

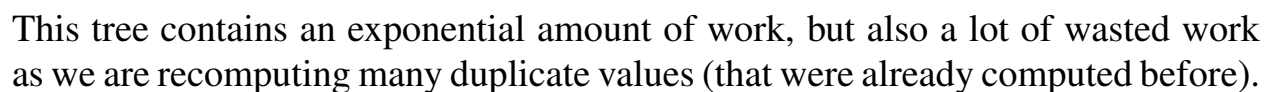
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Dynamic Programming

A naive algorithm for this could be given by the following algorithm:

For example, let us compute the recursive tree for $Fib(7)$:



For example, we are completely recomputing the entire recursive tree for $Fib(5)$ on the left most node on level 3 when it is already being computed as the right most node on level 2.

A better approach is to use *dynamic programming* where we start with an empty list and build up the answer from there (being able to remember the computation for smaller subproblems).

In the following algorithm, we initialize a list and include the two base cases in the first two entries (first index is 1). There are no recursive calls, but we are still exploiting the recursive structure of the problem.

```
1 Function  $Fib(n)$  :  
2    $T \leftarrow []$   
3    $T[1] \leftarrow 1$   
4    $T[2] \leftarrow 1$   
5   for  $i \in 2 \rightarrow n$  do  
6      $T[i] \leftarrow T[i - 1] + T[i - 2]$   
7   return  $T[n]$ 
```

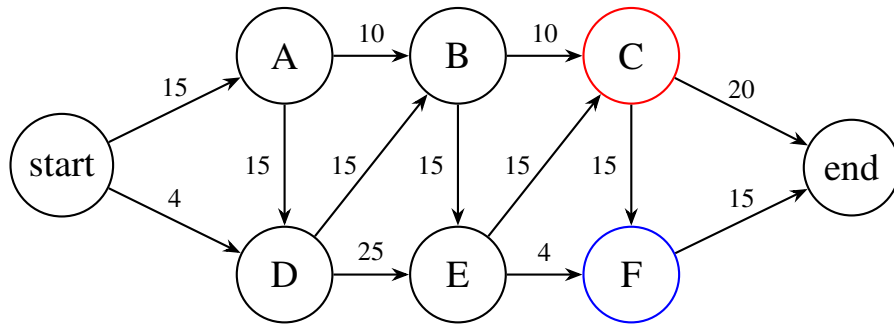
In general, dynamic programming is great because

- Exploiting recursive structure without recursion
- Solve the smaller problems first and store the solutions to avoid recalculation

Let's see another example on the next page!

Route Planning Example

Let's say you have a start and end point in a grid, and you are curious about how long it takes to get from the start to the end.



If we pick the two nodes (colored red and blue), then the time that it takes to get to the end can be defined recursively (T is the time it takes to get to a certain node):

$$T(\text{end}) = \min(T(C) + 20, T(F) + 15)$$

A recursive approach this would have a lot of wasted computation (recomputing the distance to nodes we already have a stored value for). Rather, we opt for a dynamic programming approach.

Idea: Since we need all the subproblems, calculate the subproblems first and then use those to propagate up the problem (larger and larger problems).

First, we know that $T(\text{start}) = 0$, then $T(A) = 15$ and $T(D) = \min(T(A) + 15, 4) = 4$. Repeating this logic gives us:

$$T(\text{start}) = 0$$

$$T(A) = T(\text{start}) + 15 = 15$$

$$T(D) = \min(T(\text{start}) + 4, T(A) + 15) = \min(4, 19) = 4$$

$$T(B) = \min(T(A) + 10, T(D) + 15) = \min(25, 19) = 19$$

$$T(E) = \min(T(B) + 15, T(D) + 25) = \min(34, 29) = 29$$

$$T(C) = \min(T(B) + 10, T(E) + 15) = \min(29, 44) = 29$$

$$T(F) = \min(T(E) + 4, T(C) + 15) = \min(33, 44) = 33$$

$$T(\text{end}) = \min(T(C) + 20, T(F) + 15) = \min(49, 48) = \boxed{48}$$

To formalize this example into an algorithm: the shortest path in a DAG (Directed Acyclic Graph) is:

```

1 Function ShortestPath( $G, u$ ) :
    /* topologically sort the DAG, covered in Unit 3 of this course */
2    $G_T \leftarrow \text{TopoSort}(G)$ ;           // graph is now "ordered"
3    $T \leftarrow \{\}$ ;                       // dynamic programming map <edge: time>
4    $T[s] \leftarrow 0$ ;                       // takes no time to get to start
5   for edge( $x, v$ )  $\in G$  do
    /* checking if there is a faster way get to  $v$  */
6   |    $T[v] \leftarrow \min(T[v], T[x] + \text{weight}(x, v))$ 
7   return  $T[v]$ 

```

The running time for the TopoSort is $O(n^2)$ where n is the number of vertices, and there can be as many as $|V|^2 = n^2$ edges, so the loop iterates a total of n^2 times gives $O(n^2)$ total runtime.

Dynamically Robbing Houses Example

You have houses h_1, \dots, h_n where h_i is the amount of the i^{th} house. You are a robber and cannot rob consecutive houses. What is the maximum wealth you can accrue?

For every dynamic programming problem, there are 4 parts for a total answer:

1. Table entry definition
2. Recursive relationship between entries of the table
3. Pseudocode (for the algorithm, this unit only)
4. Running time analysis

Let $T[i]$ be the maximum wealth from houses m_1, \dots, m_i (this is the table entry definition). Now, we want to know how to define $T[i]$ in terms of smaller i .

We could choose to rob the i^{th} house, but then we could not rob the previous house, so our sum would be $T[i] = h_i + T[i - 2]$. However, we could also choose not

to rob this house, so $T[i] = T[i - 1]$ and since we want to maximize how much money to make, the recursive relationship between entries would be defined as:

$$T[i] = \max(h_i + T[i - 2], T[i - 1])$$

For pseudocode, we can implement this as:

```

1 Function Robbery( $h_1, \dots, h_n$ ) :
    /* let  $T$  be an empty array that is  $n$  in length */
2      $T \leftarrow []$ 
3      $T[1] \leftarrow h_1$   $T[2] \leftarrow \max(h_1, h_2)$ 
    /* this loop is inclusive on  $n$  */
4     for  $i \in 3 \rightarrow n$  do
5          $T[i] \leftarrow \max(h_i + T[i - 2], T[i - 1])$ 
6     return  $T[n]$ 

```

This algorithm has a total runtime given by $O(n)$ because all operations are $O(1)$ while the loop is iterating $O(n)$ times meaning we are making $O(n)$ $O(1)$ operations, giving $O(n)$ final runtime.

Challenge: Hopscotch Example

Let's play a game of hopscotch! There is a straight line of squares with integers values on each square. We start at the start and eventually reach the end.

The score is the sum of the squares that you land on. *However*, you can only jump forward by 1 or 3 steps. What is the maximum score?

Try this question. Answer is on the next page.

Answer

We can represent $T[i]$ has the maximum score for the squares s_1, \dots, s_i , then we can give a recurrence relation as:

$$T[i] = \max(T[i-1], T[i-3]) + s_i$$

In pseudocode, it could be given as:

```
1 Function Hopscotch( $s_1, \dots, s_n$ ):  
    /* let  $T$  be an empty array that is  $n$  in length */  
2     $T \leftarrow []$   
3     $T[1] \leftarrow s_1$   
4     $T[2] \leftarrow s_1 + s_2$   
5     $T[3] \leftarrow s_1 + s_2 + s_3$   
6    for  $i \in 4 \rightarrow n$  do  
7         $T[i] = \max(T[i-1] + s_i, T[i-3] + s_i)$   
8    return  $T[n]$ 
```

This algorithm has a runtime of $O(n)$ since the initialization of the array takes $O(n)$ time (for n -long array) and the base case updates are all $O(1)$. The for loop iterates a total of $O(n)$ times giving a final runtime of $O(n)$.