# Basic of networks

Indiana University,

Emerson Melo

- There is a set of nodes,  $N = \{1, ..., n\}$ , where n is a finite number.
- Relationships between nodes are conceptualized in terms of binary variables, so that a relationship either exists or does not exist.
- Denote by  $g_{ij} \in \{0,1\}$  a relationship between two nodes i and j.

- The variable  $g_{ij}$  takes on a value of 1 if there exists a link between i and j and 0 otherwise.
- The set of nodes taken along with the links between them defines the network; this network is denoted by G and the collection of all possible networks on n nodes is denoted by  $\mathscr{G}$ .



The network G is undirected if  $g_{ij} = g_{ji}$  for all pair i, j



The network G is undirected if  $g_{ij} = g_{ji}$  for all pair i, j

## Definition

The network G is directed if the links have a direction



The network G is undirected if  $g_{ij} = g_{ji}$  for all pair i, j

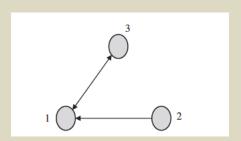
#### Definition

The network G is directed if the links have a direction

### Definition

The network G is weighted (directed or undirected) if  $g_{ij} \in \mathbb{R}_+$  for all pair of nodes











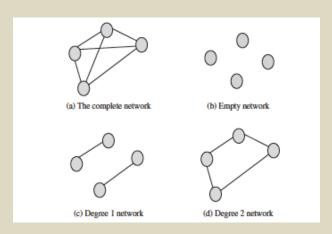
- The network G may be interpreted as a connection matrix
- Given a network G,  $G+g_{ij}$  and  $G-g_{ij}$  have the natural interpretation
- When  $g_{ij} = 0$  in G,  $G + g_{ij}$  adds the link  $g_{ij} = 1$ , while if  $g_{ij} = 0$  in G, then  $G + g_{ij} = G$ .

- Similarly, if  $g_{ij}=1$  in G,  $G-g_{ij}$  deletes the link  $g_{ij}$ , while if  $g_{ij}=0$  in G, then  $G-g_{ij}=G$
- Let  $N_i(G) = \{j \in N | g_{ij} = 1\}$  denotes the node with which node i has a link; this set will be referred to as the neighbors of i
- Let  $\eta_i(G) = |N_i(G)|$  denote the number of neighbors of node i in network G

- Let  $\mathbf{N}_1(G),\ldots,\mathbf{N}_{n-1}(G)$  be a division of nodes into distinct groups, where nodes belong to the same group if and only if they have the same number of links, i.e.,  $i,j\in\mathbf{N}_k(G)$ ,  $k=1,\ldots,n-1$ , if and only if  $\eta_i(G)=\eta_j(G)$
- With this notation in hand we can now describe a number of well-known networks.

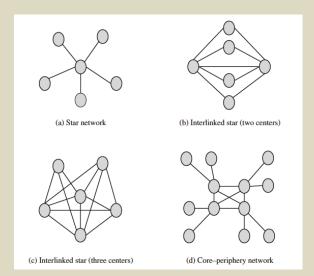
- A network G is said to be *regular* if every node has the same number of links, i.e.,  $\eta_i(G) = \eta$  for all  $i \in N$
- The *complete* network,  $G^c$  is a regular network in which  $\eta = n-1$
- The *empty* network,  $G^e$ , is a regular network in which  $\eta = 0$





- A core-periphery network structure describes the following situation. There are two groups of nodes,  $\mathbf{N}_1(G)$  and  $\mathbf{N}_k(G)$  with  $k > |\mathbf{N}_k(G)|$ . Nodes in  $\mathbf{N}_1(G)$  constitute the periphery and have a single link each and this link is with a node in  $\mathbf{N}_k(G)$
- Nodes in the set  $N_k(G)$  constitute the core and are fully linked with each other and with a subset of nodes in  $N_1(G)$





- An *interlinked stars* network consists of two groups  $\mathbf{N}_1(G)$  and  $\mathbf{N}_{n-1}(G)$  which satisfy the following condition:  $N_i(G) = \mathbf{N}_{n-1}(G)$  for  $i \in \mathbf{N}_k(G)$
- The star network is again a special case of such an architecture with  $|\mathbf{N}_{n-1}(G)|=1$  and  $|\mathbf{N}_1(G)|=n-1$
- In an interlinked star network, nodes which have n-1 links are referred to as central nodes or as hubs, while the complementary set of nodes are referred to as peripheral nodes or as spokes
- A *line* network consists of two groups of nodes  $N_1(G)$  and  $N_2(G)$ , with  $|N_1(G)| = 2$  and  $|N_2(G)| = n 2$

- The *degree* of node i is the number of i's direct connections; so  $\eta_i(G) = N_i(G)$  denotes the degree of node i in network G
- The degree distribution in a network is a vector P, where

$$P(k) = \frac{|\mathbf{N}_k(G)|}{n}$$

is the frequency/fraction of nodes with degree k

• Thus P(k) for each k, and  $\sum_{k=0}^{n-1} P(k) = 1$ 



- This degree distribution has support on  $\mathcal{D} = \{1, \dots, n-1\}$
- The average degree in network G is defined as

$$\hat{\eta}(G) = \sum_{k=0}^{n-1} P(k)k = \sum_{i \in N} \frac{\eta_i(G)}{n}$$
 (1)

- In a star network the degree distribution has support on degrees 1 and n-1, with n-1 nodes having degree 1, and 1 node having degree n-1
- The average degree in a star is  $2 \frac{2}{n}$

- An important concern in the study of networks is the variation in the degrees
- This variation has an important interpretation: degree may be related to node behavior and well-being
- One of the primary motivations for the study of networks in economics is the issue of how nodes extract advantages on account of their connections
- The variance in the degree distribution is defined as

$$Var(G) = \sum_{k=0}^{n-1} P(k) [\hat{\eta}(G) - k]^2$$
 (2)

- The degree variance in a star grows with n, while the variance in any regular network is 0 for all n
- ullet The range of degrees in network G is

$$R(G) = \max_{i \in N} \eta_i(G) - \min_{j \in N} \eta_j(G)$$
 (3)

- The range has a maximum value of n-2 and a minimum value of 0
- The range in a star is n-2, while the range in any regular network is 0

- The description of a network in terms of a degree distribution allows for an elegant way to study the addition and the redistribution of links
- The idea of adding links is captured in the relation of first-order stochastic domination
- Similarly, the idea of redistributing links is captured in the relation of mean-preserving spreads and second-order stochastic domination
- Given a degree distribution P, let the cumulative distribution function be denoted by  $\mathscr{P}: \{0, 1, \dots, n\} \longrightarrow [0, 1]$ , where

$$\mathscr{P}(\eta) = \sum_{x=0}^{\eta} P(x) \tag{4}$$



• Let P and P' be two degree distributions defined on  $\{0,1,\ldots,n\}$  and  $\mathscr P$  and  $\mathscr P'$  the two corresponding cumulative distribution functions, respectively



P first-order stochastically dominates (FOSD) P' if and only if

$$\mathscr{P}(k) \leq \mathscr{P}'(k),$$

for every  $k \in \{0, 1, 2, \dots, n-1\}$ 

#### Definition

P second-order stochastically dominates P' if and only if

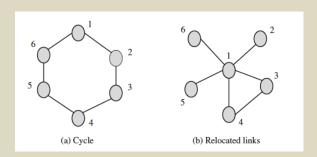
$$\sum_{k=0}^{x} \mathscr{P}(x) \le \sum_{k=0}^{x} \mathscr{P}'(x),$$

for every x

P' is a mean-preserving spread of P if and only if P and P' have the same mean and P second-order stochastically dominates P'

- A simple example of first-order shift in degree distribution arises when we move from a regular network with degree k to a regular network with degree k+1
- A simple example of a second-order shift arises when we move from a cycle with n=6 nodes to a network in which node 1 is linked to all nodes, nodes 2,3, and 4 have just this one link with node 1, while nodes 5 and 6 have two links each, a link with node 1 and a link with each other





- A major concern throughout the course will be the ways in which a node can be reached from another node in the network
- The first step in understanding this issue is the notion of walk:

   a walk is a sequence of nodes in which two nodes have a link
   between them in the network (they are neighbors)

- A node or a link may appear more than once in a walk
- A walk is the most general sequence of nodes and links possible in a network, subject to the constraint that any two consecutive nodes must have a link in the network
- The length of a walk is simply the number of links it crosses;
   this is simply equal to the number of nodes involved minus one

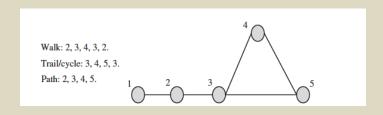
A walk in which all links are distinct is called a trail



- A trail in which there are three or more nodes and the initial and the end node are the same is called a cycle
- A trail in which every node is distinct is called a path
- Formally, there is a *path* between two distinct nodes i and j either if  $g_{ij}=1$  or if there is a set of distinct intermediate nodes  $j_1,\ldots,j_n$  such that  $g_{ij_1}=g_{j_1j_2}=\cdots g_{j_nj}=1$

- Next example represents a network with n = 5
- A possible walk in this network is 2,3,4,3,2
- This walk contains the links  $g_{23}$  and  $g_{34}$  twice and the nodes 2 and 3 also appear twice each
- This walk is therefore not a trail
- A possible trail in the network is 3,4,5,3
- However, since node 3 appears twice this trail is not a path
- A possible path in this network is 2, 3, 4, 5





- Two nodes belong to the same component if and only if there exists a path between them
- A network is connected if there exists a path between any pair of nodes  $i,j\in N$
- It follows that there exists only one component in a connected network

- In the case of an unconnected network, the components can be ordered in terms of their size, and the network has a giant component if, informally speaking, the largest component covers a relatively large fraction of the nodes while all other components are small
- The notion of minimality plays an important role in networks

- Intuitively speaking, minimality of a network reflects the idea that no link is "superfluous"
- A component is said to be minimal if the deletion of any single link in the component breaks the component into two components
- A network is said to be minimal if the deletion of any single link in the network increases the number of components by 1

- The geodesic distance between two nodes i and j in network G is the length of the shortest path between them, and will be denoted by d(i, j:G)
- If there is no path between i and j in network G, then by convention set  $d(i,j:G)=\infty$
- ullet When G is connected , the average distance between nodes of a network G is

$$d(G) = \frac{\sum_{i \in N} \sum_{j \in N} d(i, j; G)}{n(n-1)}$$
 (5)



- The centrality of a node in a network captures a number of ideas relating to the *prominence* of a node in a network
- Degree centrality captures the relative prominence of a node vis-a-vis other nodes in terms of the degree
- The (standardized) degree centrality of a node i in network G
  is simply the degree of this node divided by the maximum
  possible degree:

$$C_d(i;G) = \frac{\eta_i(G)}{n-1} \tag{6}$$



- We now turn to a measure of centrality which is based on proximity
- ullet The total distance from node i to all other nodes in the network G is

$$\sum_{j\neq i} d(i,j;G)$$

• This distance will be related to the number of nodes in a network and to facilitate comparison across networks of different size, it is useful to normalize the measure by multiplying with the minimum possible total distance, which is n-1



• The *closeness centrality* of node *i* in network *G* is defined as

$$C_c(i;G) = \frac{n-1}{\sum_{j \neq i} d(i,j;G)} \tag{7}$$

- This measure of centrality has a natural analogue at the aggregate network level
- This measure is built upon differences across nodes in a network and is normalized to account for maximum attainable differences



- The measure of closeness centrality is based solely on the length of the shortest paths between nodes in a network
- In some contexts it is quite possible that links are not perfectly reliable and so the number of paths of different lengths may all matter
- More generally, it is possible that actions of a person may have implications for the actions of her neighbors, which may in turn feedback on the initial individual, and so on



- These considerations motivate the study of a notion of centrality which allows for a richer range of direct and indirect influences in a network
- Bonacich (1972) developed such a measure of centrality and we now turn to it
- Bonacich's measure is one of the most popular measures in network economic



- ullet Consider the adjacency (connection) matrix  ${f G}$  of network G
- In this matrix an entry in a square corresponding to a pair  $\{i,j\}$  signifies the presence or absence of a link
- Let  $G^k$  be the kth power of the matrix
- The 0 power matrix  $\mathbf{G}^0 = \mathbf{I}$  , the  $n \times n$  identity matrix
- In  $\mathbf{G}^k$ , an entry  $g_{ij}^k$  measures the "number" of walks of length k that exist between players i and j in network

Introduction The Economic of networks



## Example

Consider a network with three players, 1,2, and 3. Suppose links take on values of 0 and 1, and let the network consist of two links,  $g_{12} = g_{23} = 1$ . We can find **G** and **G**<sup>2</sup>

Introduction The Economic of networks



	1	2	3
1	0	1	0
2	1	0	1
3	0	1	0

1 Basic definitions



	1	2	3
1	1	0	1
2	0	2	0
3	1	0	1



- Thus there is one walk of length 2 between 1 and 1 and between 3 and 3, but two walks of length 2 between 2 and 2
- There are no other walks of length 2 in this network



- Let  $a \ge 0$  be a scalar and let **I** be the identity matrix
- Define the matrix  $\mathbf{M}(G,a)$  as follows:

$$\mathbf{M}(G,a) = [\mathbf{I} - a\mathbf{G}]^{-1} = \sum_{k=0}^{\infty} (a\mathbf{G})^k$$
 (8)

- ullet This expression is well-defined so long as a is sufficiently small
- The entry  $m_{ij}(G,a) = \sum_{k=0}^{\infty} a^k g_{ij}^k$  counts the total number of walks in G from i to j, where walks of length k are weighted by a factor  $a^k$

Introduction The Economic of networks



• Given parameter a, the Bonacich centrality vector is defined as

$$C_B(G,a) = [\mathbf{I} - a\mathbf{G}]^{-1} \cdot \mathbf{1} = \mathbf{M}(G,a) \cdot \mathbf{1}, \tag{9}$$

where 1 is the (column) vector of 1s



• In particular, the Bonacich centrality of player i is

$$C_B(i;G,a) = \sum_{j=1}^n m_{ij}(G,a)$$
 (10)

 This measure of centrality counts the total number of (suitably weighted) walks of different lengths starting from i in network G.



• To see this note that (10) can be rewritten as follows:

$$C_B(i;G,a) = m_{ii}(G,a) + \sum_{j \neq i} m_{ij}(G,a)$$
 (11)

• Since  $\mathbf{G}^0 = \mathbf{I}$ , it follows that  $m_{ii}(G,a) \geq 1$  and so for every player i in any network,  $C_B(i;G,a) \geq 1$ . It is exactly equal to 1 when a=0



- The Bonacich centrality of a node can also be expressed as a function of the centrality of its neighbors
- Let ho(G) be the (largest) eigenvalue of the adjacency matrix  ${f G}$
- The Bonacich centrality of a node can then be defined as

$$C_B(i;G,a) = \frac{1}{\rho(G)} \sum_{j \in N} g_{ij} C_B(j;G,a)$$
 (12)