Assignment 3 Report ::: PINTOS

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Installing Pintos on Ubuntu

:: Refer --> https://web.stanford.edu/class/cs140/projects/pintos/pintos 1.html

- # Understanding how Pintos program flow works::
- 1.1 Understanding GDB debugger to track programs:
- :: Refer --> http://math.hws.edu/eck/cs431/f16/lab4/index.html

(GDB)- We will need to work in two terminals.

- In the first terminal execute the pintos --run function:
- Eg: pintos --gdb -v -- run alarm_multiple
- In the second terminal make sure you're in the ~/src/threads/build/ directory
- Execute:: pintos-gdb kernel.o
- To use GDB always use the second terminal now::
- To debug execute :: debugpintos
- Some instances of GDB are as follows::
 - "break function_name" to apply the breakpoint
 - "continue" or "c" { to continue till one loop}
 - "next" or "n" for the next line
 - "p variable_name" to print the variable

Detailed explanation is given in the link mentioned...

- -- Understand makefiles -> see all the makefiles in the directories and understand the flow of the programs.
- -- For changes try changing /src/utils/Makefile . Also
- -- Some tips: you may change the CC to make compatibility accordingly set the FLAGS.

- * Makefile clean operation here comes in handy.
- -- It clears all the unnecessary .o files created while testing.

```
## Test cases for part 2 : preemption of threads
```

-----How to run a program for checking the outputs?

Note:: Path to all the testcases we need :: ~/pintos/src/threads/build/tests/threads/*.

To test them individually we can use ::

```
--->> $ make
```

~/pintos/src/threads/build/tests/threads/tests/threads/test-name-from-list.result

Moreover, understanding the flow of program and the structure of the list_elem was the trickiest.

- --Refer list elem structure in **list.h** and **list.c** for functions {/src/lib/kernel/list.h}
- --Refer thread structure in threads.h and thread.c for functions {/src/threads/threads.h}
- --Refer list entry() macro from **list.c** ---- the hardest and the most generic function of the OS.

```
#define list_entry(LIST_ELEM, STRUCT, MEMBER)

((STRUCT *) ((uint8_t *) &(LIST_ELEM)->next \
- offsetof (STRUCT, MEMBER.next)))
```

- --Implementation of List.c and use of functions like list insert ordered(...) {} .
- --Know and understand what thread_yield(), schedule() does. {}
- --Understand the drawback of the function such as "list_pushback(...)" over list_insert_ordered(...)

Structures of Data_Types:

```
struct list_elem
{
   struct list_elem *prev; /* Previous list element. */
   struct list_elem *next; /* Next list element. */
};
```

```
struct list
  struct list_elem head; /* List head. */
  struct list_elem tail; /* List tail. */
};
struct thread
 {
  tid_t tid;
                          /* Thread identifier. */
  enum thread_status status;
                                   /* Thread state. */
  char name[16];
                             /* Name (for debugging purposes). */
                            /* Saved stack pointer. */
  uint8_t *stack;
  int priority;
                          /* Priority. */
  struct list_elem allelem;
                                /* List element for all threads list. */
  /* Shared between thread.c and synch.c. */
  struct list_elem elem;
                              /* List element. */
#ifdef USERPROG
  /* Owned by userprog/process.c. */
                              /* Page directory. */
  uint32_t *pagedir;
#endif
  /* Owned by thread.c. */
                              /* Detects stack overflow. */
  unsigned magic;
  //Wakeup_ticks for inserting in wakeup queue
 // Added new
  int wakeup_ticks;
 };
```


Test cases for part 2 : preemption of threads

1.1 Alarm-single and Alarm-zero

--->> \$ make \sim /pintos/src/threads/build/tests/threads/tests/threads/alarm-single.result

It passes without making any change. It gives a brief idea of how to check for the program flow in Pintos.

We used GDB and putting several breakpoints on functions. 1.2 Alarm-multiple, alarm-negative, alarm-simultaneous -Understand what these cases do. -Refer to the output expected from their corresponding test-case.ck file ---> Refer {/src/devices/timer.c} to understand what it does, - the function timer_sleep() is called. - Most important is to understand what busy_wait means in this context and how is it different from busy_wait from what we refer in semaphore (Galvin) - We have to make changes in timer sleep(). -How? - In {./thread.h} structure we insert an element to keep track of the wakeup ticks. - Refer what timer_ticks(), ticks variables are used for... - Insted of using the busy_wait while loop for thread yield, we use the added extra variable in thread structure (int wakeup_ticks) to keep track of when to wakeup the thread. Every thread is initialized with wakeup_ticks = timer_ticks() + ticks; - Use this to insert list in ordered fashion using -- list_insert_ordered(...) . - For the comparator function write your own. { first get to observe what this function actually does. Now the 3rd argument is comparator function } - The function compares the wakeup_ticks someway like this---int insert_sleep_at (struct list_elem *x , struct list_elem *y , void* aux UNUSED) {

struct thread *first = list entry(x , struct thread , elem); ----- list entry usage

struct thread *last = list_entry(y , struct thread , elem);

```
if(first->wakeup_ticks < last->wakeup_ticks) ---- using wakeup_ticks to push_back the threads
in a sorted order...
  return 1:
 return 0;
## Test cases for part 3: priority scheduling
1.Alarm-priority
→ $ make ~/pintos/src/threads/build/tests/threads/tests/threads/alarm-priority.result
- use the testing command as mentioned above to look at the way it's meant to be implemented.
- we make changes in /src/threads/thread.c
- a similar helper functions as we used in inserting thread according to wakeup time is (we can
observe this by looking at the desired output)
- in thread_unblock() use this function to pushback according to the thread priority.
- Here the function list_insert_ordered is useful..
***********************
int priority_accordance ( struct list_elem* x , struct list_elem* y , void* aux UNUSED)
 struct thread *first = list_entry( x , struct thread , elem );
 struct thread *last = list_entry( y , struct thread , elem );
 if(first->priority > last->priority)
  return 1:
 return 0;
```

}

2. Priority-change

- → \$ make ~/pintos/src/threads/build/tests/threads/tests/threads/alarm-change.result
- use the testing command as mentioned above to look at the way it's meant to be implemented.
- we make changes in /src/threads/thread.c
- the main observation is the usage of the function "thread_set_priority(new_priority)"
- it dynamically changes the priority~
- after "creating thread" check if ready list changed and effects preemption of currently running thread
- so in thread create use a function which does so.
- we implemented void priority_ordering();
- checking if current thread priority is less than front of ready list(max priority thread).
- if so then we yield the thread.

```
****
void priority_ordering()
{

struct thread *t = thread_current();
enum intr_level intt = intr_disable ();

if(!list_empty(&ready_list))
{
    struct thread *y = list_entry( list_front(&ready_list) , struct thread , elem );

if(t->priority <= y->priority)
{
    thread_yield();
}
}
intr_set_level(intt);
```

3. 4. Priority-fifo and Priority-sema

- → \$ make ~/pintos/src/threads/build/tests/threads/tests/threads/priority-fifi.result
- → \$ make ~/pintos/src/threads/build/tests/threads/tests/threads/priority-sema.result
- use the testing command as mentioned above to look at the way it's meant to be implemented.
- we make changes in /src/threads/sync.c and sync.h
- Understand the modules of semaphores. The struct semaphore, struct lock and struct condition.
- Try to analyze how semaphores and locks are implemented practically, (we have read the theory)
- notice we have a concept like counting semaphore implemented in priority-sema...
- **sync.c** is the major function we tries to debug and atlast made some additions there.
- notice the function sema down(...)
- here for priority-sema and fifo just one addition is enough.
- use the comparator function used in thread to insert the elements in the list based on threads priority.
- add the thread.c header in #include of the code to refer to the function... {}
- use the list_insert_ordered using the comparator function priority_accordance(...) mentioned above!
- after this check if the thread needs to be replaced with the higher_priority thread...
- this can be done using thread_block()
- another way is to use thread_yield().

5. Priority-condvar

- → \$ make ~/pintos/src/threads/build/tests/threads/tests/threads/priority-condvar.result
- understand the structure of struct semaphore_elem and how this test is different from priority-sema

- this uses Binary semaphore instead of Counting semaphore. - the function in sync.c {cond_wait()} is used here. - function cond signal() working understanding is important - main observation is to see that the list in struct condition; --> waiters is needed to be sorted. - the list contains entry of structure - > semaphore_elem *************************** struct semaphore elem int priority_semaphore; // Added to the structure struct list_elem elem; /* List element. */ struct semaphore semaphore; /* This semaphore. */ **}**; - instead of sorting we use an indirect approach. This exploits and reuses the entire code - in the structure semaphore_elem; { as shown above } - to keep track of priority of track we insert an element { int priority semaphore } - We assign it with the current thread's priority; in the function void cond_wait(...) waiter.priority_semaphore = thread_current()->priority; ******************************* int priority_sema (struct list_elem* x, struct list_elem* y,void* aux) struct semaphore_elem *f = list_entry(x,struct semaphore_elem,elem); struct semaphore_elem *I = list_entry(y,struct semaphore_elem,elem); if(f->priority_semaphore > I->priority_semaphore) return 1: return 0;

- we used this comparator function to insert list ordered fashion in the function { void codn_wait() }
MLFQS
read theory