



CS773-2025-Spring: Computer Architecture for Performance and Security

Lecture 5: No Flush Only Conflicts



ON SILENT MODE PLEASE

CASPER

No sharing ??

What If I do not share anything with you ??



Do not worry, I have Amazon Prime



Whaaaaaaf?!



amazon.com
Prime

Sorry:

amazon.com
Prime+Probe



CASPER

Prime + Probe



Step 0: Spy *fills* the entire shared cache



LLC



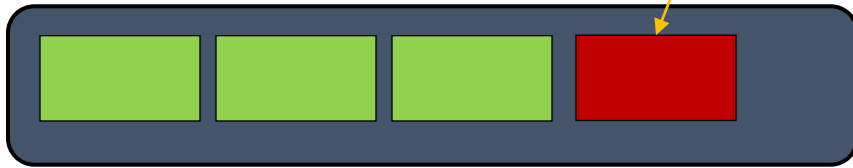
CASPER

Prime + Probe



Step 0: Spy *fills* the entire shared cache

Step 1: Victim *evicts* cache blocks while running

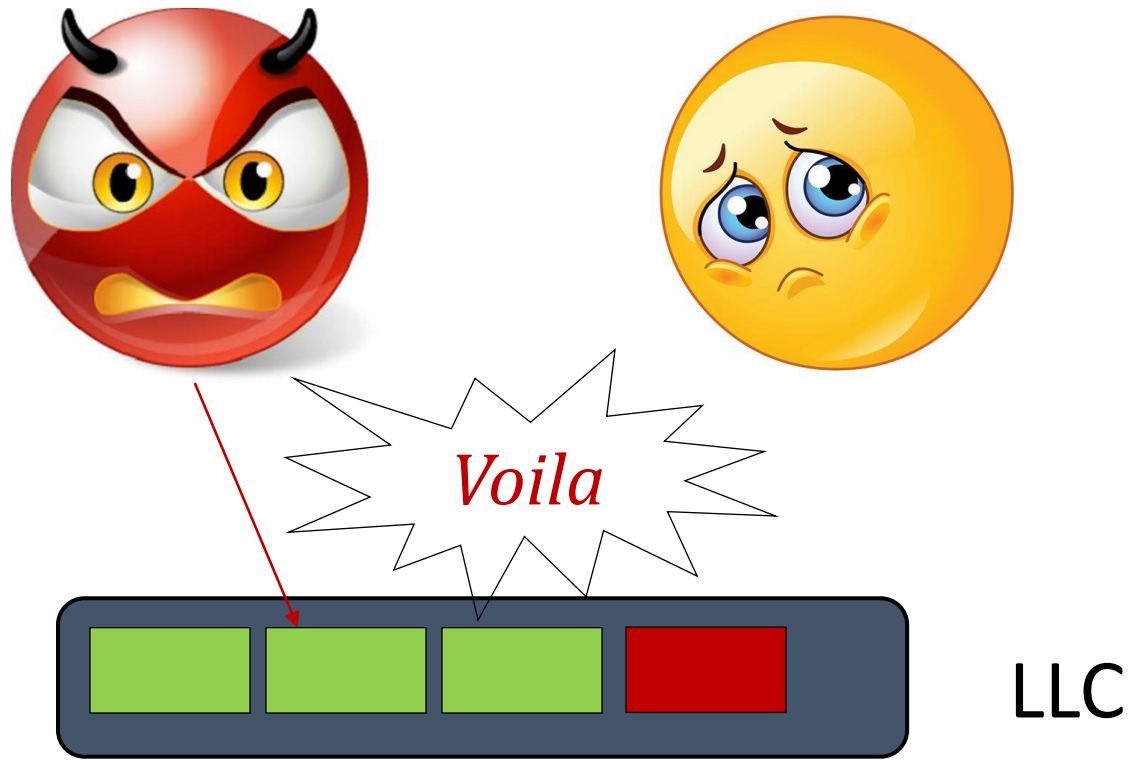


LLC



CASPER

Prime + Probe



Step 0: Spy *fills* a cache set

Step 1: Victim *evicts* cache blocks while running

Step 2: Spy *probes* the cache set

If misses then victim has accessed the set

Covert Channel



LLC

Step 0: Receiver gets data from L1 (fast, bit “0”)

Step 1: Sender thrashes LLC and back-invalidates L1

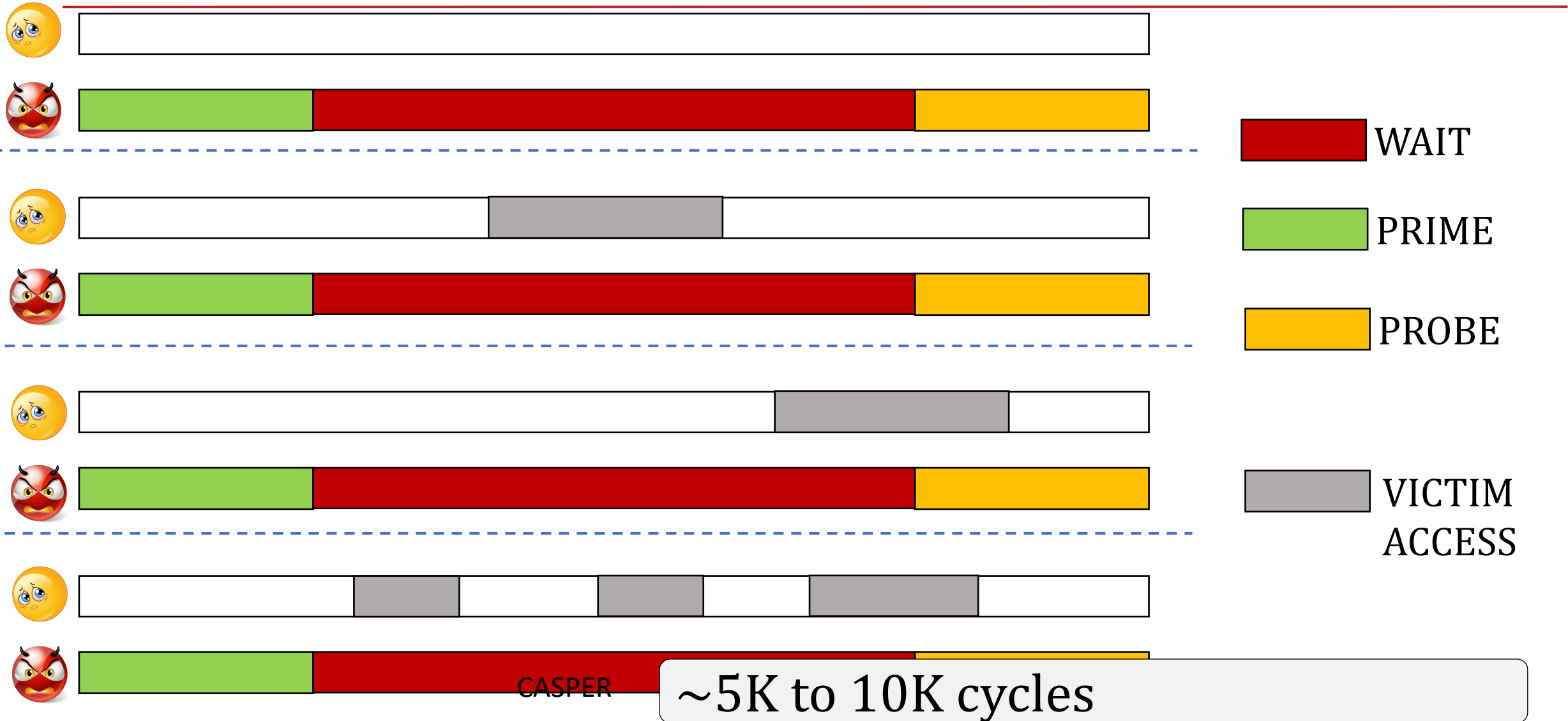
Step 2: Receiver gets data from DRAM (slow, bit “1”)



CASPER

Difficult to mount in a non-inclusive cache ☹

Notion of Time Gap



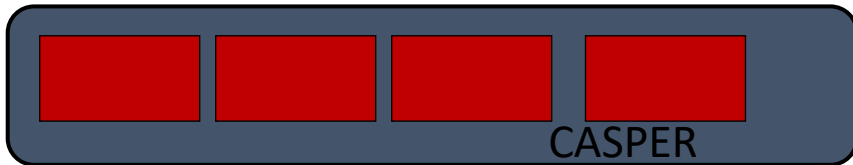
Inclusiveness



L1/L2



LLC

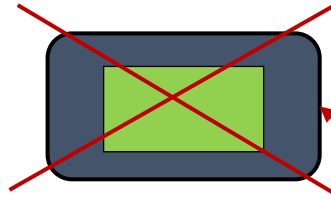


Miss

Inclusiveness

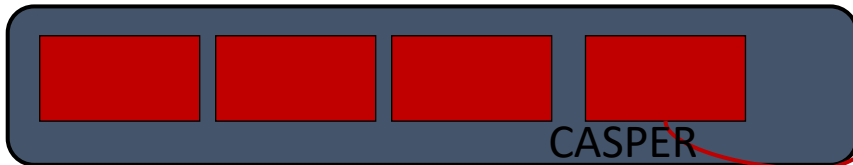


L1/L2



Cross-core back-invalidation

LLC



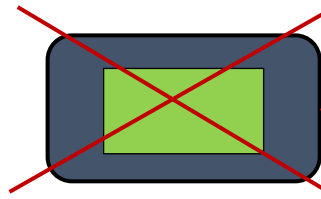
Miss

Inclusiveness



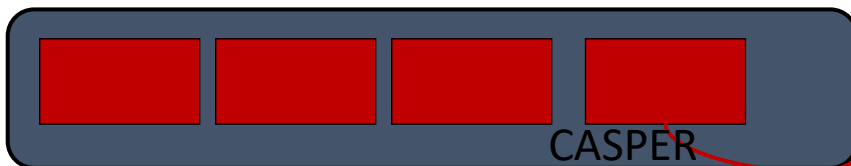
Attacker knows whether victim has accessed a set or not

L1/L2



Miss

LLC



Miss

Cache Occupancy Attack

```
int Trace[T*1000];
loop {
    counter = 0;
    t_begin = time();

    do {
        // count iterations
        counter++;
        // memory accesses
        for (i=0; i<size; i++) {
            tmp = buffer[i * 64]
        }
    } while(time() - t_begin < P);

    Trace[t_begin] = counter;
}
```

The attacker takes a parameter of period length P as input. It then constructs a trace, where each element in the trace measures how many iterations of the inner-most loop were executed every P milliseconds.

In the sweep-counting attack's code, the loop body contains an increment operation, memory accesses to a large buffer, and a call to the `time()` function.

Note that the buffer's size matches the size of the last-level cache so that one completion of the inner loop sweeps the entire last-level cache.

The counter value can thus be used to infer how many of the accessed cache lines reside in the cache.

What is the deal

Within a fixed window, you can have extremes of counter values.

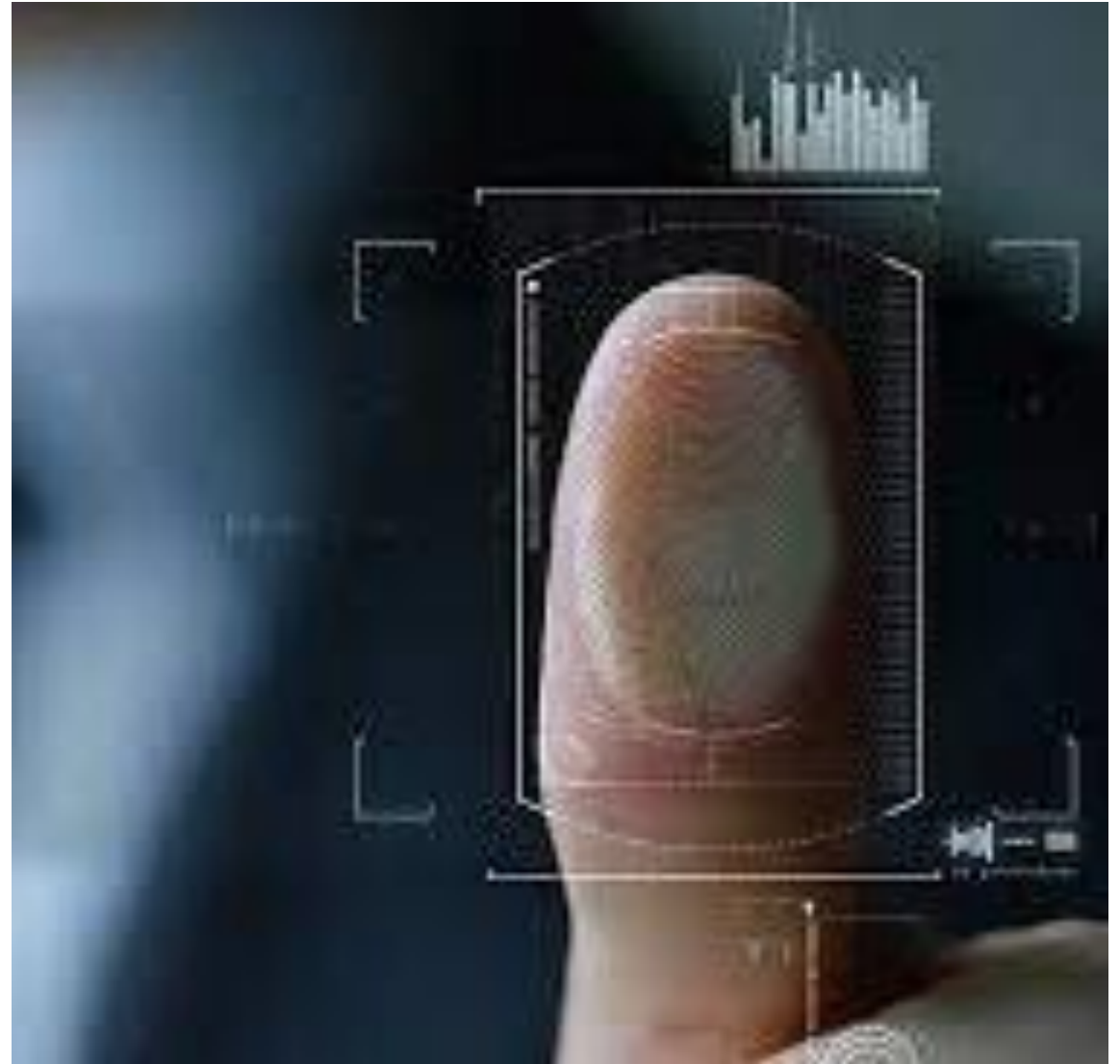
- (a) If the count is small, which means someone has kicked out a large part of LLC. So, occupancy of attacker is low, and it reached the window quickly, but the counter value is small.
- (b) If the count is large, which means no one else has kicked out the LLC lines, and the attacker has occupied the LLC completely. So, larger counter value.

There is mystery, please remind me in April. What is the utility?

Memorygrams 😊

Physical Fingerprinting

Identifying a person using his/her fingerprint, which is unique



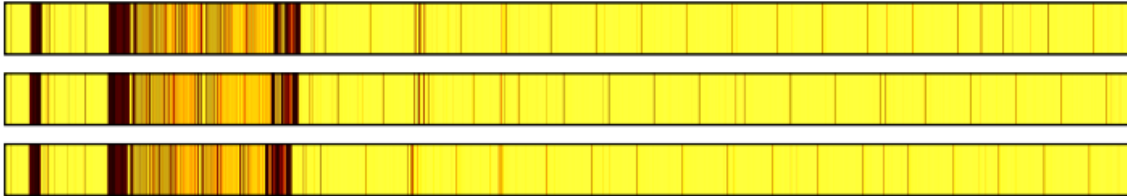


Digital Fingerprinting

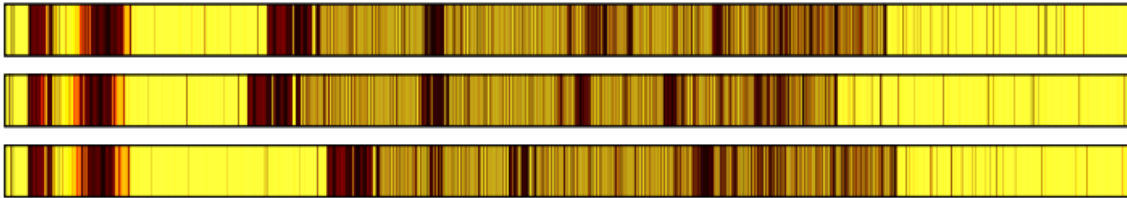
Many types - device, browser, website fingerprinting etc.

Here it is: Cache based website Fingerprinting

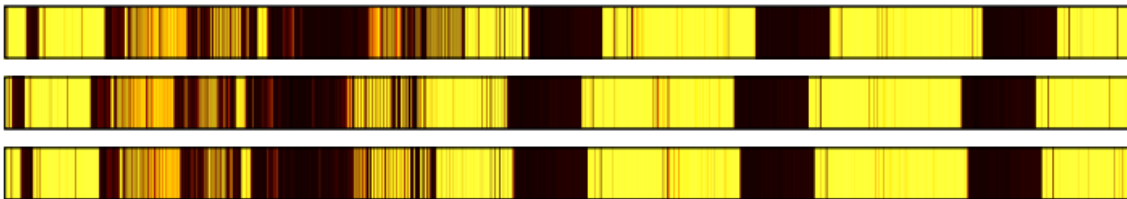
Wikipedia



Github



Oracle



Examples of memorygrams. Time progresses from left to right, shade indicates the number of evictions. (Darker shades correspond to more eviction.)

Use cases: Close World and Open World

- Close World

- Attacker needs to distinguish between a finite set of web pages.
- But attacker should have knowledge of a complete list of sites the victim visits.

- Open World

- Attacker wants to monitor access to set of sensitive websites, and classify them with high accuracy.

A photograph of a red pawn and a group of yellow pawns on a wooden surface. The red pawn is on the left, and the yellow pawns are on the right. The text "Back to Conflict based attack" is overlaid in the center.

Back to Conflict based attack

Thrashing Entire LLC (Prime+Probe): Questions of interest

Extremely Slow pre-attack step: Think about an 8MB/16MB LLC

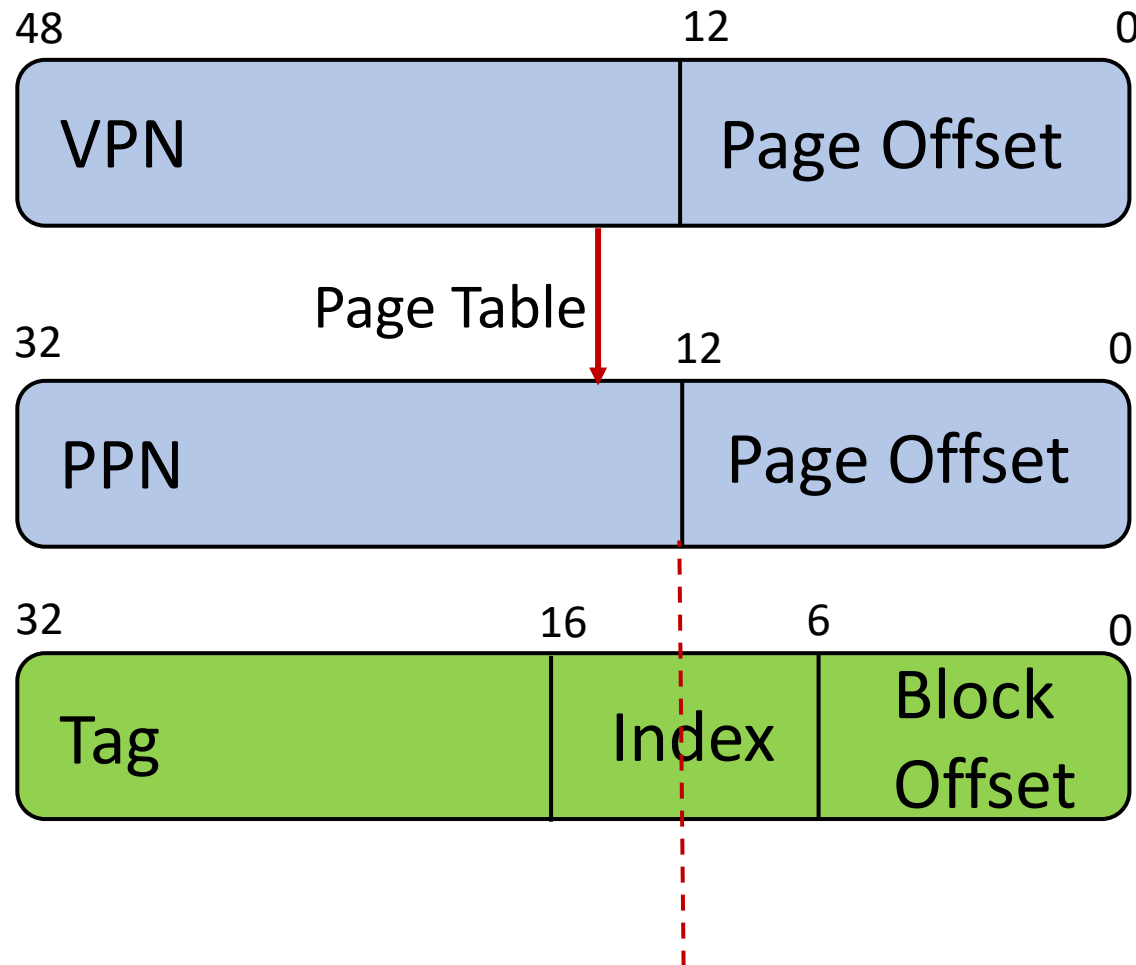
Why not thrash a group of addresses that are mapped to the same set?

Is there an algorithm to find out the same? Eviction set algorithm?

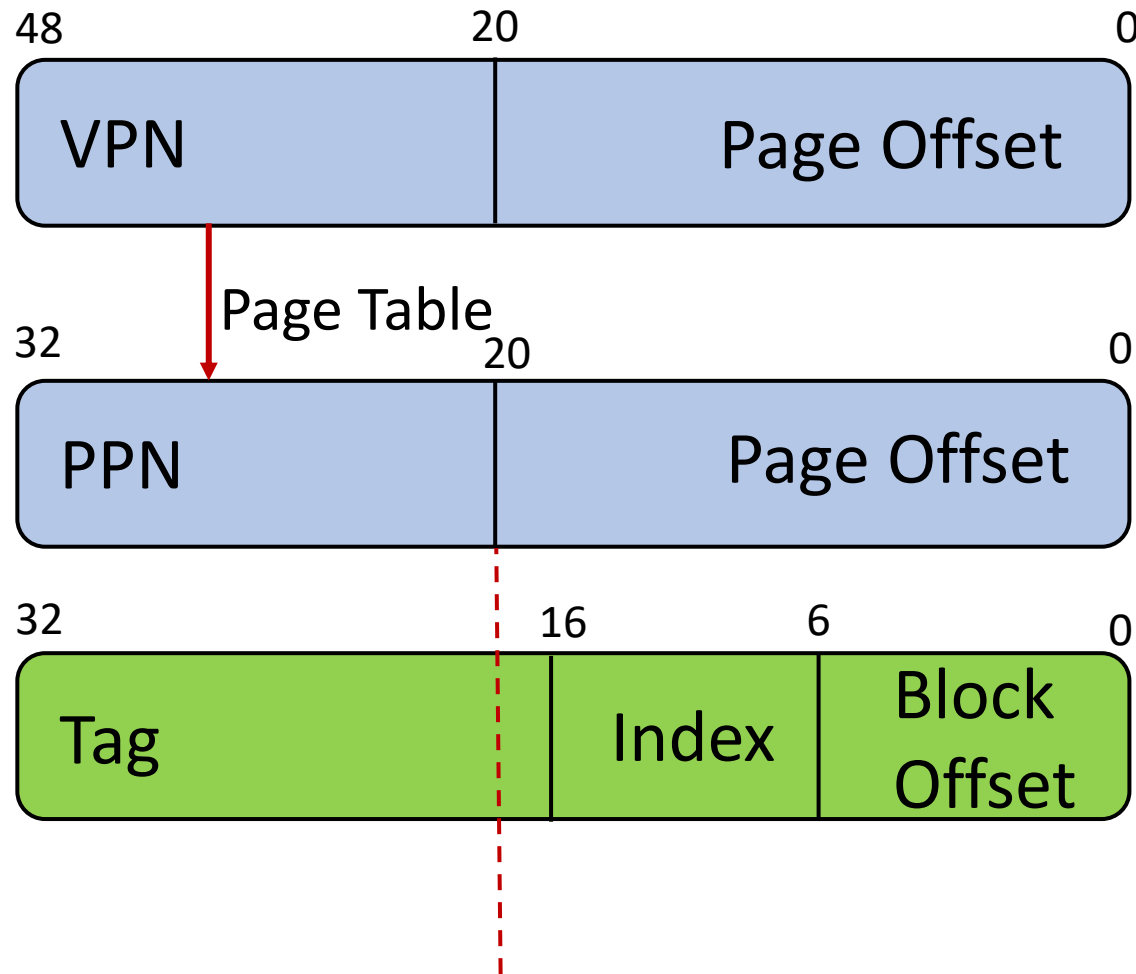
But what about virtual to physical address translation? LLC will have the physical address.

How to trigger requests that will go to the same set bypassing L1 and L2?

Attacker cannot control: LLC with 1024 sets



What if we have huge pages



Awesome. Now attacker can control all the accesses to a particular set.

What About?

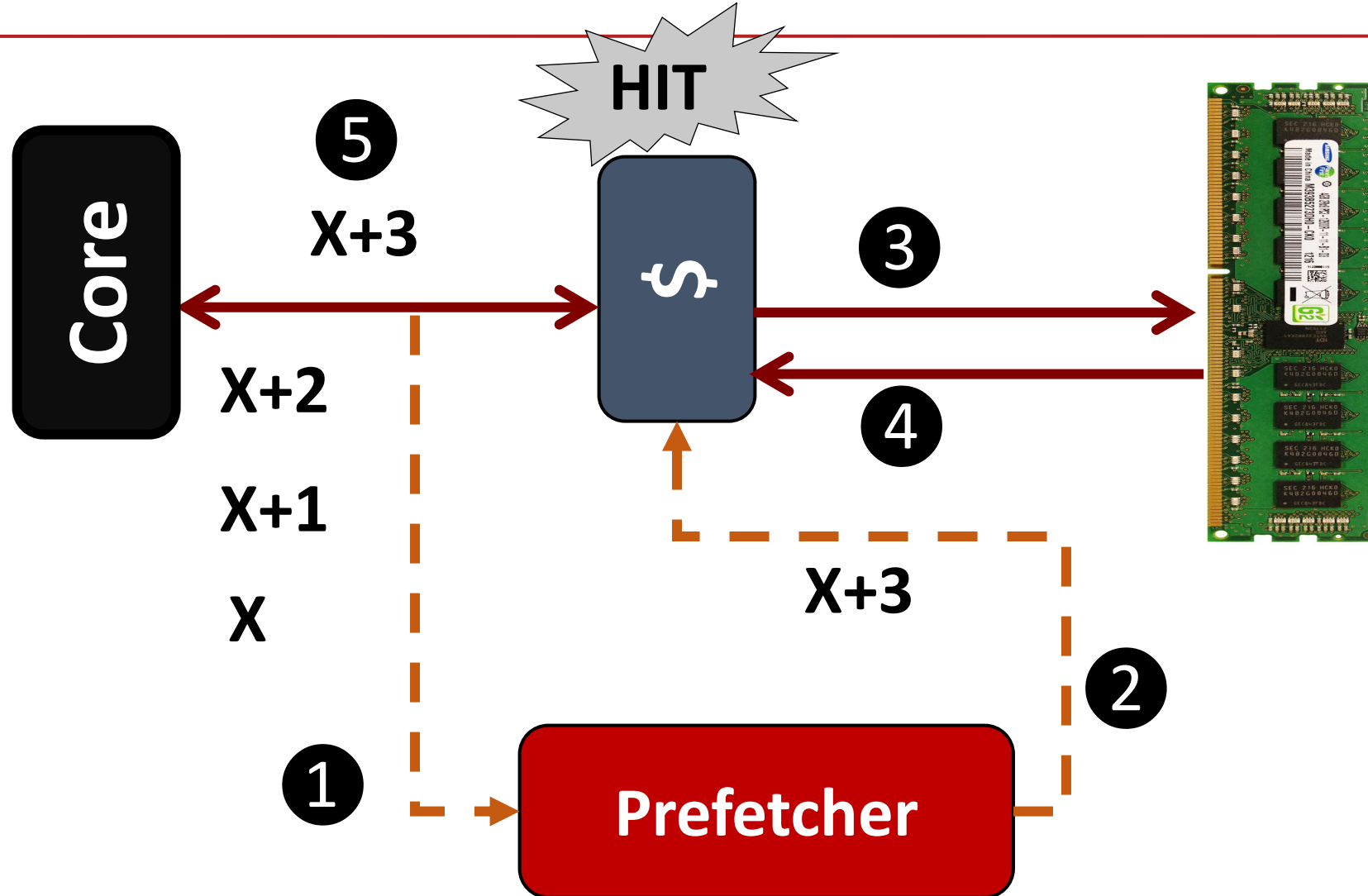
Effect of cache replacement policy at the LLC?

What if it is adaptive?

What if attacker's access pattern is predictable?

A hardware prefetcher can affect the eviction set creation process?

Hardware Prefetching



Effect of cache replacement policy at the LLC?

Fool the replacement policy too.

What if attacker's access pattern is predictable?

Fool the prefetcher too.

HOW GOOD IS
THE ATTACKER?

ASSUMPTIONS

AGILITY
(BANDWIDTH)

ADAPTIVE

ACCURACY

STEALTHY
(DETECTOR
CANNOT DETECT)