

# Distance Vector Routing

- a) The least-cost route between any two nodes is the route with **minimum distance**.
- b) Each node maintains a vector(table) of **minimum distances** to every node.
- c) The table at **each node also guides the packets** to the desired node by showing the showing the next hop routing.

Example:

Assume each **node as the cities**.

**Lines as the roads** connecting them.

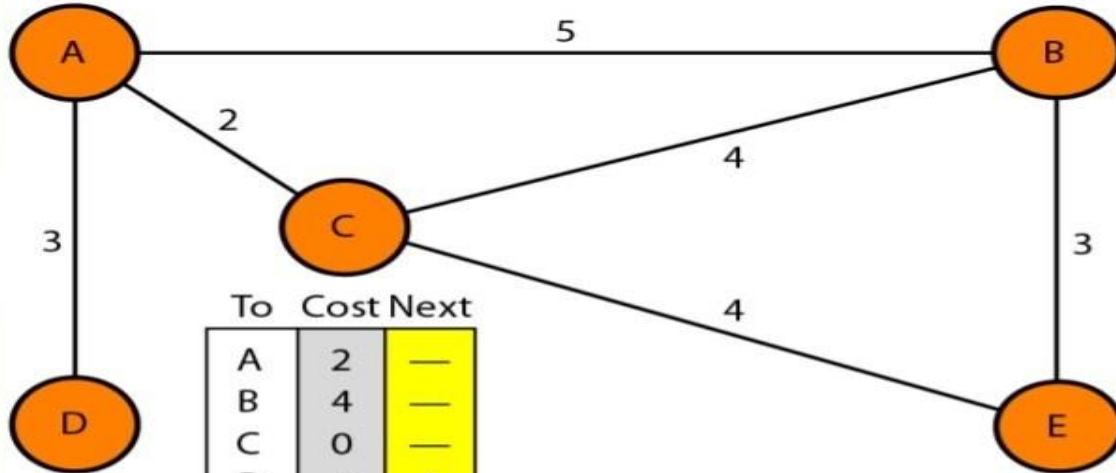
## Final Distance vector routing tables

To	Cost	Next
A	0	—
B	5	—
C	2	—
D	3	—
E	6	C

A's table

To	Cost	Next
A	3	—
B	8	A
C	5	A
D	0	—
E	9	A

D's table



To	Cost	Next
A	2	—
B	4	—
C	0	—
D	5	A
E	4	—

C's table

To	Cost	Next
A	5	—
B	0	—
C	4	—
D	8	A
E	3	—

B's table

To	Cost	Next
A	6	C
B	3	—
C	4	—
D	9	C
E	0	—

E's table



# Initialization

- a) The table in figure are stable.
- b) Each node knows how to reach any other node and their **cost**.
- c) At the beginning, each node know the cost of itself and its immediate neighbor.[those node directly connected to it.]
- d) Assume that each node send a message to the immediate neighbors and find the distance between itself and these neighbors.
- e) The distance of any entry that is not a neighbor is marked as **infinite**(unreachable).

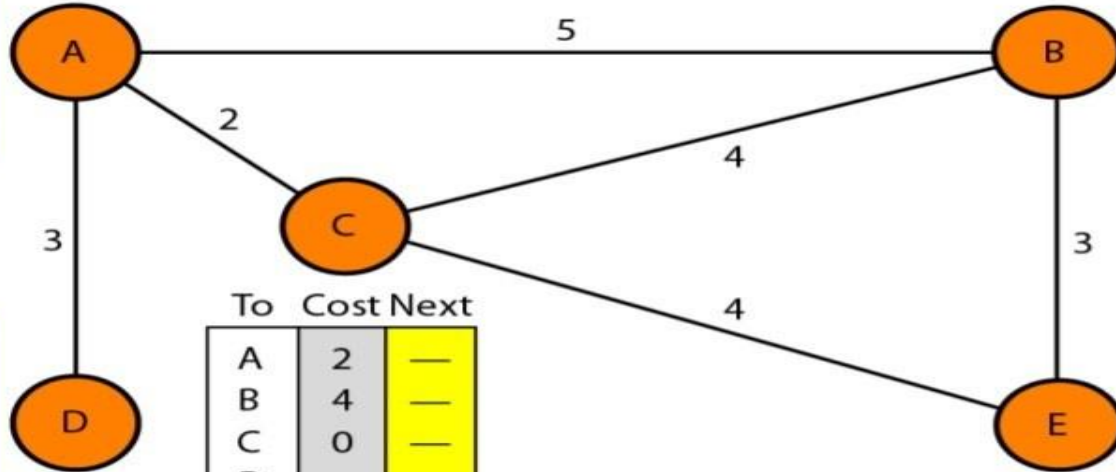
## Initialization of tables in distance vector routing (DVR)

To	Cost	Next
A	0	—
B	5	—
C	2	—
D	3	—
E	$\infty$	—

A's table

To	Cost	Next
A	3	—
B	$\infty$	—
C	$\infty$	—
D	0	—
E	$\infty$	—

D's table



To	Cost	Next
A	2	—
B	4	—
C	0	—
D	$\infty$	—
E	4	—

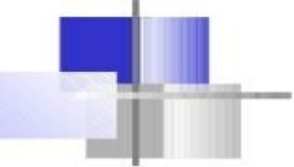
C's table

To	Cost	Next
A	5	—
B	0	—
C	4	—
D	$\infty$	—
E	3	—

B's table

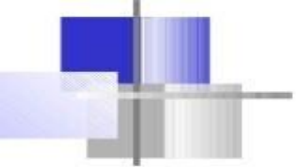
To	Cost	Next
A	$\infty$	—
B	3	B
C	4	C
D	$\infty$	—
E	0	D

E's table



# Sharing

- a) Idea is to share the information between neighbors.
- b) The node A does not know the distance about E, but node C does.
- c) If node C share it routing table with A, node A can also know how to reach node E.
- d) On the other hand, node C does not know how to reach node D, but node A does.
- e) If node A share its routing table with C, then node C can also know how to reach node D.
- f) Node A and C are immediate neighbors, can improve their routing tables if they help each other.



# Sharing

# Contd.,

- a) How much of the table must be shared with each neighbor?
- b) The third column of the table(next hop) is not useful for the neighbor.
- c) When the neighbor receives a table, this column needs to be replaced with the **sender's name**.
- d) If any of the rows can be used, the next node column filled with sender of the table.
- e) Therefore, a node can send only the **first two column** of its table to any neighbor.





# When to share

- a) Periodic Update: A node sends its table, normally every 30s, in a periodic update, it depends on the protocol that is using DVR.
- b) Triggered Update: A node sends its two-column routing table to its neighbors anytime there is a change in its routing table.
- c) This is called triggered update the change can result from the following:
  - ✓ A node receives a table from a neighbor, resulting in changes in its own table after updating.
  - ✓ A node detects some failure in the neighboring links which results in a distance change to infinity.



# Distance Vector Routing (DVR)

- a) 3 keys to understand how this algorithm works:
- **Sharing knowledge about the entire AS.** Each router shares its knowledge about the entire AS with neighbours. It sends whatever it has.
  - **Sharing *only with immediate neighbours*.** Each router sends whatever knowledge it has thru **all** its interface.
  - **Sharing at *regular intervals*.** sends at fixed intervals, e.g. every 30 sec.
- a) Problems: Tedious comparing/updating process, slow response to infinite loop problem, huge list to be maintained!!