



**Sunbeam Institute of Information Technology
Pune and Karad**

Module - Concepts of Operating System

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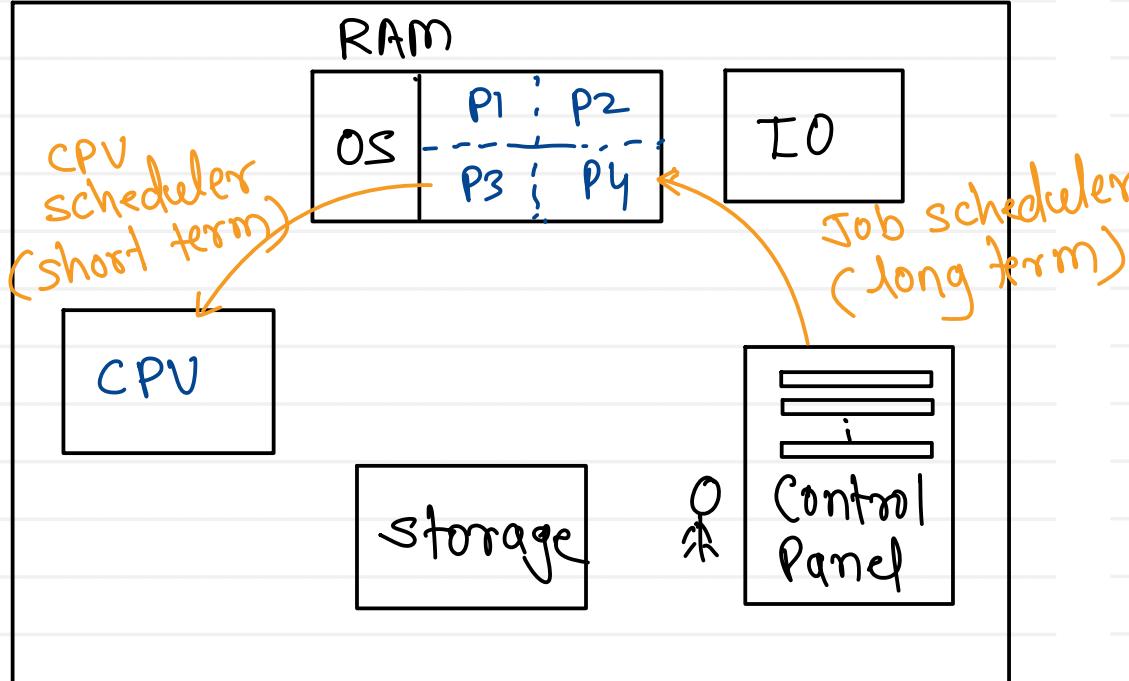


Types of Operating system

- 1 > Desktop OS (GPOS)
- 2 > Server OS
- 3 > Embedded OS
- 4 > Real time OS (FreeRTOS / Xenomai)
- 5 > Handheld OS
- 6 > Distributed OS



Types of Operating system



1. Resident monitor
2. Batch systems
3. Multi programming system
 - multiple programs are loaded into RAM (memory)
 - CPU burst : time spent by process on CPU
 - IO burst : time spent by process for IO
 - CPU burst > IO burst : CPU bound process
 - IO burst > CPU burst : IO bound process
 - mixture of CPU bound & IO bound processes was loaded inside RAM

Degree of multi programming :
No. of processes which can be loaded into RAM

Types of Operating system

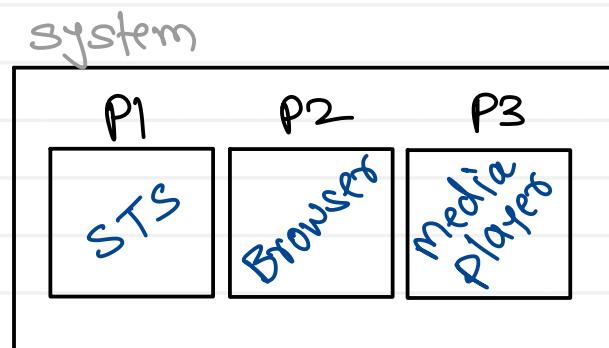
4. Time sharing system : (Multi tasking system)

- CPU time is shared in all the processes of RAM

Response time < 1 sec

- multiple tasks are getting executed in single system concurrently.

i) Process based multitasking

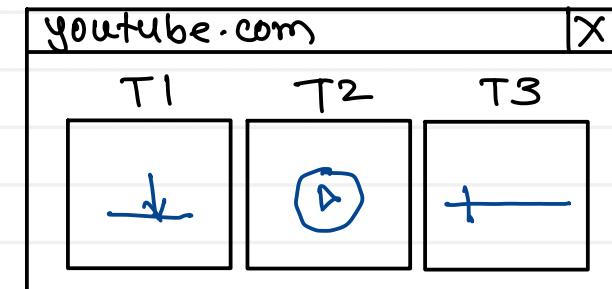


- system wide multitasking

- light weight process

ii) Thread based multitasking. (multithreading)

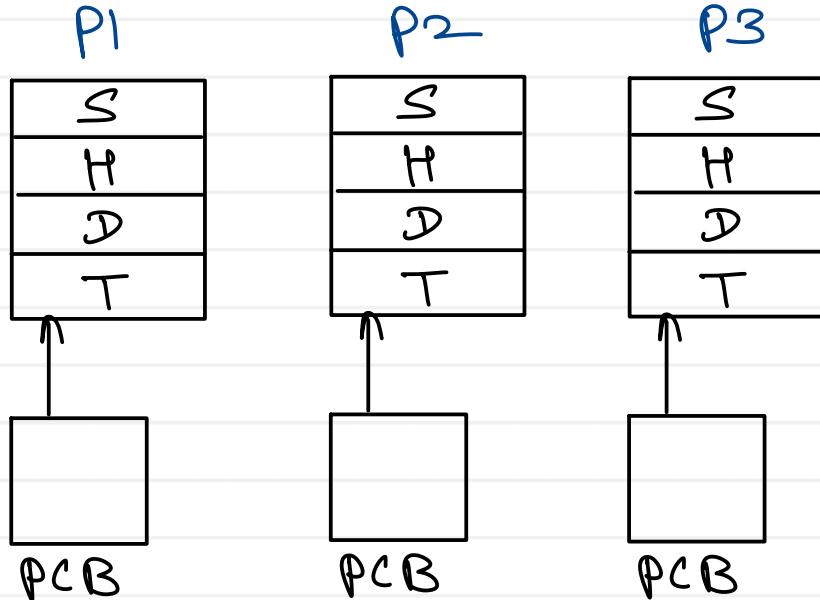
Browser



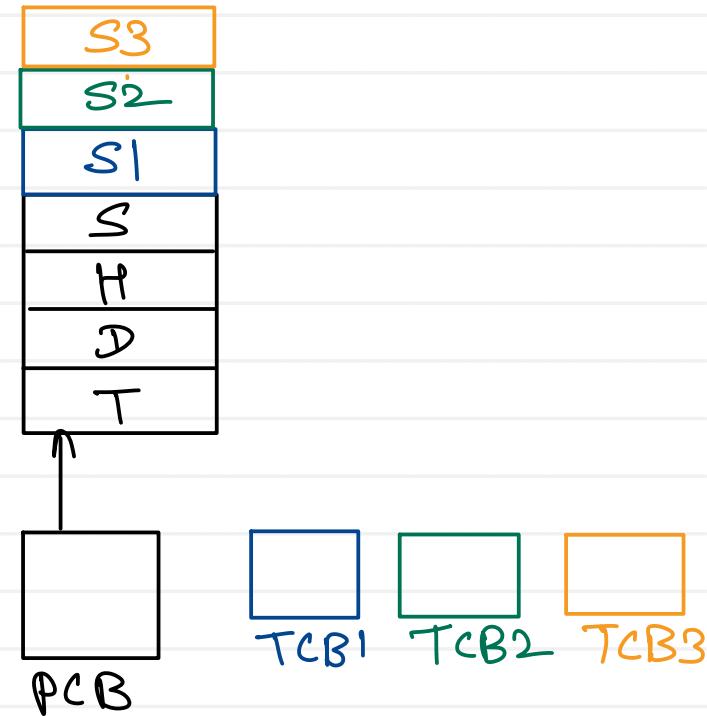
- multitasking within process

Process Vs Thread

- processes are created in system.



- threads are created always inside process



- for process all sections are created in user space + PCB is created in kernel space.

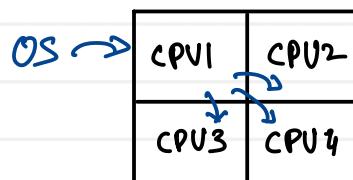
- for thread only stack section is created in user space & TCB is created in kernel space.
- remaining sections of user space are shared with parent process.

Types of Operating system

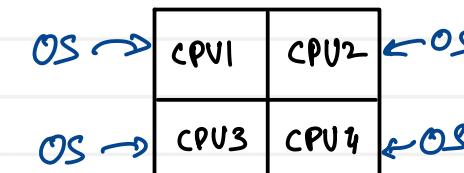
5. Multiprocessing systems :

- multiple processing units (CPU) are putted on single chip, which is called as multiprocessor / multi core.
- OS can take advantage of this, by scheduling process for each core
- multiple processes can be executed parallelly . (parallel systems)

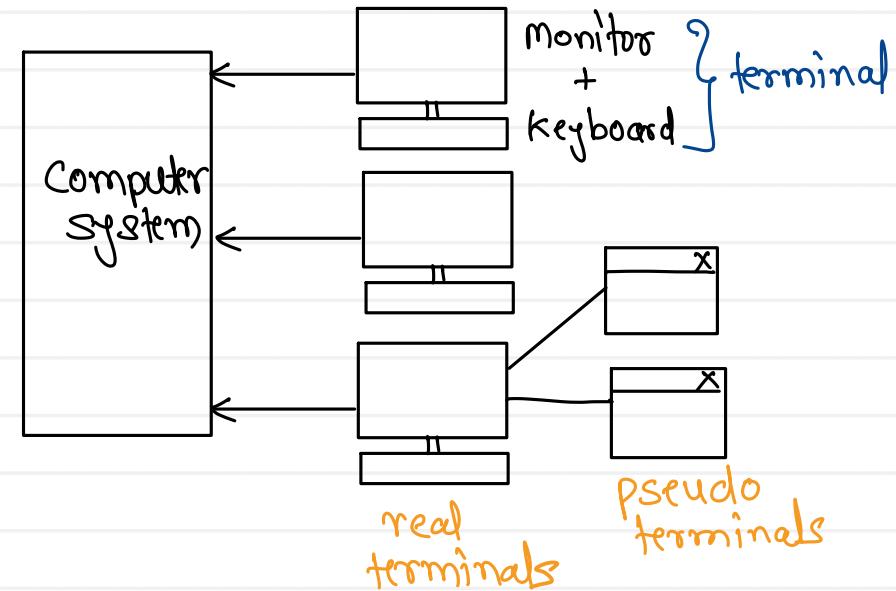
i. Asymmetric multi processing



Windows Vista
linux 2.6



6. Multiuser system:



ps	uname	who
ps -e	uname -r	tty
top	uname -a	whoami

Context switching

Execution context :

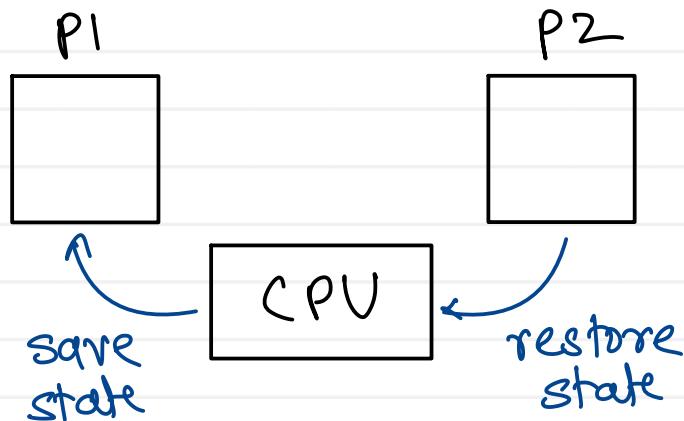
- values of CPU registers

↑
small memory elements
present inside CPU

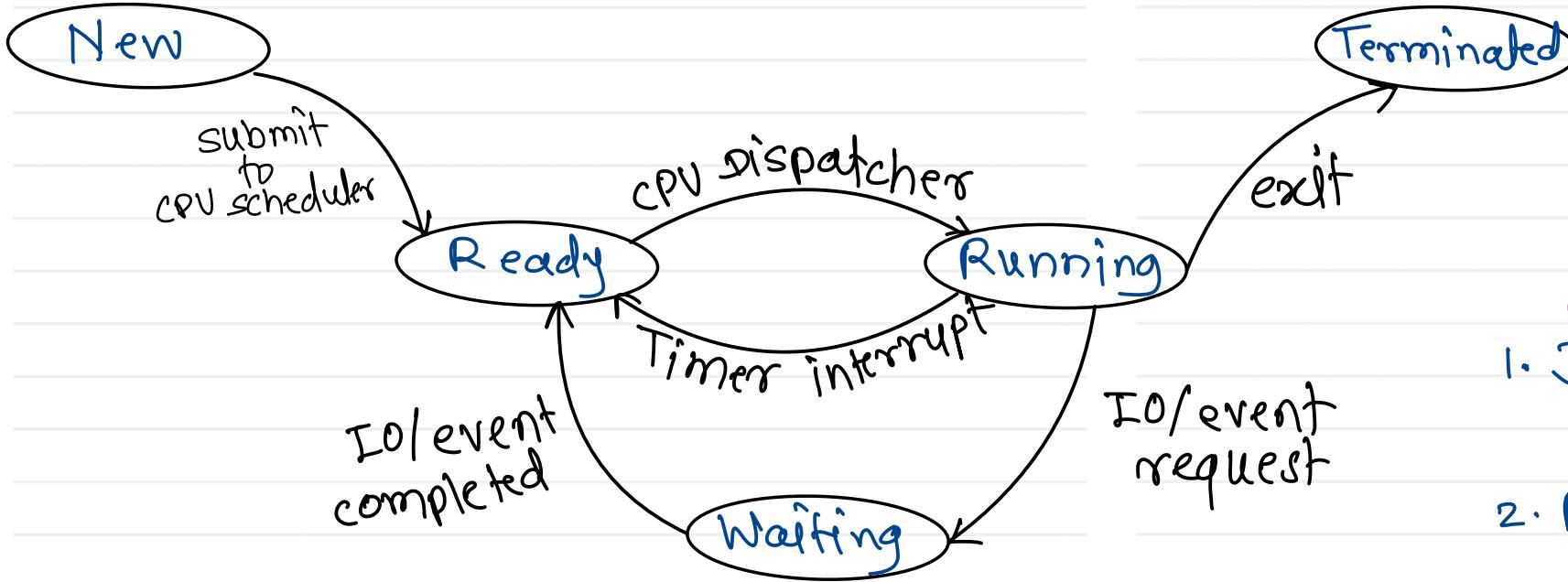
CPU dispatcher loads process on CPU. Execution context of selected from its PCB is loaded on CPU

Context switching :

- changing the process of CPU
- execution context of running process is saved into its PCB & execution context of selected process is restored on CPU.



Process life cycle



OS data structures

1. Job queue / process list:
 - all processes of RAM
2. Ready queue:
 - processes which are in ready state.
3. Waiting queue(s):
 - processes which are in waiting state.

CPU scheduling

Cases in which CPU scheduler is called

1. Running → Terminated } voluntarily
2. Running → waiting
3. Running → Ready } forcefully
4. Waiting → Ready

CPU scheduling criterias:

1. CPU Utilization (Max)
2. Throughput (Max)
 - amount of work done per unit
3. Waiting time (Min)
 - time spent by process in ready queue
$$WT = TAT - CPU \ burst$$
4. Response time (Min)
 - time from arrival to first time execution
$$RT = start \ time - AT$$
5. Turn Around Time (TAT) (Min)
 - time spent by process inside RAM

$$TAT = DT - AT$$

Types of scheduling :

1. Non pre emptive scheduling:
CPU access is always given voluntarily (co operative process)
2. Pre emptive scheduling:
CPU access is given forcefully,

CPU scheduling algorithms:

1. FCFS
2. SJF
3. Priority
4. RR

FCFS (First Come First Serve) (Non preemptive)

Process	Arrival	CPU Burst
P1	0	24
P2	0	3
P3	0	3

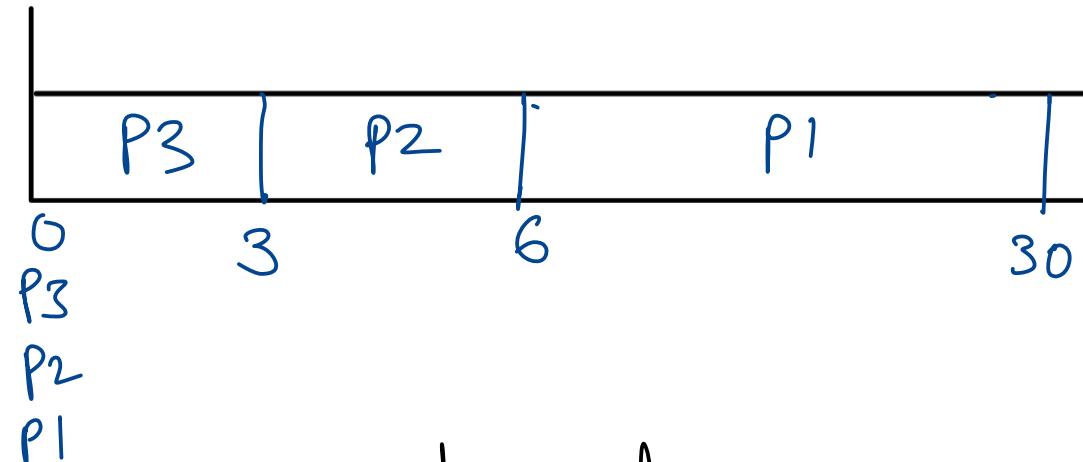
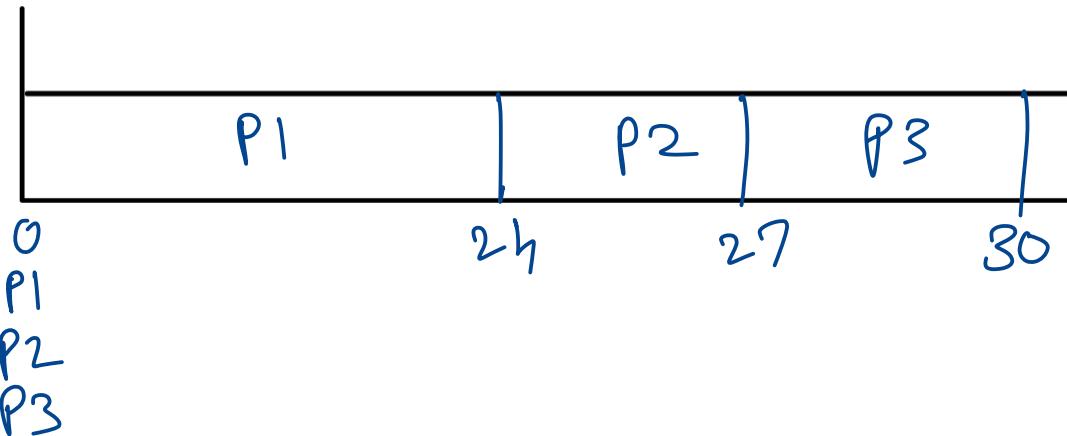
WT RT TAT

0	0	24
24	24	27
27	27	30

Process	Arrival	CPU Burst
P3	0	3
P2	0	3
P1	0	24

WA RT TAT

0	0	3
3	3	6
6	6	30

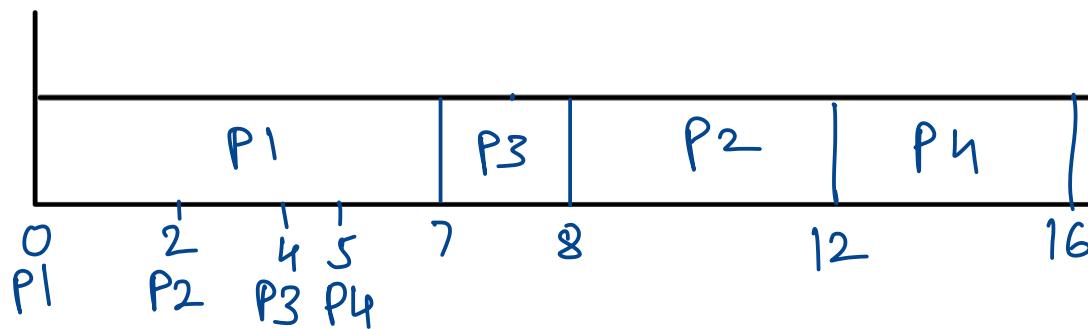


Conway effect : due to arrival of longer processes all other processes need to wait

SJF (Shortest Job First)

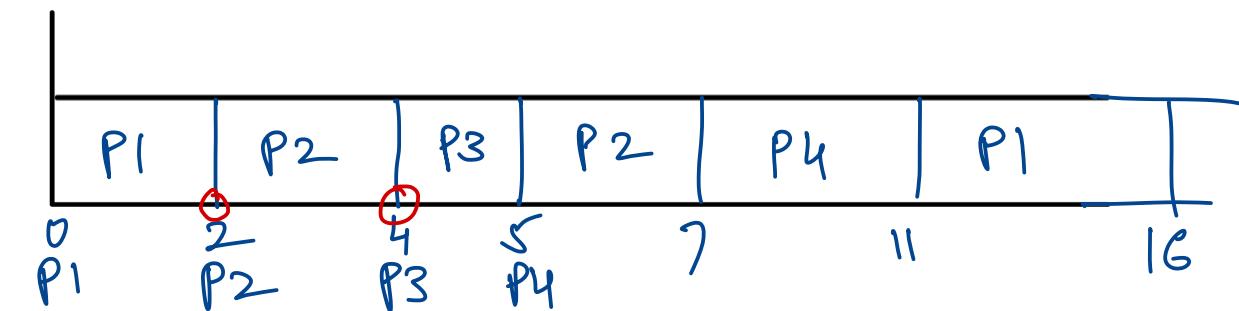
(Non Preemptive)

Process	Arrival	CPU Burst	WT	RT	TAT
P1	0	7	0	0	7
P2	2	4	6	6	10
P3	4	1	3	3	7
P4	5	4	7	7	11



Shortest Remaining Time First
(Preemptive)

Process	Arrival	CPU Burst	remaining time
P1	0	7	5
P2	2	4	2
P3	4	1	0
P4	5	4	

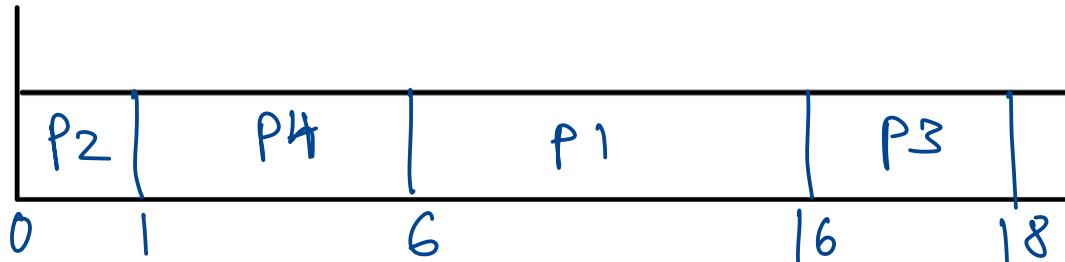


starvation :

due to longer CPU burst processes do not get chance to execute on CPU. (Needs to wait for longer duration)
— there is no solution for starvation in SJF

Priority

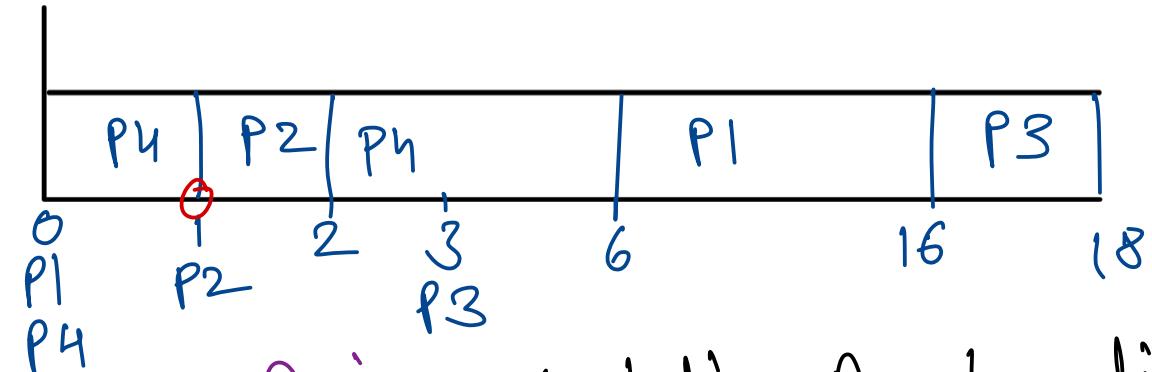
Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	0	1	1 (H)
P3	0	2	4 (L)
P4	0	5	2



Starvation :

- due to less priority, process do not get chance to execute on CPU (need to wait for longer duration)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	1	1	1
P3	3	2	4
P4	0	5	2



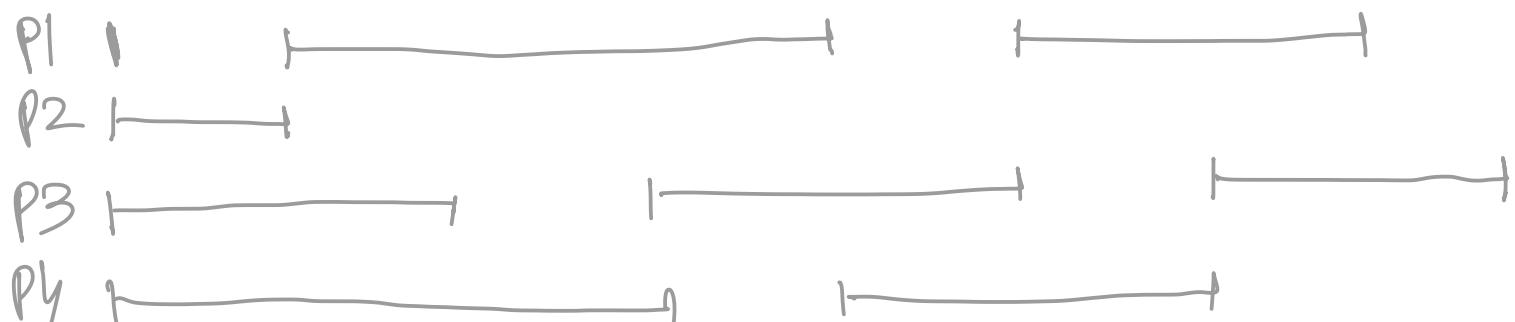
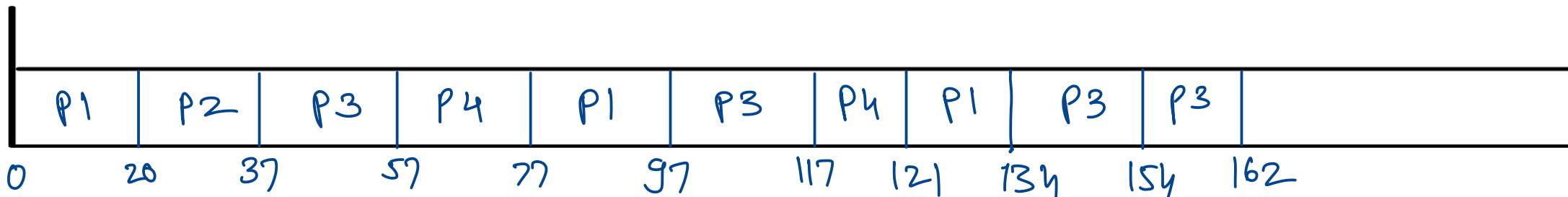
Aging : (solution for starvation)

- priority of starved process is increased gradually, so that it will get scheduled.

RR (Round Robin) (preemptive)

Process	CPU Burst
P1	53
P2	17
P3	68
P4	24

33, 13 X 0 + 57 + 24
 X 20
 48, 28, 8 37 + 40 + 17
 4 X 57 + 40

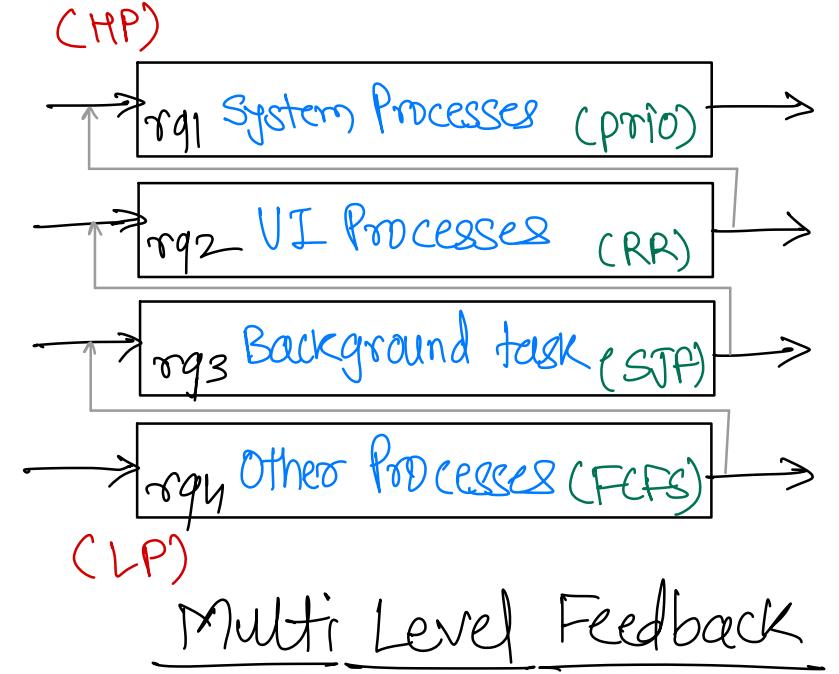
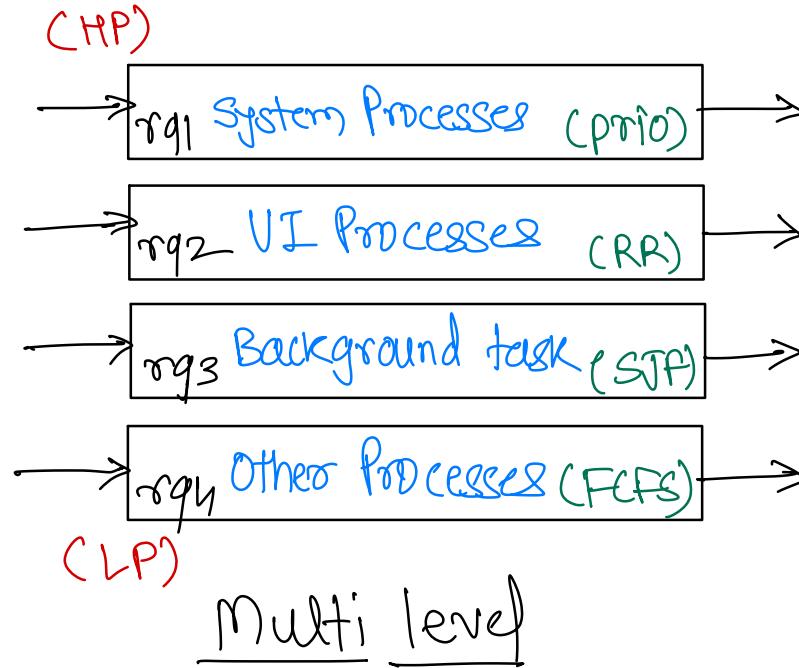


Time Quantum: CPU time slice
 $TQ = 20$

$TQ = 100$
 ↳ behave like FCFS

$TQ = 4$
 ↳ CPU overhead will increase

Multi Level Ready Queue





Thank you!!!

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