Y \mid V^{1}V^{2}) - R_{L}\\ R_{L} \leqslant complicated \ expression\\ \eta \geqslant complicated \ expression \end{cases}$$

[Bross-Lapidoth-05]

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Bross-Lapidoth Region



Comparison with CL

- BL region includes CL region
- In some cases the inclusion is strict

Capacity

$$R_{BL} = \bigcup_{\substack{P_{U}P_{\chi^{1}|U}P_{\chi^{2}|U}\\g_{1}:\mathcal{X}^{1}\times\mathcal{Y}\rightarrow\mathcal{Y}^{1},g_{2}:\mathcal{X}^{2}\times\mathcal{Y}\rightarrow\mathcal{Y}^{2}\\stuff}} \begin{cases} R_{1}\leqslant (1+\eta)^{-1}I(X^{1}\wedge Y\mid UX^{2}V^{1})\\ R_{2}\leqslant (1+\eta)^{-1}I(X^{2}\wedge Y\mid UX^{1}V^{2})\\ R_{1}+R_{2}\leqslant (1+\eta)^{-1}I(X^{1}X^{2}\wedge Y\mid V^{1}V^{2})-R_{L}\\ R_{L}\leqslant complicated\ expression\\ \eta\geqslant complicated\ expression \end{cases}$$

[Bross-Lapidoth-05]



Block Markov superposition codes

Cover-Leung Scheme

- ► Decoder decodes after one block
- ► Uses one auxiliary random variable



Block Markov superposition codes

Cover-Leung Scheme

- Decoder decodes after one block
- Uses one auxiliary random variable



Bross-Lapidoth Scheme

- ▶ Decoder decodes after two blocks
- Uses three auxiliary random variable and two partial functions



Block Markov superposition codes

Cover-Leung Scheme

- Decoder decodes after one block
- **►** Uses one auxiliary random variable

Bross-Lapidoth Scheme

- ► Decoder decodes after two blocks
- Uses three auxiliary random variable and two partial functions



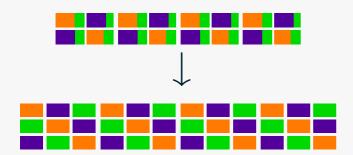


Questions

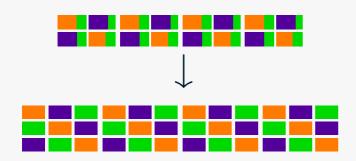
- ➤ Are these the best coding schemes that decode after one and two blocks?
- ► Can we improve the achievable region by using more auxiliary random variables?



Venkataramanan-Pradhan scheme



Venkataramanan-Pradhan scheme



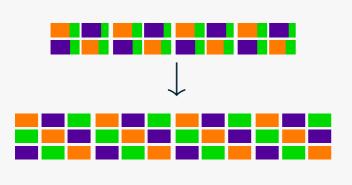
Capacity

$$R_{VP} = \bigcup \left\{ \begin{aligned} R_1 &\leqslant I(X^1 \wedge Y \mid UWX^2V^2) - [I(V^1 \wedge X^2 \mid YV^2WU) - I(U \wedge Y)]^+ \\ R_2 &\leqslant I(X^2 \wedge Y \mid UWX^1V^1) - [I(V^2 \wedge X^1 \mid YV^1WU) - I(U \wedge Y)]^+ \\ R_1 + R_2 &\leqslant I(X^1X^2 \wedge Y) - I(W \wedge Y \mid U) \end{aligned} \right\}$$

union over $P_U P_{V^1V^2} P_{X^1|UV^1} p_{X^2|UV^2} Q_{\tilde{V}^1|X^1UV^1V^2Y} Q_{\tilde{V}^1|X^1UV^1V^2Y}$, [Venkataramanan-Pradhan-09]



Venkataramanan-Pradhan scheme



Comparison with earlier schemes

- ▶ VP region includes CL region
- ► In some cases the VP region strictly improves over BL region

Capacity

$$R_{VP} = \bigcup \left\{ \begin{aligned} R_1 &\leqslant I(X^1 \wedge Y \mid UWX^2V^2) - [I(V^1 \wedge X^2 \mid YV^2WU) - I(U \wedge Y)]^+ \\ R_2 &\leqslant I(X^2 \wedge Y \mid UWX^1V^1) - [I(V^2 \wedge X^1 \mid YV^1WU) - I(U \wedge Y)]^+ \\ R_1 + R_2 &\leqslant I(X^1X^2 \wedge Y) - I(W \wedge Y \mid U) \end{aligned} \right\}$$

union over $P_U P_{V^1V^2} P_{X^1|UV^1} p_{X^2|UV^2} Q_{\tilde{V}^1|X^1UV^1V^2Y} Q_{\tilde{V}^1|X^1UV^1V^2Y}$, [Venkataramanan-Pradhan-09]



Questions

Structure of optimal coding scheme

Given that we are decoding after n-blocks, can we improve performance if we use a more sophisticated coding scheme?

Questions

Structure of optimal coding scheme

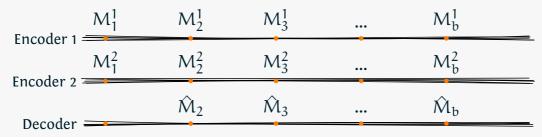
Given that we are decoding after n-blocks, can we improve performance if we use a more sophisticated coding scheme?

Nature of auxiliary random variables

Are the auxiliary random variables intrinsic to the problem or are they an artifact of our solution approach?

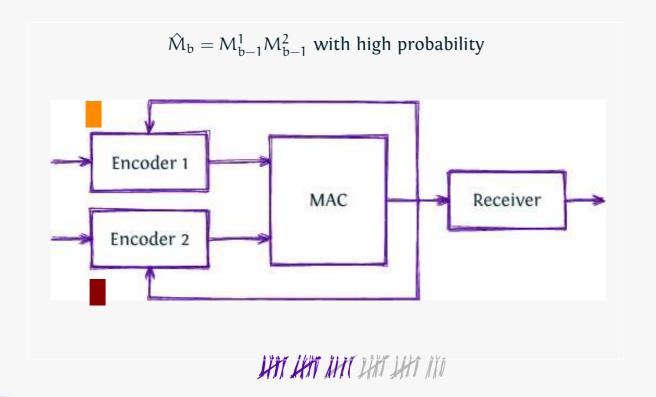
Outline

- 1. Overview of MAC with feedback
 - Cover-Leung scheme
 - Bross-Lapidoth scheme
 - ► Venkataramanan-Pradhan scheme
- 2. Formulation as a decentralized control problem with delayed sharing
 - delay = 1
 - delay = 2
- 3. Conclusion

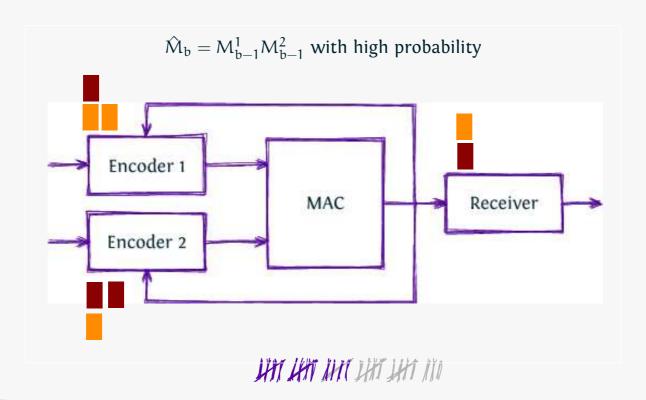


$$\hat{M}_b = M_{b-1}^1 M_{b-1}^2$$
 with high probability

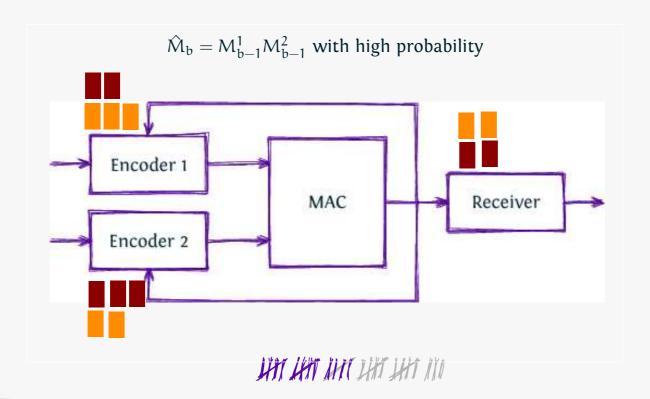
Encoder 1	M_1^1	M_2^1	M_3^1		M_b^1	
Encoder 2		M_2^2	M_3^2	•••	M_b^2	
Encoder 2		\hat{M}_2	\hat{M}_3		\hat{M}_{b}	
Decoder	_	-	•	•	•	



Encoder 1	M_1^1	M_2^1	M_3^1		M_b^1	
Encoder 2		M_2^2	M_3^2	•••	M_b^2	
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Decoder	_	-	•	•	•	



Encoder 1	M_1^1	M_2^1	M_3^1	•••	M_b^1	
Encoder 2		M_2^2	M_3^2	•••	M_b^2	
Encoder 2	 -	Μa	M ₃		Ń	
Decoder		1717	1713	•••	1416	



Encoder 1

$$\overline{M}_{[1:b-1]}, Y_{[1:b-1]},$$
 M_b^1

 $\overline{M}_{[1:b-2]}, Y_{[1:b]}$

Decoder

$$\overline{M}_{[1:b-1]}, Y_{[1:b-1]}, M_b^2$$

Encoder 2



Info struc at block level (general n)

Encoder 1

$$\overline{M}_{[1:b-n]}, Y_{[1:b-1]},$$
 $M^1_{[b-n+1:b]}$

 $\overline{M}_{[1:b-n-1]}, Y_{[1:b]}$

Decoder

$$\overline{M}_{[1:b-n]}, Y_{[1:b-1]},$$
 $M^2_{[b-n+1:b]}$

Encoder 2

Block Markov superposition coding at block level

Setup

- ▶ i.i.d. Messages at each encoders $\overline{M}_b = (M_b^1, M_b^2)$
- Fixed delay decoding \widehat{M}_b is trying to reproduce \overline{M}_{b-n} .

Block Markov superposition coding at block level

Setup

Additional information

- ightharpoonup i.i.d. Messages at each encoders ightharpoonup The encoder know \overline{M}_{b-n} $\overline{M}_b = (M_b^1, M_b^2)$
- **Fixed delay decoding** \hat{M}_b is trying to reproduce \overline{M}_{b-n} .
- ► The decoder knows \overline{M}_{b-n-1}

Block Markov superposition coding at block level

Setup

Additional information

- ▶ i.i.d. Messages at each encoders ▶ The encoder know \overline{M}_{b-n} $\overline{M}_b = (M_b^1, M_b^2)$
- ightharpoonup Fixed delay decoding \hat{M}_b is trying to reproduce \overline{M}_{b-n} .
- ightharpoonup The decoder knows \overline{M}_{h-n-1}

Coding scheme

- **Encoder**: $X_h^i = f_h^i(\overline{M}_{[1:b-n]}, Y_{[1:b-1]}, M_{[h-n+1:h]}^i)$
- ightharpoonup Decoder: $\hat{M}_b = d(\overline{M}_{[1:b-n-1]}, Y_{[1:b]})$
- $ightharpoonup Error: \rho(\overline{M}_{b-n}, \hat{M}_b)$

Block Markov superposition coding at block level

Setup

Additional information

- ▶ i.i.d. Messages at each encoders ▶ The encoder know \overline{M}_{b-n} $\overline{M}_b = (M_b^1, M_b^2)$
- **Fixed delay decoding** \hat{M}_b is trying to reproduce \overline{M}_{b-n} .
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Coding scheme

- **Encoder**: $X_h^i = f_h^i(\overline{M}_{[1:b-n]}, Y_{[1:b-1]}, M_{[h-n+1:h]}^i)$
- **Decoder:** $\hat{M}_b = d(\overline{M}_{[1:b-n-1]}, Y_{[1:b]})$
- \triangleright Error: $\rho(\overline{M}_{b-n}, \hat{M}_b)$
- **Objective**: Choose coding scheme to minimize $\sum \mathbb{E}[\rho(\overline{M}_{b-n}, \hat{M}_b)]$

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Transforming information structures

Encoder 1

$$\overline{M}_{[1:b-n]}, Y_{[1:b-1]},$$
 $M^1_{[b-n+1:b]}$

 $\overline{M}_{[1:b-n-1]}, Y_{[1:b]}$

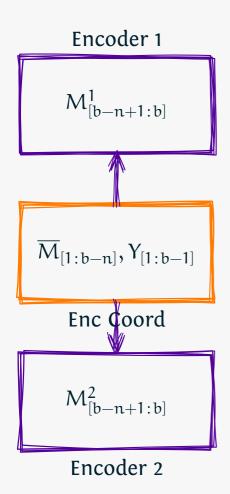
Decoder

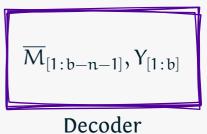
 $\overline{M}_{[1:b-n]}, Y_{[1:b-1]},$ $M^2_{[b-n+1:b]}$

Encoder 2

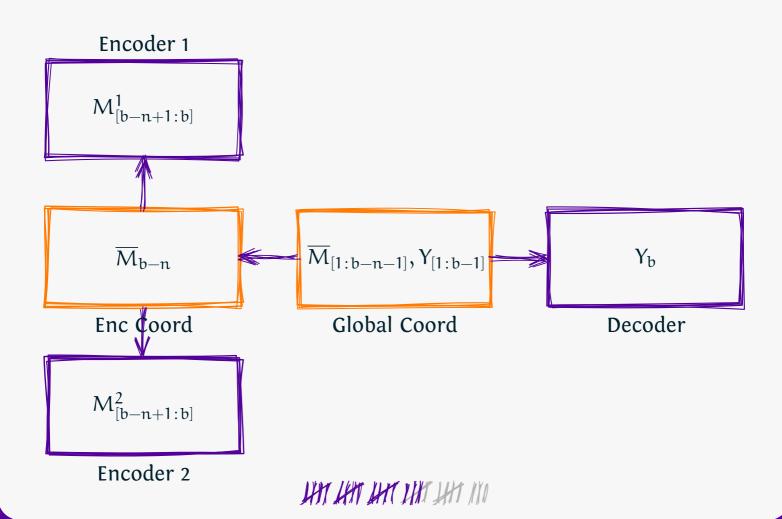


Transforming information structures





Transforming information structures



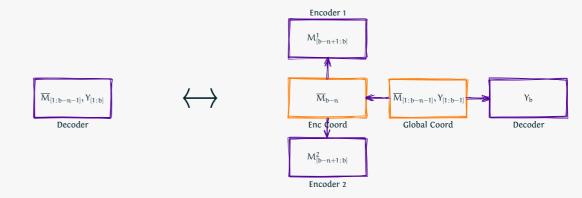
Centralized planning in decentralized systems



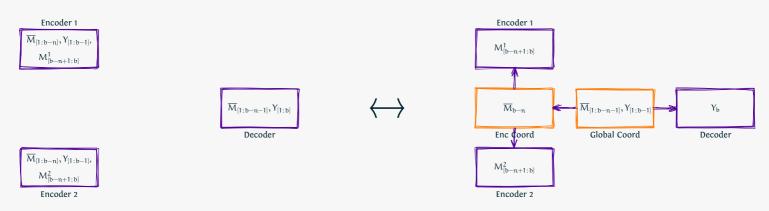


 $\overline{M}_{[1:b-n]}, Y_{[1:b-1]},$ $M^2_{[b-n+1:b]}$

Encoder 2

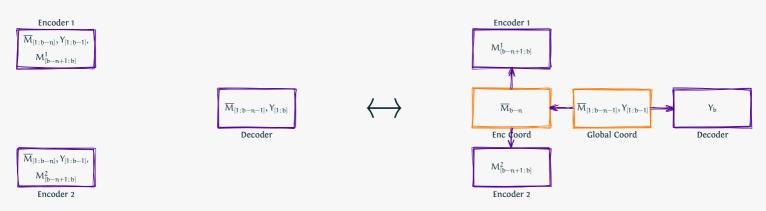


Centralized planning in decentralized systems



Hierarchical delayed sharing info structures

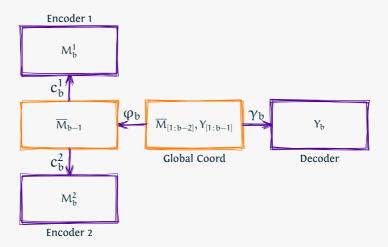
Centralized planning in decentralized systems

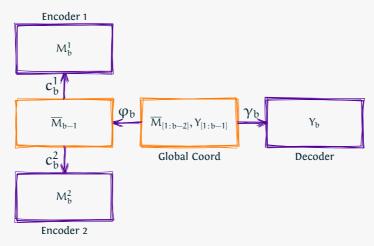


Hierarchical delayed sharing info structures

- ► Look at the system from global coordinator's point of view
- Show that the two optimization problems are equivalent
- Derive structure of optimal control policies for the coordinator







Global coordinator

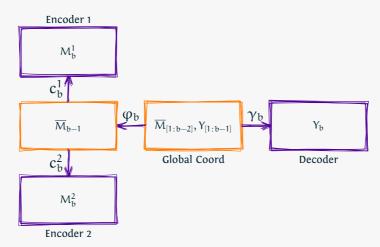
$$(\phi_b, \gamma_b) = \psi_b(\overline{M}_{[1:b-2]}, Y_{[1:b-1]})$$

where

$$\phi_{\mathfrak{b}}: \overline{\mathbb{M}} \to \big(\; (\mathbb{M}^1 \to \mathfrak{X}^1), (\mathbb{M}^2 \to \mathfrak{X}^2) \,\big)$$

$$\gamma_b: \mathcal{Y} \to \widehat{\mathfrak{M}}$$





Global coordinator

Encoder coordinator

$$(\varphi_b, \gamma_b) = \psi_b(M_{[1:b-2]}, Y_{[1:b-1]})$$

$$(c_b^1, c_b^2) = \phi_b(\overline{M}_{b-1})$$

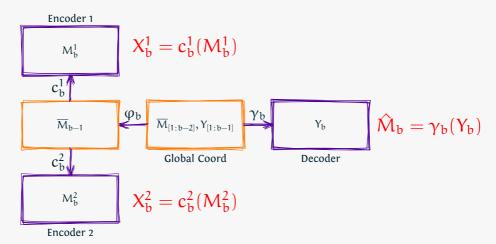
where

where

$$\phi_{b}: \overline{\mathbb{M}} \to \left((\mathbb{M}^{1} \to \mathbb{X}^{1}), (\mathbb{M}^{2} \to \mathbb{X}^{2}) \right)$$

$$c_{\mathrm{b}}^{\mathrm{i}}: \mathcal{M}^{\mathrm{i}}
ightarrow \mathfrak{X}^{\mathrm{i}}$$

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$$c_{
m b}^{
m i}: {\mathbb M}^{
m i} o {\mathfrak X}^{
m i}$$

$$\gamma_b: \mathcal{Y} \to \widehat{\mathfrak{M}}$$

IM HAT HAT HAT HAT MI

$$\left(\ \overline{M}_{[1:b-2]}, Y_{[1:b-1]}, \phi_{[1:b-1]}, \gamma_{[1:b-1]} \right) \xrightarrow{\text{info state}} (Y_{b-1}, c_{b-1}^1, c_{b-1}^2)$$

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Controlled Markov process

$$\begin{split} \mathbb{P}(Y_b, c_b^1, c_b^2 \mid \overline{M}_{[1:b-2]}, Y_{[1:b-1]}, \phi_{[1:b]}, \gamma_{[1:b]}) \\ &= \sum_{\overline{M}_{b-1,b}} \mathbb{P}(Y_b \mid c_b^1(M_b^1) c_b^2(M_b^2)) \times \mathbb{I}[(c_b^1, c_b^2) = \phi_b(\overline{M}_{b-1})] \\ &\times \frac{\mathbb{P}(Y_{b-1} \mid c_{b-1}^1(M_{b-1}^1) c_{b-1}^2(M_{b-1}^2)) \mathbb{P}(\overline{M}_{b-1})}{\sum_{\underline{M}_{b-1}} \mathbb{P}(Y_{b-1} \mid c_{b-1}^1(\underline{M}_{b-1}^1) c_{b-1}^2(\underline{M}_{b-1}^2)) \mathbb{P}(\underline{M}_{b-1})} \\ &= \mathbb{P}(Y_b, c_b^1, c_b^2 \mid Y_{b-1}, c_{b-1}^1, c_{b-1}^2, \phi_b) \end{split}$$



$$\left(\ \overline{M}_{[1:b-2]}, Y_{[1:b-1]}, \phi_{[1:b-1]}, \gamma_{[1:b-1]} \right) \xrightarrow{\text{info state}} (Y_{b-1}, c_{b-1}^1, c_{b-1}^2)$$

Controlled Markov process

$$\mathbb{P}(Y_b, c_b^1, c_b^2 \mid \overline{M}_{[1:b-2]}, Y_{[1:b-1]}, \varphi_{[1:b]}, \gamma_{[1:b]})$$

$$= \mathbb{P}(Y_b, c_b^1, c_b^2 \mid Y_{b-1}, c_{b-1}^1, c_{b-1}^2, \varphi_b)$$

$$\left(\ \overline{M}_{[1:b-2]}, Y_{[1:b-1]}, \phi_{[1:b-1]}, \gamma_{[1:b-1]} \right) \xrightarrow{\text{info state}} (Y_{b-1}, c_{b-1}^1, c_{b-1}^2)$$

© Controlled Markov process

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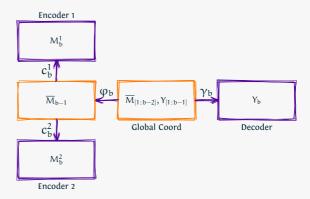
$$= \mathbb{P}(Y_b, c_b^1, c_b^2 \mid Y_{b-1}, c_{b-1}^1, c_{b-1}^2, \phi_b)$$

Expected cost per unit time

$$\begin{split} \mathbb{E}[\rho(\overline{M}_{b-1}, \hat{M}_b) \mid \overline{M}_{[1:b-2]}, Y_{[1:b-1]}, \phi_{[1:b]}, \gamma_{[1:b]}] \\ = \mathbb{E}[\rho(\overline{M}_{b-1}, \hat{M}_b) \mid Y_{b-1}, c_{b-1}^1, c_{b-1}^2, \phi_b, \gamma_b] \end{split}$$

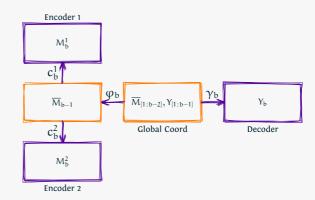


($Y_{b-1}, c_{b-1}^1, c_{b-1}^2$) is a controlled Markov process with control action (φ_b, γ_b).



- ($Y_{b-1}, c_{b-1}^1, c_{b-1}^2$) is a controlled Markov process with control action (φ_b, γ_b).
- Without loss of optimality

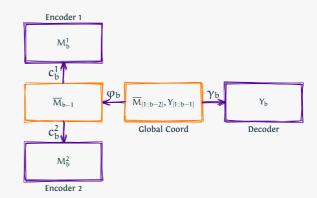
$$(\phi_b,\gamma_b)=\psi_b(Y_{b-1},c_{b-1}^1,c_{b-1}^2)$$



- ($Y_{b-1}, c_{b-1}^1, c_{b-1}^2$) is a controlled Markov process with control action (φ_b, γ_b) .
- Without loss of optimality

$$(\phi_b,\gamma_b)=\psi_b(Y_{b-1},c_{b-1}^1,c_{b-1}^2)$$

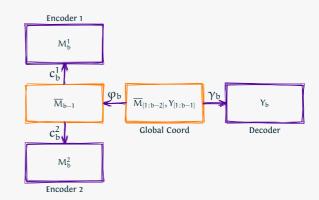
$$X_{b}^{i} = f_{b}^{i}(\overline{M}_{b-1}, Y_{b-1}, c_{b-1}^{1}, c_{b-1}^{2}, M_{b}^{1}), \quad i = 1, 2$$



- ($Y_{b-1}, c_{b-1}^1, c_{b-1}^2$) is a controlled Markov process with control action (φ_b, γ_b).
- Without loss of optimality

$$(\phi_b,\gamma_b)=\psi_b(Y_{b-1},c^1_{b-1},c^2_{b-1})$$

$$\begin{split} X_b^i &= f_b^i(\overline{M}_{b-1}, Y_{b-1}, c_{b-1}^1, c_{b-1}^2, M_b^1), \quad i = 1, 2 \\ &= f_b^i(U_b, M_b^i) \end{split}$$

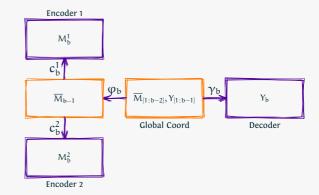


- ($Y_{b-1}, c_{b-1}^1, c_{b-1}^2$) is a controlled Markov process with control action (φ_b, γ_b) .
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- \bullet Fresh information : M_b^i



- ($Y_{b-1}, c_{b-1}^1, c_{b-1}^2$) is a controlled Markov process with control action (φ_b, γ_b).
- Without loss of optimality

$$(\phi_b,\gamma_b)=\psi_b(Y_{b-1},c_{b-1}^1,c_{b-1}^2)$$

$$\begin{split} X_b^i &= f_b^i(\overline{M}_{b-1}, Y_{b-1}, c_{b-1}^1, c_{b-1}^2, M_b^1), \quad i = 1, 2 \\ &= f_b^i(U_b, M_b^i) \end{split}$$

- Fresh information : M_h^i
- $\begin{array}{c} \text{ Auxiliary variable } U_b \equiv (\overline{M}_{b-1}, Y_{b-1}, c_{b-1}^1, c_{b-1}^2) \\ \equiv \big(\underbrace{\mathbb{P}(\overline{M}_{b-1}|Y_{b-1}, c_{b-1}^1(\cdot), c_{b-1}^2(\cdot))}_{\text{decoder's uncertainity list}}, \underbrace{\overline{M}_{b-1}}_{\text{index}} \big) \\ \end{array}$

Encoder 1

$$M_b^1$$
 C_b^1
 \overline{M}_{b-1}
 C_b^2

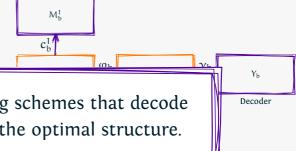
Global Coord

Decoder

 M_b^2

Encoder 2

- ($Y_{b-1}, c_{b-1}^1, c_{b-1}^2$) is a controlled Markov process with control action (ϕ_b, γ_b) .
- Without loss of optimality



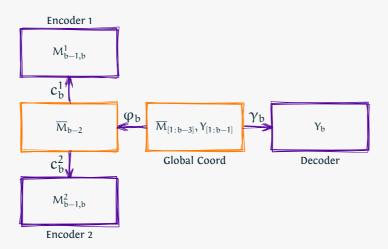
Encoder 1

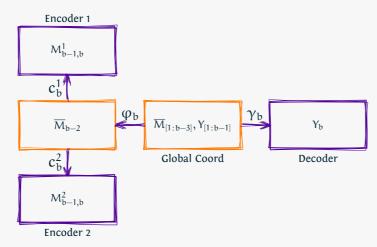
Within the class of all block Markov coding schemes that decode after one block, Cover-Leung scheme has the optimal structure.

$$\begin{split} X_b^i &= f_b^i(\overline{M}_{b-1}, Y_{b-1}, c_{b-1}^1, c_{b-1}^2, M_b^1), \quad i = 1, 2 \\ &= f_b^i(U_b, M_b^i) \end{split}$$

- \odot Fresh information : M_b^i
- $\begin{array}{c} \text{ Auxiliary variable } U_b \equiv (\overline{M}_{b-1}, Y_{b-1}, c_{b-1}^1, c_{b-1}^2) \\ \equiv \big(\underbrace{\mathbb{P}(\overline{M}_{b-1}|Y_{b-1}, c_{b-1}^1(\cdot), c_{b-1}^2(\cdot))}_{\text{decoder's uncertainty list}}, \underbrace{\overline{M}_{b-1}}_{\text{index}} \big) \\ \end{array}$







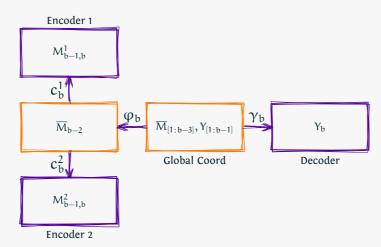
Global coordinator

$$(\phi_b, \gamma_b) = \psi_b(\overline{M}_{[1:b-3]}, Y_{[1:b-1]})$$

where

$$\begin{split} \phi_b : \overline{\mathcal{M}} &\to \big(\ (\mathcal{M}^1 \times \mathcal{M}^1 \to \mathcal{X}^1), (\mathcal{M}^2 \times \mathcal{M}^2 \to \mathcal{X}^2) \, \big) \\ \gamma_b : \mathcal{Y} &\to \hat{\mathcal{M}} \end{split}$$





Global coordinator

$$(\varphi_b, \gamma_b) = \psi_b(\overline{M}_{[1:b-3]}, Y_{[1:b-1]})$$

 $\varphi_{b}:\overline{\mathcal{M}}\to ((\mathcal{M}^{1}\times\mathcal{M}^{1}\to\mathcal{X}^{1}),(\mathcal{M}^{2}\times\mathcal{M}^{2}\to\mathcal{X}^{2}))$

where

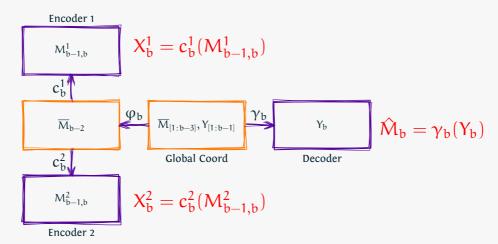
$$\gamma_{\mathrm{b}}: \mathcal{Y}
ightarrow \widehat{\mathfrak{M}}$$

Encoder coordinator

$$(c_b^1, c_b^2) = \varphi_b(\overline{M}_{b-2})$$

where

$$c_b^i: \mathcal{M}^i imes \mathcal{M}^i o \mathcal{X}^i$$



Global coordinator

$$(\varphi_b, \gamma_b) = \psi_b(\overline{M}_{[1:b-3]}, Y_{[1:b-1]})$$

 $\varphi_{b}:\overline{\mathcal{M}}\to ((\mathcal{M}^{1}\times\mathcal{M}^{1}\to\mathcal{X}^{1}),(\mathcal{M}^{2}\times\mathcal{M}^{2}\to\mathcal{X}^{2}))$

where

$$\gamma_b: \mathcal{Y} \to \hat{\mathfrak{M}}$$

Encoder coordinator

$$(c_b^1, c_b^2) = \varphi_b(\overline{M}_{b-2})$$

where

$$c_b^i: \mathcal{M}^i imes \mathcal{M}^i o \mathcal{X}^i$$

$$\left(\begin{array}{l} \overline{M}_{[1:b-3]}, Y_{[1:b-1]}, \phi_{[1:b-1]}, \gamma_{[1:b-1]} \right) \xrightarrow{\text{info state}} (Y_{b-2,b-1}, c_{b-1}^1, c_{b-1}^2, \widehat{c}_{b-2}^1, \widehat{c}_{b-2}^1) \\ \\ \text{where } \widehat{c}_{b-2}^i(\cdot) = c_{b-2}^i(M_{b-3}^i, \cdot), \ i=1,2. \end{array}$$

$$\left(\ \overline{M}_{[1:b-3]}, Y_{[1:b-1]}, \phi_{[1:b-1]}, \gamma_{[1:b-1]} \right) \xrightarrow{\text{info state}} (Y_{b-2,b-1}, c_{b-1}^1, c_{b-1}^2, \widehat{c}_{b-2}^1, \widehat{c}_{b-2}^1)$$
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$$\begin{split} \mathbb{P}(Y_{b-1,b},c_b^1,c_b^2,\hat{c}_{b-1}^1,\hat{c}_{b-1}^2 \mid \overline{M}_{[1:b-2]},Y_{[1:b-1]},\phi_{[1:b]},\gamma_{[1:b]}) \\ &= \mathbb{P}(Y_b,c_b^1,c_b^2,\hat{c}_{b-1}^1,\hat{c}_{b-1}^2 \mid Y_{b-2,b-1},c_{b-1}^1,c_{b-1}^2,\hat{c}_{b-2}^1,\hat{c}_{b-2}^1,\phi_b) \end{split}$$

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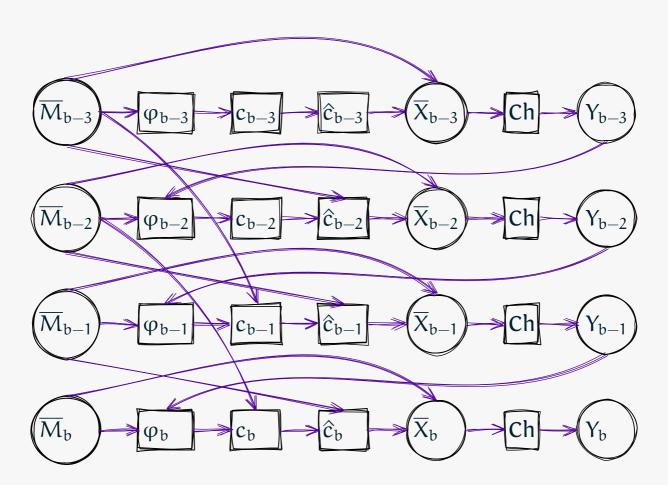
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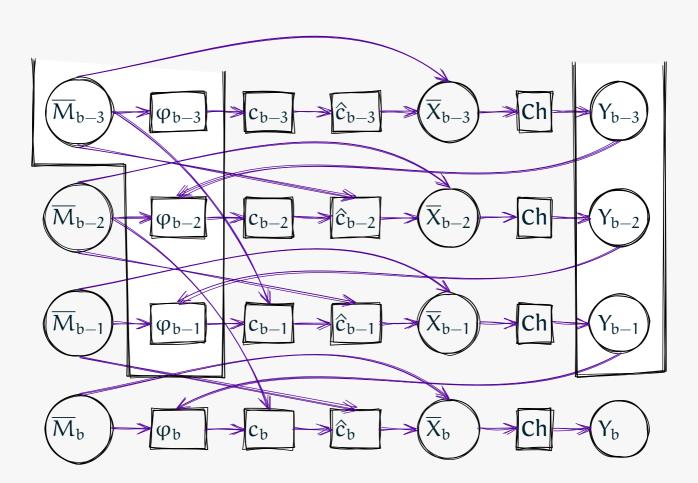
$$\begin{split} \mathbb{P}(Y_{b-1,b},c_b^1,c_b^2,\hat{c}_{b-1}^1,\hat{c}_{b-1}^2 \mid \overline{M}_{[1:b-2]},Y_{[1:b-1]},\phi_{[1:b]},\gamma_{[1:b]}) \\ &= \mathbb{P}(Y_b,c_b^1,c_b^2,\hat{c}_{b-1}^1,\hat{c}_{b-1}^2 \mid Y_{b-2,b-1},c_{b-1}^1,c_{b-1}^2,\hat{c}_{b-2}^1,\hat{c}_{b-2}^1,\hat{c}_{b-2}^2,\phi_b) \end{split}$$

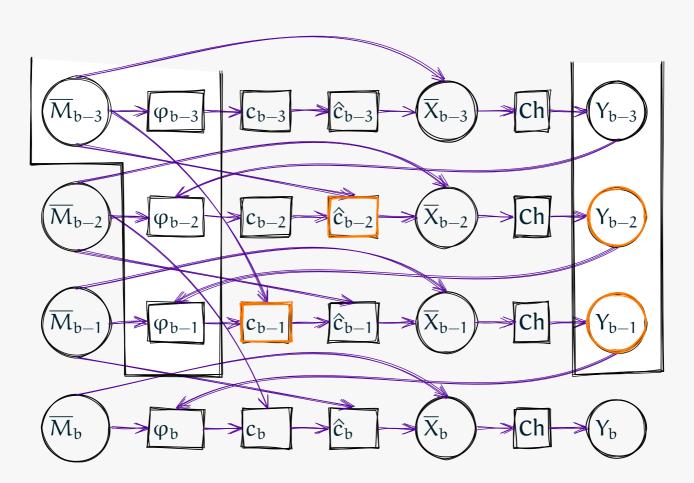
Expected cost per unit time

$$\begin{split} \mathbb{E}[\rho(\overline{M}_{b-2}, \hat{M}_b) \mid \overline{M}_{[1:b-3]}, Y_{[1:b-1]}, \phi_{[1:b]}, \gamma_{[1:b]}] \\ = \mathbb{E}[\rho(\overline{M}_{b-2}, \hat{M}_b) \mid Y_{b-2,b-1}, c_{b-1}^1, c_{b-1}^2, \hat{c}_{b-2}^1, \hat{c}_{b-2}^2, \phi_b, \gamma_b] \end{split}$$









(Solution $(Y_{b-2,b-1},c_{b-1}^{1,2},\hat{c}_{b-2}^{1,2})$ is a controlled Markov process with control action (ϕ_b,γ_b) .

- $(Y_{b-2,b-1},c_{b-1}^{1,2},\hat{c}_{b-2}^{1,2})$ is a controlled Markov process with control action (ϕ_b,γ_b) .
- Without loss of optimality $(\phi_b, \gamma_b) = \psi_b(Y_{b-1}, c_{b-1}^{1,2}, \hat{c}_{b-2}^{1,2})$

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- **Without loss of optimality** $(\phi_b, \gamma_b) = \psi_b(Y_{b-1}, c_{b-1}^{1,2}, \hat{c}_{b-2}^{1,2})$
- This is equivalent to

$$X_b^i = f_b^i(\overline{M}_{b-2}, Y_{b-2,b-1}, c_{b-1}^{1,2}, \hat{c}_{b-2}^{1,2}, M_{b-1,b}^i), \quad i = 1, 2$$

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- **©** Fresh information : M_h^i
- **a** Aux variable $V_b^i \equiv (M_{b-1}^i, c_{b-1}^{1,2}, Y_{b-1}, \overline{M}_{b-2})$
- and Aux variable $U_b \equiv (\overline{M}_{b-2}, Y_{b-2}, \hat{c}_{b-2}^{1,2})$

- 6 $(Y_{b-2,b-1},c_{b-1}^{1,2},\hat{c}_{b-2}^{1,2})$ is a controlled Markov process with control action (ϕ_b,γ_b) .
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- \bullet Fresh information : M_b^i
- **a** Aux variable $U_b \equiv (\overline{M}_{b-2}, Y_{b-2}, \hat{c}_{b-2}^{1,2})$

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- Fresh information : M_h^i



- Without loss of optimality $(\phi_b,\gamma_b)=\psi_b(Y_{b-1},c_{b-1}^{1,2},\hat{c}_{b-2}^{1,2})$
- This is aguivalant to

Within the class of all block Markov coding schemes that decode after two blocks,

Venkataramanan-Pradhan scheme has the optimal structure.

- Tresh information: /vth



Summary

Main Results

- ► CL has optimal structure for delay=1
- ▶ VP has optimal structure for delay-2

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Generalizations

- ► More than 2 blocks of delay
- More than 2 users

Summary

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- ▶ VP has optimal structure for delay-2

Main point

- Auxiliary random variables are related to information states
- Systematic approaches to derive info states can be used for auxiliary random variables

Generalizations

- ► More than 2 blocks of delay
- ► More than 2 users

Choice of info structures

- ► Block Markov superposition schemes enforce a specific information structure
- Stochastic control can identify whether an info structure is tractable (in many cases without explicitly solving the system)

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Conclusion

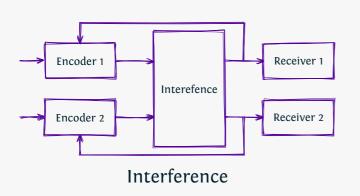
Horses for courses

- Use information theory to determine the form of coding scheme within blocks
- ▶ Use stochastic control to determine the form of coding scheme across blocks

Conclusion

Horses for courses

- Use information theory to determine the form of coding scheme within blocks
- ▶ Use stochastic control to determine the form of coding scheme across blocks



Future directions Receiver 1 Encoder Broadcast Receiver 2 **Broadcast** Relav Encoder Channel Channel Receiver Relay

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Thank you