

Lecture 1: The Basics of Optimization

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Disclaimer: *These notes aggregate content from several texts and have not been subjected to the usual scrutiny deserved by formal publications. If you find errors, please bring to the notice of the Instructor.*

In this lecture, we discuss some basic optimization problems.

	Variable	Solution space	Solution complexity
Continuous optimization	take any real value (continuous)	infinite	polynomial in the size of the problem
Discrete optimization	discrete	finite	exponential (e.g., knapsack)

1.1 What is optimization?

An optimization problem involves finding the minimum value attained by a function subject to some constraints, i.e.,

$$\min_{x \in \mathcal{C}} f(x),$$

where $f(x)$ is the objective function and \mathcal{C} is the constraint set.

Example. Minimize $(x - 2)^2$ with the constraint that $x \in [0, 1] \cup [4, 7]$.

Here, the objective function is $f(x) = (x - 2)^2$ and the constraint set is $\mathcal{C} = [0, 1] \cup [4, 7]$. This is a one-variable function, and we can easily see from the plot in figure 1.1 that $x^* = 1$ is the optimal value of x .

Example. We want to optimally place a warehouse, so that the sum of the Euclidean distances between the warehouse and the cities is minimized.

Let the cities be located at $\mathbf{y}_1, \dots, \mathbf{y}_m$ and the warehouse at \mathbf{x} . We want to find the following:

$$\min_{\mathbf{x} \in \mathcal{C}} \sum_{i=1}^m \|\mathbf{x} - \mathbf{y}_i\|_2,$$

where \mathcal{C} is the set of all points where we want our warehouse to be, and $\|\cdot\|_2$ represents the L^2 norm, defined by

$$\|\mathbf{x}\| := \sqrt{x_1^2 + x_2^2},$$

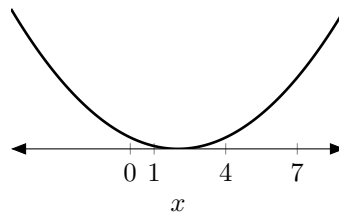


Figure 1.1: A plot of $f(x) = (x - 2)^2$.

where $\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$.

Example (Image de-blurring). We consider grayscale images of size $m \times n$, where each pixel has an intensity value in $[0, 1]$. The input image is $\mathbf{y} = [y_{i,j}]^{m \times n}$ and the desired output is \mathbf{x} .

$$\min_{\mathbf{x} \in [0,1]^{m \times n}} \left(\sum_{i=0}^{m-1} \sum_{j=0}^{n-1} \|y_{i,j} - (k * \mathbf{x})_{i,j}\| + \lambda \sum_{i=0}^{m-1} \sum_{j=0}^{n-1} ((\mathbf{x}_{i,j} - \mathbf{x}_{i,j+1})^2 + (\mathbf{x}_{i+1,j} - \mathbf{x}_{i,j})^2) \right)$$

Here, λ and k are the hyperparameters, which are determined by experiment. The optimal value of λ and k depend on the image. For example, medical images may require different λ and k as compared to images of trees.

Example (Machine learning; curve-fitting). We have inputs (x_i, y_i) , where $i \in [n]$. We want to find

$$\min_{\Theta} \sum_{i=1}^n \ell(h_{\theta}(x_i), y_i),$$

where $\ell(\cdot, \cdot)$ is a loss function, $h_{\Theta}(x) = w_0 + w_1x + w_2x^2$ is the hypothesis (so $h_{\theta}(x_i)$ is the hypothesized point), and $\Theta = (w_0, w_1, w_2)$. An example of a loss function is

$$\ell(y_1, y_2) = (y_1 - y_2)^2.$$

1.1.1 Linear Programming: An optimization type

A linear program is a simple optimization type where the objective function and constraints are given by linear relationships.

Example (Political Winning). Investing money to win the election

Consider a political scenario where a political party P needs to invest money to win elections. There are 3 different demographic classes, Class 1, Class 2, and Class 3, and four issues to be addressed: A, B, C, and D. There is a specific pattern in which people from different Classes respond to different issues. The following table contains the number of votes gained/lost by the political party per unit money spent, on the respective Class and issue:

	Classes		
Issues	Class ₁	Class ₂	Class ₃
$x_1 \rightarrow A$	-2	5	3
$x_2 \rightarrow B$	8	2	-5
$x_3 \rightarrow C$	0	0	10
$x_4 \rightarrow D$	10	0	2
Population	100000	200000	50000
Majority	50000	100000	25000

The population of different classes along with the majority in each Class required by the party to win the elections are also present in the table. The aim of the political party is to minimize the total amount of money it needs to invest, yet get the required majority across each Class. Let x_1, x_2, x_3, x_4 be the amount of money the party invests in issues A, B, C, and D respectively. Hence, this problem becomes the following optimization problem:

We want to minimize $x_1 + x_2 + x_3 + x_4$ and hence, the money spent on the election.

$$-2x_1 + 8x_2 + 0x_3 + 10x_4 \geq 50000 \quad (1.1)$$

$$5x_1 + 2x_2 + 0x_3 + 0x_4 \geq 100000 \quad (1.2)$$

$$3x_1 + 5x_2 + 10x_3 + 2x_4 \geq 25000 \quad (1.3)$$

Here, $x_1, x_2, x_3, x_4 \geq 0$. Let us look at the optimal solution $x_1^*, x_2^*, x_3^*, x_4^*$.

$$x_1^* = \frac{2050000}{111}, \quad x_2^* = \frac{425000}{111}, \quad x_3^* = 0, \quad x_4^* = \frac{625000}{111}.$$

The optimal value of $x_1 + x_2 + x_3 + x_4 = x_1^* + x_2^* + x_3^* + x_4^* = 3100000/111$.

After multiplying and adding the equations as $(1.1)*\frac{25}{222} + (1.2)*\frac{46}{222} + (1.3)*\frac{14}{222}$, we get

$$x_1 + x_2 + \frac{140}{222}x_3 + x_4 \geq \frac{3100000}{111} \quad (1.4)$$

$$\because x_1 + x_2 + x_3 + x_4 \geq x_1 + x_2 + \frac{140}{222}x_3 + x_4 \geq \frac{3100000}{111}$$

\therefore This proves that our given solution is truly the optimal solution

1.1.2 Standard Form of Linear Program

Let \mathbf{x} be the vector containing variables to optimize and \mathbf{c} be the vector of constants

$$\mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \quad \mathbf{c} = \begin{bmatrix} c_1 \\ c_2 \\ \vdots \\ c_n \end{bmatrix}$$

Then we can write the standard form of linear program as follows :

Maximize $\mathbf{c}^T \mathbf{x} = c_1x_1 + c_2x_2 + \dots c_nx_n$

subject to constraints :

$$\mathbf{A}_{m \times n} \mathbf{x}_{n \times 1} \leq \mathbf{b}_{m \times 1} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_m \end{bmatrix}$$

$$\mathbf{x} \geq \mathbf{0}$$

Here \geq means element wise greater than equal to i.e.

$$\begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \geq \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

$$\forall i, \quad x_i \geq 0$$

The aforementioned problem is commonly referred to as the Primal problem, and it is accompanied by a corresponding dual problem.

Primal Problem (P1)

$$\begin{array}{ll} \text{maximize} & \mathbf{c}^T \mathbf{x} \\ \text{subject to} & \mathbf{Ax} \leq \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0} \end{array}$$

Dual Problem (P2)

$$\begin{array}{ll} \text{minimize} & \mathbf{b}^T \mathbf{y} \\ \text{subject to} & \mathbf{A}^T \mathbf{y} \geq \mathbf{c} \\ & \mathbf{y} \geq \mathbf{0} \end{array}$$

1.1.3 Weak Duality Principle

Let \mathbf{x} and \mathbf{y} represent feasible solutions, i.e., solutions that satisfy all the constraints, for the Primal and Dual problems, respectively then

$$\mathbf{b}^T \mathbf{y} \geq \mathbf{c}^T \mathbf{x}$$

Proof :

$$\mathbf{Ax} \leq \mathbf{b}$$

$$\mathbf{x}^T \mathbf{A}^T \leq \mathbf{b}^T$$

As $\mathbf{y} \geq \mathbf{0}$, multiplying it on both sides does not change the inequality

$$\mathbf{x}^T \mathbf{A}^T \mathbf{y} \leq \mathbf{b}^T \mathbf{y}$$

We know $\mathbf{A}^T \mathbf{y} \geq \mathbf{c}$, so

$$\mathbf{b}^T \mathbf{y} \geq \mathbf{x}^T \mathbf{c}, \text{ which can also be written as } \mathbf{b}^T \mathbf{y} \geq \mathbf{c}^T \mathbf{x}.$$

Thus weak duality principle provides relation between the solutions of Primal and Dual problems.

1.1.4 Strong Duality Theorem

If a linear programming problem has an optimal solution, so does its dual. If \mathbf{x}^* and \mathbf{y}^* are the optimal solutions of Primal and Dual problems respectively then

$$\mathbf{b}^T \mathbf{y}^* = \mathbf{c}^T \mathbf{x}^*$$