

CS/IS F214 Logic in Computer Science

MODULE: PREDICATE LOGIC

Expressing using Predicates: Examples

Experessing using Predicates: Horn-Clause Form and Prolog

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Specifications using Predicate Logic

- Examples:
 - No student attended every lecture
 - ∀X ¬ (∀Y student(X) ∧ lecture(Y) --> attended(X,Y))
 - Exercises:
 - Rewrite this (without changing the meaning)
 - i. by replacing the <u>outermost quantifier</u>
 with an <u>existential quantifier</u>.
 - ii. by replacing the <u>innermost quantifier</u> with an <u>existential quantifier</u>.
 - iii. such that the <u>scope of the inner</u> quantifier is the smallest



Specifications using Predicate Logic

- Examples:
 - i. No student attended every lecture
 - $\forall X \neg (\forall Y \text{ student}(X) \land \text{lecture}(Y) --> \text{attended}(X,Y))$
 - ii. No lecture was attended by every student
 - $\forall Y \neg (\forall X \text{ student}(X) \land \text{lecture}(Y) \rightarrow \text{attended}(X,Y))$
 - Exercises:
 - Does (i) imply (ii)?
 - ii. Does (ii) imply (i)?
 - **iii.** If you swap variables X and Y in the quantifier positions will the formula in (ii) remain the same?
 - **IV.** If you swap all occurrences of variables X and Y will the formula in (ii) remain the same?



Exercises (from the text book: Huth & Ryan).

- Write the following in Predicate Logic
 - **i.** All red things are in the box
 - ∀T red(T) --> inBox(T)
 - *ii.* Only red things are in the box
 - ∀T inBox(T) --> red(T)



Exercises (adapted from the text book: Huth & Ryan).

- Write the following in Predicate Logic
 - Raj is Jaya's cousin
- Given a cousin relation defined, this would do:
 - cousin("Raj","Jaya")
- Otherwise one has to define a cousin relation.



Exercises (adapted from the text book: Huth & Ryan).

- Define a cousin property using father, mother, sister, brother:
 - ∀Any1 ∀C cousin(Any1,C) <--
 ∃Pa ∃Pc parent(Any1,Pa) ∧ parent(C,Pc) ∧ sibling(Pa,Pc)
 - $\forall A \forall B \text{ sibling}(A,B) < -- \text{ brother}(A,B) \lor \text{ sister}(A,B)$
 - ∀Any1 ∀Pa parent(Any1,Pa) <-- mother(A,Pa) ∨ father(A,Pa)



Exercises (from the text book: Huth & Ryan).

- Write the following in Predicate Logic
 - All brothers are siblings
- Given a predicate brother(X,Y) and sibling(X,Y) to indicate
 X's brother is Y and X's sibling is Y,
 - $\forall X \ \forall Y \ brother(X,Y) \longrightarrow sibling(X,Y)$.



Exercises (from the text book: Huth & Ryan).

- Write the following in Predicate Logic
 - An attacker can persuade a server that a successful login has occurred even if it hasn't.
- Given the predicates
 - attacker(X) // X is an attacker
 - server(S) // S is a server
 - persuade(A, X, Y) // A can persuade X about Y
 - login(S, token(S, T))
 - // attempt to login into S results in token(S,T) where T is TRUE or FALSE
- the statement above can be encoded as:
 - ∀S ∀A attacker(A) ∧ server(S) -->
 (∀T login(S,token(S,T)) --> persuade(A,S,token(S,TRUE)))



Encoding in Predicate Logic: Example

- Axioms of Group (G,+):
 - Closure:
 - Associativity:
 - Existence of Identity
 - Existence of Inverse



Encoding in Predicate Logic: Example

- Definition of natural numbers (incomplete):
 - $\forall X \text{ (equals(X,0)} \lor \exists Y \text{ equals(X, succ(Y)) --> natural(X)}$
- Note: We must insist Y to be a natural number this will require a recursive definition. End of Note
 - Exercise: Write such a recursive definition.



Prolog Programming and Horn Clauses

- Prolog uses Horn Clauses as the basis for programming:
 - A Horn Clause is an implication with zero or more <u>antecedents</u> and one <u>implicand</u>:
 - All antecedents and the implicand are predicates.
 - i.e. a Horn Clause is of the form:
 - $p_1(T_{11},...,T_{1K1}) \wedge p_2(T_{21},...,T_{2K2}) \wedge ... \wedge p_m(T_{m1},...,T_{mKm}) --> q(T_{q1},...,T_{qKq})$
 - where each p_i is a predicate name and each T_{ij} is a term: i.e.
 - a constant
 - a variable or
 - a function term
- A single predicate $p(T_1,...,T_K)$ is a degenerate implication:
 - TRUE --> $p(T_1,...,T_K)$

Prolog Programming and Horn Clauses

- In Prolog, a typical Horn Clause of the form
 - $p_1(T_{11},...,T_{1K1}) \wedge p_2(T_{21},...,T_{2K2}) \wedge ... \wedge p_m(T_{m1},...,T_{mKm}) --> q(T_{q1},...,T_{qKq})$
- is represented as:
 - $q(T_{q1}...,T_{qKq}) := p_1(T_{11},...,T_{1K1}), p_2(T_{21},...,T_{2K2}), ...,p_m(T_{m1},...,T_{mKm}).$ i.e.
 - the implicand is on the left most end,
 - the antecedents occur on the right of :- (read this as <--),
 - the commas separating the antecedents indicate conjunction, and
 - there is a period ending the clause.

Prolog Programming and Horn Clauses

- A Prolog program is a conjunction of Horn Clauses and the conjunction is implicit:
 - i.e. syntactically, a Prolog program is a list of Horn Clauses.
- A degenerate clause (with a single predicate) is referred to as a *fact* in Prolog: e.g.
 - nat(0).
- A typical Horn Clause is referred to as a rule in Prolog: e.g.
 - nat(s(X)):-nat(X).
 - Note: A fact is a special form of a rule. End of Note.
- Argue that these two rules form a specification of <u>natural</u> <u>numbers</u> in Prolog.
- Exercise:
 - Specify the <u>addition</u> operation in Prolog.

