

# Student Answer Script View



MIT MPL - BTech-M Sc-MCA - II-IV and VI Semester - Midterm Examination - Mar 2025      Answer Sheet

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<b>Course:</b>	MIT-OPE
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<b>Subject Name:</b>	APPLIED GRAPH THEORY
<b>Exam Date:</b>	08-Mar-2025

Q.No : 1)

If  $(4, 3, 2, 2, x)$  is a degree sequence of a simple graph, then the value of  $x$  is

0

1

2

3

Q.No : 2)

In a social network, if the sum of degrees of all users is 100, how many connections (edges) are there in the network?

50

52

54

56

Q.No : 3)

If  $G$  is a disconnected graph of order 101, then the minimum degree of  $G$  is

less than 50

less than 51

less than 52

less than 53

Q.No : 4)

In a self-complementary graph with 101 vertices, the number of edges is

2525

2526

2527

2528

Q.No : 5)

The number of cut vertices in a Hamiltonian graph of order 10 is

0

2

4

6

Q.No : 6)

The vertex independence number of a path graph with 15 vertices is

8

6

4

2

Q.No : 7)

The vertex covering number of a cycle with 30 vertices is

14

15

16

17



Q.No : 8)

The number of cut edges in a tree with 24 vertices is

22

23

24

25

Q.No : 9)

If  $G$  is a connected graph with diameter 10, then diameter of the complement graph of  $G$  is at most

4

3

5

6

Q.No : 10)

Maximum number of edges in a bipartite graph G with 14 vertices is

49

51

54

58

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**Q.No : 11)**

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(i) Prove that a connected graph on  $n$  vertices with  $n - 1$  edges is a tree.

(ii) A tree  $T$  with 50 pendant vertices has an equal number of vertices of degree 2, 3, 4 and 5 and no vertices of degree greater than 5. Determine the order of  $T$ .

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(ii) Let no of vertices of degree 2, 3, 4, 5  
 $x$ .

Let total no of vertices be  $n$ .

$$\Rightarrow n = 50 + x + x + x + x$$

$$n = 50 + 4x \quad \text{————— (1)}$$

$$\sum_{i=1}^n \deg(v_i) = 2(n-1)$$

$$50 + 2x + 3x + 4x + 5x = 2(n-1)$$

$$50 + 14x = 2n - 2$$

$$52 + 14x = 2n$$

$$26 + 7x = n \quad \text{————— (2)}$$

From (1) and (2),

$$50 + 4x = 26 + 7x$$

$$24 = 3x$$

$$\boxed{x = 8}$$

$$\Rightarrow n = 50 + 4(8)$$

$$\boxed{n = 82}$$

$$\begin{aligned} \text{order of } T &= 50 + 14x \\ &= \underline{\underline{162}} \end{aligned}$$



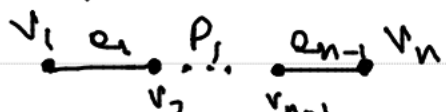




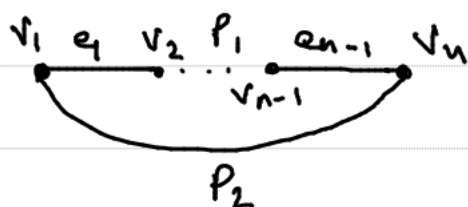
(i) A connected graph on  $n$  vertices is a tree if there are no cycles present in the graph.

A graph with  $n$  vertices needs at least  $n-1$  edges to make a cycle.

Let  $e_1$  and  $e_{n-1}$  be two edges with vertices  $v_1$  and  $v_n$  on the path  $P_1$  between  $v_2$  and  $v_{n-1}$ .

 It has  $n-1$  edges and no cycles.

Let  $P_2$  be another path between  $v_1$  and  $v_n$ .



$v_1 \xrightarrow{e_1} v_2 \xrightarrow{P_1} v_{n-1} \xrightarrow{e_{n-1}} v_n \xrightarrow{P_2} v_1$  makes a cycle with  $n$  edges.

This is a contradiction.

Hence a connected graph on  $n$  vertices with  $n-1$  edges is a tree.



$$\text{Vector } v = (3, 3, 2, 2, 2, 2, 2)$$

$$\textcircled{1} \quad v = (\overset{v_1}{\underline{3}}, \overset{v_2}{3}, \overset{v_3}{2}, \overset{v_4}{2}, \overset{v_5}{2}, \overset{v_6}{2}, \overset{v_7}{2})$$

$$v_1 = (\overset{v_2}{2}, \overset{v_3}{1}, \overset{v_4}{1}, \overset{v_5}{2}, \overset{v_6}{2}, \overset{v_7}{2}) \longrightarrow v_1$$

$$\textcircled{2} \quad v = (\overset{v_2}{2}, \overset{v_5}{2}, \overset{v_6}{2}, \overset{v_7}{2}, \overset{v_3}{1}, \overset{v_4}{1})$$

$$v_1 = (\overset{v_5}{1}, \overset{v_6}{1}, \overset{v_7}{2}, \overset{v_3}{1}, \overset{v_4}{1}) \longrightarrow v_2$$

$$\textcircled{3} \quad v = (\overset{v_7}{2}, \overset{v_5}{1}, \overset{v_6}{1}, \overset{v_3}{1}, \overset{v_4}{1})$$

$$v_1 = (\overset{v_5}{0}, \overset{v_6}{0}, \overset{v_3}{1}, \overset{v_4}{1}) \longrightarrow v_7$$

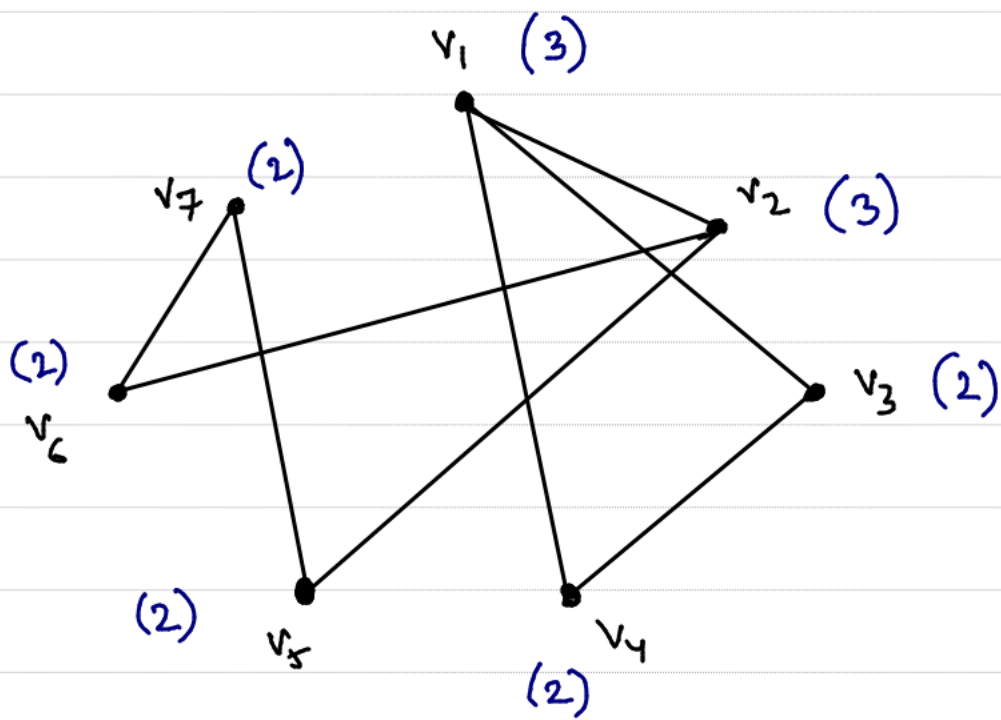
$$\textcircled{4} \quad v = (\overset{v_3}{1}, \overset{v_4}{1}, \overset{v_5}{0}, \overset{v_6}{0})$$

$$v_1 = (\overset{v_4}{0}, \overset{v_5}{0}, \overset{v_6}{0}) \longrightarrow v_3$$

Since all are 0, the vector is graphical

Graph:



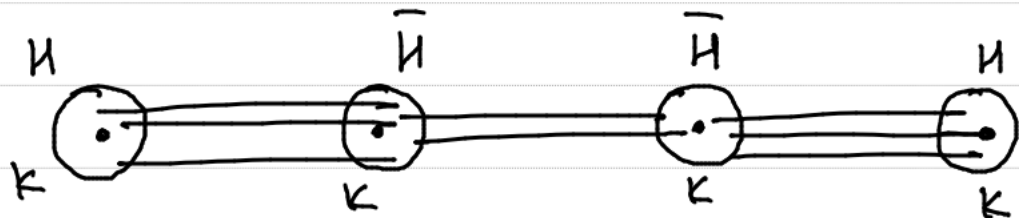




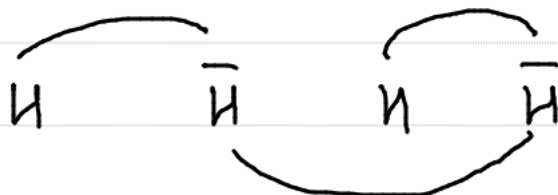
A self-complementary graph can only be of order  $4k$  or  $4k+1$  for some integer  $k$ .

If order  $= 4k$ ,

Let  $H$  be a graph and  $\bar{H}$  be its complementary graph.

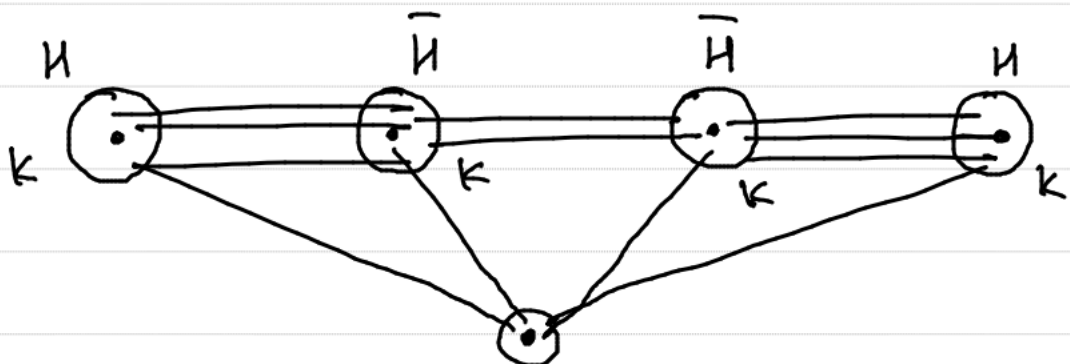


Thus, the graph is



If order  $= 4k+1$ ,

Let  $H$  be a graph and  $\bar{H}$  be its complementary graph.



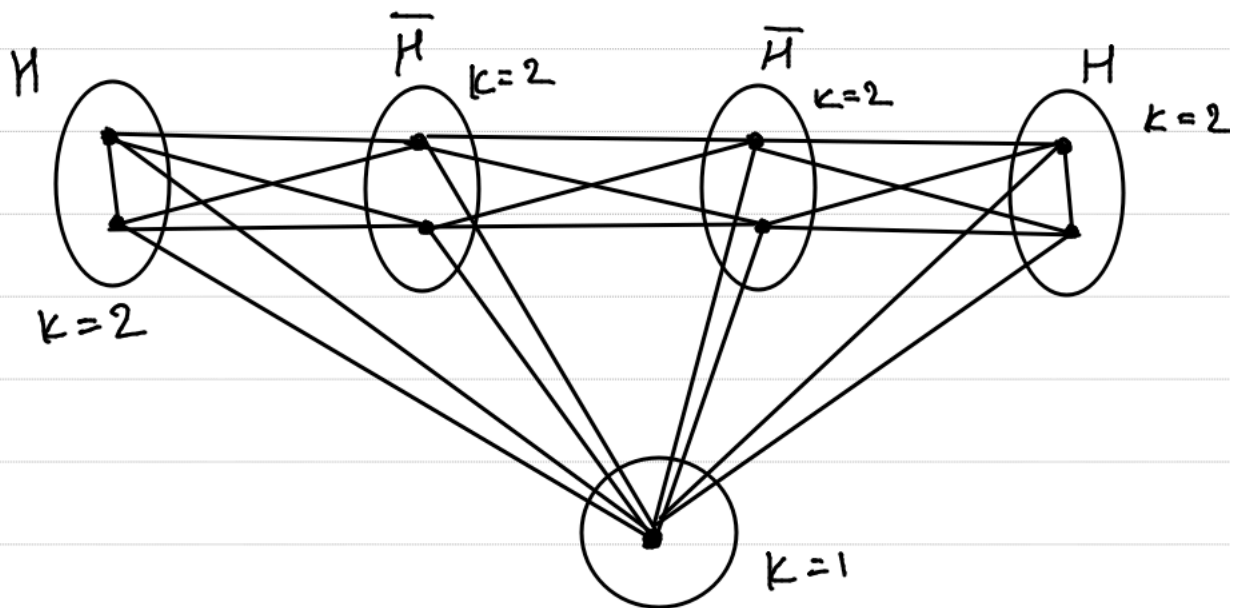




For graphs of order  $4k+2$  and  $4k+3$ , a self-complementary graph cannot be drawn

$$9 \text{ vertices} = 4k+1$$

$$k=2$$





**Q.No : 14)**

If  $G$  is Eulerian, then show that all its vertices have even degree. Give an example of a Hamiltonian graph that is not Eulerian.



Eulerian graph is when a graph contains an Eulerian cycle.

Hamiltonian graph is when a spanning set the graph exists.

→ If any vertex has an odd degree, it is have a cycle.

⇒ It won't have an ~~Eulerian~~ cycle.

⇒ It is a contradiction.

Hence, all vertices must



**Q.No : 15)**

Let  $G$  be graph on  $n$  vertices, with vertex covering number  $\alpha_0$  and vertex independence number  $\beta_0$ . Prove that  $\alpha_0 + \beta_0 = n$ .





Vertex covering number  $\alpha_0$  is the minimum cardinality of a vertex covering set. A vertex covering set when all the edges are connected to at least one of the vertices.

Vertex independence number  $\beta_0$  is the maximum cardinality of a vertex independence set. A vertex independence set when no two vertices in the set are adjacent to each other.

Using an example:



vertex cover =  $\{v_2, v_4\}$   
 $\alpha_0 = 2$

Vertex independence set =  $\{v_1, v_3, v_5\}$   
 $\beta_0 = 3$

$$n = 5$$

$$\alpha_0 + \beta_0 = 5$$

Hence,  $\alpha_0 + \beta_0 = n$



**Q.No : 16)**

Prove that the chromatic polynomial  $P_{C_n}(x)$  of a cycle  $C_n$  is  $(x-1)^n + (-1)^n(x-1)$  for  $n \geq 3$ .



For an even cycle, no of colours used =  
 For an odd cycle, no of colours used = 3

Even cycle,  $n = \text{even}$

$v_1$  would get colour 1.

$\Rightarrow v_2$  would get colour 2.

$\Rightarrow v_3$  would get colour 1.

Similarly,  $v_n$  would get colour 2.

$v_1$  and  $v_n$  are adjacent but have different

Odd cycle,  $n = \text{odd}$

$v_1$  would get colour 1.

$\Rightarrow v_2$  would get colour 2.

$\Rightarrow v_3$  would get colour 1.

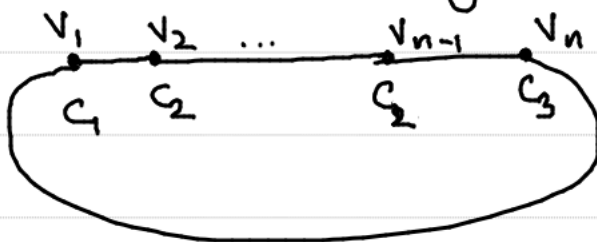
Similarly,  $v_{n-1}$  would get colour 2.

Similarly,  $v_n$  would get colour 1.

$v_1$  and  $v_n$  are adjacent but have same col

This would be a contradiction.

Hence we assign a third colour to  $v_n$ .



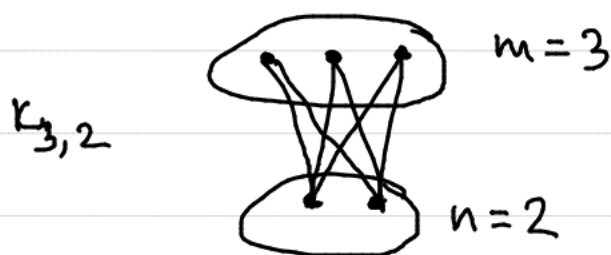


<b>Q.No : 17)</b>
Show that a bipartite graph $G$ has no odd cycles.





A bipartite graph is one where all the  $v$  can be divided into two sets where the vertices in each set are not adjacent to other vertices in the same set.

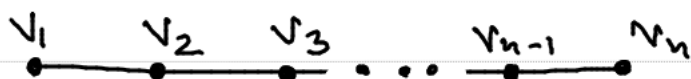


If a cycle graph has odd cycles,

$\Rightarrow$  no of vertices = odd =  $n$

$\Rightarrow$  no of edges =  $n$

Let  $G$  be a graph with vertices  $v_1, v_2$   
Let the two bipartite sets be  $S_1, S_2$



Let  $v_1$  belongs to  $S_1$ .

$\Rightarrow v_2$  belongs to  $S_2$ .

$\Rightarrow v_3$  belongs to  $S_1$ .

Similarly,  $v_{2n} \in S_2$  and  $v_{2n+1} \in S_1$



If  $n = \text{odd}$ ,

$v_n$  would belong to  $S_1$

$v_1$  and  $v_n$  are adjacent and  $v_1$  is also in  $S_1$ .

Thus,

$$S_1 = \{v_1, v_3, \dots, v_n\}$$

$$S_2 = \{v_2, v_4, v_6, \dots, v_{n-1}\}$$

Thus,  $v_1$  and  $v_n$  would form a cycle since they are in the same set.

This is a contradiction.

Hence, a bipartite graph cannot have odd cycles.



**Q.No : 18)**

Draw a simple graph with vertex connectivity 2, edge connectivity 3 and minimum degree 4.



Vertex connectivity =  $\kappa(G) = 2 = \kappa_a$

Edge connectivity =  $\lambda(G) = 3 = \kappa_b$

Min degree =  $\delta(G) = 4 = \kappa_d$

Graph:

$\kappa$

|||

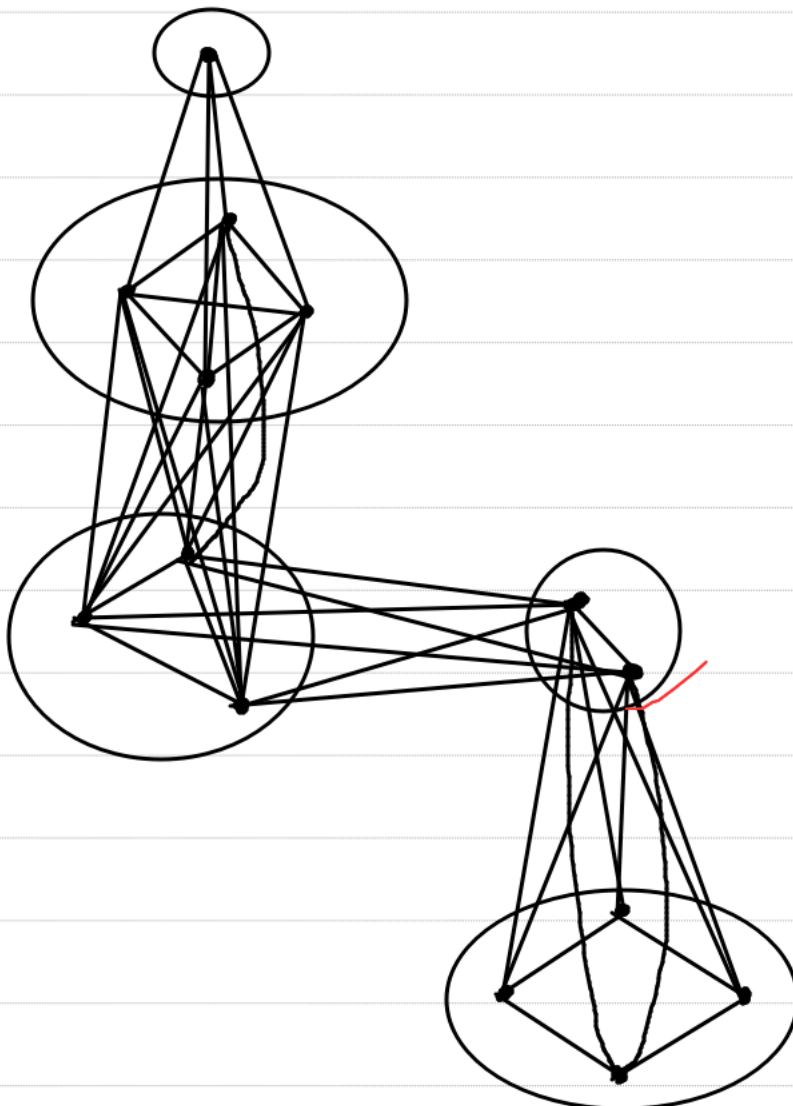
$\kappa_d$

|||

$\kappa_b \equiv \kappa_a$

|||

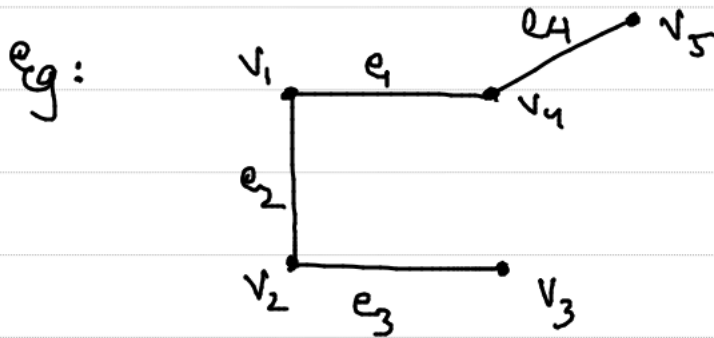
$\kappa_d$





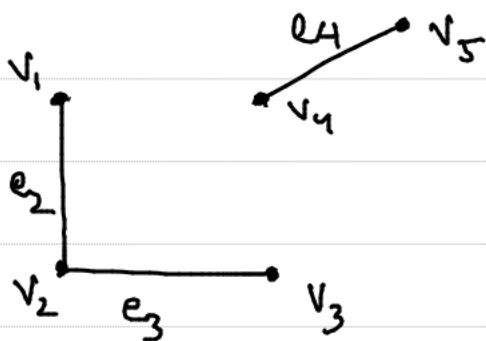


A cut edge in a graph is any edge w<sup>h</sup> on deleting increases the number of components in the graph.

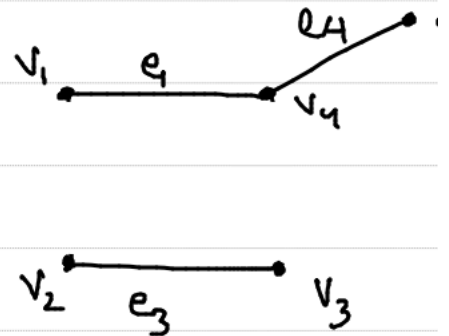


Here either  $e_1$  or  $e_2$  is a cut edge.

cut edge =  $e_1$



cut edge =  $e_2$



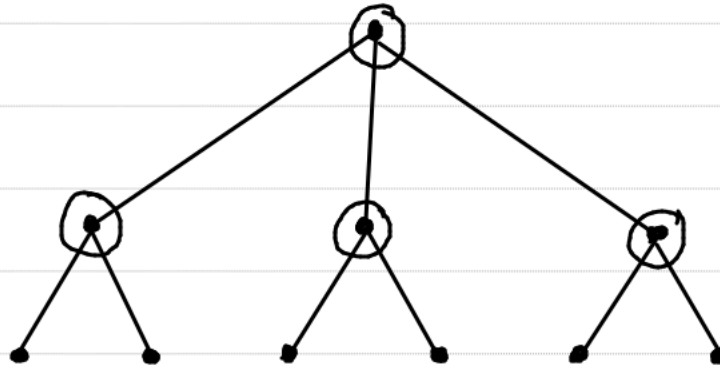
No of vertices = 10

No of pendant vertices = 6

Tree :







There are 4 cut vertices.

Since there are 6 pendant vertices, they can't be cut vertices.

Hence no of cut vertices =  $10 - 6 = 4$

