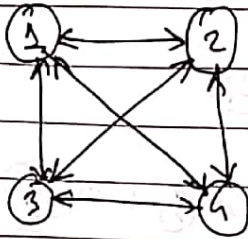


Travelling Salesperson problem:- Dynamic Programming..

salesperson has to visit each city at once & return to the same starting city, with minimum distance.

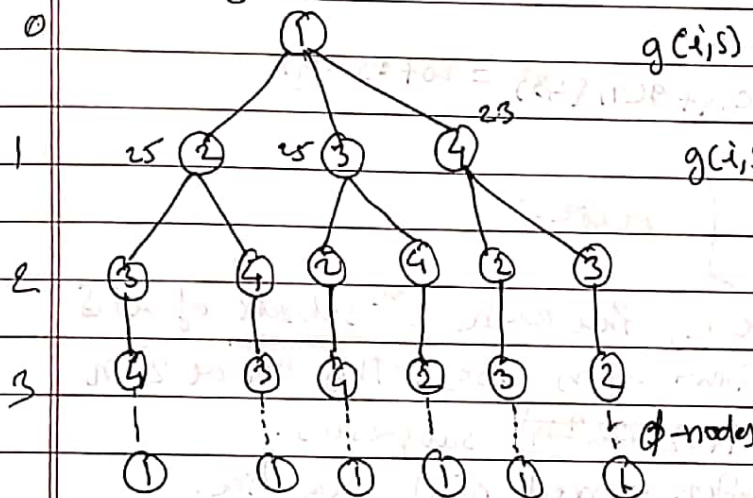


$n=4$ cities.

	1	2	3	4
1	0	10	15	20
2	5	0	9	10
3	6	13	0	12
4	8	8	9	0

Cost adjacency matrix.

starting city = 1.



$$g(i, S) = \min_{j \in S} \{ c_{ij} + g(j, S - \{j\}) \}$$

$g(i, S)$ - is the shortest path starting from i & going through all the vertices in S & terminate at vertex 1.

③ $g(2, \emptyset)$ - cost of getting from 2 to 1 = $c_{21} + 0 = 5$ $j=1$.

$g(3, \emptyset) = c_{31} = 6$

$g(4, \emptyset) = c_{41} = 8$

$|S|=0$

$g(2, \{3\}) = c_{23} + g(3, \emptyset) = 9 + 6 = 15$

$g(2, \{4\}) = c_{24} + g(4, \emptyset) = 10 + 8 = 18$

$|S|=1$.

② $g(3, \{2\}) = c_{32} + g(2, \emptyset) = 13 + 5 = 18$

$g(3, \{4\}) = c_{34} + g(4, \emptyset) = 12 + 8 = 20$

$g(4, \{3\}) = c_{43} + g(3, \emptyset) = 9 + 6 = 15$

$g(4, \{2\}) = c_{42} + g(2, \emptyset) = 8 + 5 = 13$

① $g(2, \{3, 4\}) = \begin{cases} c_{23} + g(3, \{4\}) = 9 + 20 = 29 & j=3 \\ c_{24} + g(4, \{3\}) = 10 + 15 = 25 & j=4 \end{cases}$ $|S|=2$

$\Rightarrow 25$ ✓ min

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$$g(3, \{2, 4\}) = \begin{cases} c_{32} + g(2, \{4\}) = 15 + 18 = 31 \\ c_{34} + g(4, \{2\}) = 12 + 13 = 25 \end{cases} \quad \min$$

L min

= 25

$$c_{42} + g(2, \{3\}) = 8 + 15 = 23$$

$$g(4, \{2, 3\}) = \begin{cases} c_{42} + g(2, \{3\}) = 8 + 15 = 23 \\ c_{43} + g(3, \{2\}) = 9 + 18 = 27 \end{cases}$$

$$g(1, \{2, 3, 4\}) = \begin{cases} \text{J=2} \\ c_{12} + g(2, \{3, 4\}) = 10 + 25 = 35 \\ \text{J=3} \\ c_{13} + g(3, \{2, 4\}) = 15 + 25 = 40 \\ \text{J=4} \\ c_{14} + g(4, \{2, 3\}) = 20 + 23 = 43 \end{cases} \quad \min$$

15/3

Route:-

1-2-4-3-1

For a problem of size n , there can be 2^n subsets of set S & each contains minimum of n fibres. Thus there are $2^n \cdot n$ subproblems.

To solve each subproblem we need, $O(n)$ linear time.

$$\therefore T(n) = 2^n \cdot n \cdot n = 2^n \cdot n^2 \Rightarrow O(2^n \cdot n^2)$$