INSTRUCTION DIVISION FIRST SEMESTER 2018-2019 Course Handout Part II

In addition to part-I (General Handout for all courses appended to the time table) this portion gives further specific details regarding the course

Course No.: CS / IS F214

Course Title: Logic in Computer Science

Instructor-in-Charge: Shan Sundar Balasubramaniam (email: sundarb)

Instructors:

Shan Sundar Balasubramaniam (email: sundarb) - Tut. Sections 1 & 3

Jagat Sesh Chella (email: jagatsesh) - Tut. Sections 2 & 4

1. a. Scope:

The purpose of this course is to introduce the formal study of Logic to Computer Science undergraduates. Within this context the course introduces *first order propositional and predicate logics* as well as special topics such as *modeling*, *program verification*, and *temporal logics*.

Natural Deduction as a proof system for propositional and predicate logics is covered. **Soundness** and **Completeness** of proofs are covered but only for propositional logic. The course also covers some applications in modeling computing systems and reasoning about programs - in particular, **model checking** based on temporal logics and program verification using **Floyd-Hoare logic**. The relationship between formal logic and pragmatics of computing is highlighted via a few specific topics: (i) the **satisfiability** problem in propositional logic and its computational complexity (ii) **Horn-clause problem solving** as a basis for programming.

b. Course Objectives

On completion of this course the student shall be able to

- (i) write statements and proofs in first order predicate logic,
- (ii) explain the limitations of <u>first order predicate logic</u> and identify the need for higher-order logics or temporal logics in specific contexts
- (iii) state and argue <u>soundness and completeness</u> properties of logics
- (iv) write grammars for the syntax of logics
- (v) write algorithms for verifying satisfiability / validity where applicable





- (vi) relate problems in logic to problems in computation
- (vii) use logics to formally specify and verify computational properties.

2. Text Book:

T1: Michael Huth and Mark Ryan. *Logic in Computer Science – Modelling and Reasoning about Systems*. Cambridge University Press. 2nd Edition. 2004.

3. Course Plan:

3.a. Modules

Module #	Topics
I	Introduction to Logics and Proofs. Overview of the interplay between Logic and Computing.
II	Propositional Logic: Natural Deduction, Syntax and Semantics, Soundness and Completeness, Satisfiability and Validity – Forms and Algorithms.
III	Predicate Logic: Syntax and Semantics, Logic Programming, Natural Deduction, Limitations of First Order Logic.
IV	Program Verification: Floyd-Hoare Logic – Pre-conditions, Post-conditions, and Loop Invariants; Verification of imperative programs – Partial and Total correctness.
V	Temporal Logic: LTL, CTL, and Model Checking

3. b. Lecture Schedule:

Lecture	Module	Topic	Learning Outcome(s)	
#	#		[The student will be able to:]	
L1	l.a	Why study Logic?	 state a few reasons to study logi and proofs 	cs -
L1	I.b	A broad and selective History of Logic and Proofs (HLP) - Part I: Euclid's Formalization of Geometry. Hilbert's Program and the Formalization	 state typical issues and debates in formalizing proofs (and mathematical arguments in gene state the distinction between 	



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		problem. Russell's efforts in formalization and Russel's paradox .		axioms, statements, and proofs	
L2	I.c	HLP – Part II: Brouwer's intuitionism – Intuitionist's Critique on Proof by	•	distinguish between a constructive proof and a non-constructive proof	
				•	
		Contradiction and (the use of) Law of Excluded Middle	•	state the intuitionist position /	
				critique on non-constructive proofs	
L2	I.d	HLP – Part III: Godel's Incompleteness	•	state soundness and completeness	-
		result and its impact on the		requirements of proof systems and	
		formalization of mathematics	•	explain – at a high level – the	
				implications of Godel's	
				incompleteness result	
L2	l.e	<u>Logic and Computing (L&C) – Part I</u> :	•	informally define the notion of	-
		Godel, Church, and Turing:		"computability"	
		Computability – Systems for defining	•	state the equivalence of schemes /	
		computability; Church-Turing Thesis .		mechanisms for computability	
L3	I.f	L&C – Part II: Logic and Computing:	•	state the ${\bf Halting\ Problem}$ and argue	
		Non-computable Problems – Example:		that it is not computable	
		Halting			
L4	l.g	<u>L&C – Part III</u> : Time Complexity,	•	define complexity classes P and NP	-
		Complexity Classes, Is P = NP?	•	state the implications of P=NP or P!=NP	
L4	I.h	<u>L&C – Part IV:</u> Boolean Satisfiability	•	verify satisfiability of Boolean	-
		(SAT) and Horn-Clause Satisfiability		expressions and Horn-style	
		problems; their relationship to the		expressions.	
		classes NP and P;			
L4	l.i	<u>L&C – Part V</u> : Logic and Programming:	•	describe the relation between logic	-
		Logic Programming and Prolog.		and programming at a high level	
L5	II.a	Propositional Logic: Natural	•	explain the notational conventions	T1 Sec.
		Deduction as a Proof System.		used in Natural Deduction	1.2.1
L5	II. b	Propositional Logic: Natural	•	apply proof rules involving	T1 Sec.
		Deduction: Conjunction Rules		conjunction	1.2.1
L5	II.c	Propositional Logic: Natural	•	apply <i>modus ponens</i> and implication	T1 Sec.
		Deduction: Implication Rules		introduction	1.2.1
L6	II.d	Propositional Logic: Natural	•	apply proof rules involving	T1 Sec.
		Deduction: Disjunction Rules		disjunction	1.2.1

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II.e II.f	Propositional Logic: Natural Deduction: Negation Rules Propositional Logic: Natural Deduction: Double Negation Rules	•	apply proof rules involving negation apply proof rules involving double	T1 Sec. 1.2.1 T1 Sec.
	Propositional Logic: Natural	•	apply proof rules involving double	
	_	•	apply proof rules involving double	T1 Sec.
II.g	Deduction: Double Negation Rules			4 0 4
II.g			negation	1.2.1
	Propositional Logic: Natural	•	apply <i>modus tollens</i> , LEM, and PbC	T1 Sec.
	Deduction: Derived Rules: <i>Modus</i>		in proofs	1.2.2
	Tollens, Law of Excluded Middle	•	derive modus tollens, LEM, and PbC	
	(LEM), Proof by Contradiction (PbC).		from basic proof rules	
II.h	Syntax: Context Free Grammars (CFG):	•	define CFG s for simple constructs	T1 Sec.
	Notation and Examples			1.3
II.i	Context Free Grammars: Parse Trees:	•	illustrate and explain the internal	T1 Sec.
	Examples		structure of context-free constructs	1.3
II.j	Propositional Logic: Syntax: Well-	•	state the formal syntax of	T1 Sec.
	formed-formulas and grammar.			1.3
II.k	Propositional Logic: Syntax: Different	•		T1 Sec.
	forms of grammar: Precedence and			1.3
	Order of evaluation: Rules to handle		•	
	precedence and over-ride precedence		•	
l.j	Proof Techniques: Mathematical	•	articulate the generic structure /	-
	Induction			
			induction	
I.k	Proof Techniques: Structural	•	relate mathematical induction to	-
	Induction and Examples		structural induction	
1.1	Grammars, Parse Trees and Structural	•	illustrate the relation between	T1 Sec.
	Induction – Approach		languages and inductive proofs	1.4.2
		•		
			•	
l.m	Structural Induction - Examples	•	<u> </u>	T1 Sec.
	·		•	1.4.2
II.I	Propositional Logic: Semantics: Truth	•		T1 Sec.
	•		·	1.4.1
		•	•	
			completeness results w.r.t.	
1			·	
			propositional logic	
II.m	Propositional Logic: Semantics:	•	propositional logic argue or verify the equivalence (or	T1 Sec.
	II.i II.j II.k I.j	II.h Syntax: Context Free Grammars (CFG): Notation and Examples II.i Context Free Grammars: Parse Trees: Examples II.j Propositional Logic: Syntax: Well- formed-formulas and grammar. II.k Propositional Logic: Syntax: Different forms of grammar: Precedence and Order of evaluation: Rules to handle precedence and over-ride precedence I.j Proof Techniques: Mathematical Induction I.k Proof Techniques: Structural Induction and Examples I.l Grammars, Parse Trees and Structural Induction — Approach I.m Structural Induction - Examples	II.h Syntax: Context Free Grammars (CFG): Notation and Examples II.i Context Free Grammars: Parse Trees: Examples II.j Propositional Logic: Syntax: Wellformed-formulas and grammar. II.k Propositional Logic: Syntax: Different forms of grammar: Precedence and Order of evaluation: Rules to handle precedence and over-ride precedence I.j Proof Techniques: Mathematical Induction I.k Proof Techniques: Structural Induction and Examples I.l Grammars, Parse Trees and Structural Induction — Approach I.m Structural Induction - Examples I.m Structural Induction - Examples	II.h Syntax: Context Free Grammars (CFG): Notation and Examples II.i Context Free Grammars: Parse Trees: Examples illustrate and explain the internal structure of context-free constructs illustrate and explain the internal structure of context-free constructs illustrate and explain the internal structure of context-free constructs illustrate and explain the internal structure of context-free constructs illustrate and explain the internal structure of context-free constructs illustrate and explain the internal structure of context-free constructs illustrate the formal syntax of propositional logic as a CFG write grammar rules to define the syntax of propositional logic with a specific order of evaluation articulate the generic structure / template of proofs by mathematical induction induction articulate the generic structure / template of proofs by mathematical induction illustrate the relation between languages and inductive proofs explain structural induction as a proof technique write proofs using structural induction induction II.i Propositional Logic: Semantics: Truth Tables, Soundness and Completeness interpret sentences in propositional logic state the soundness and induction interpret sentences in propositional logic interpret sentences in propositional logic state the soundness and induction interpret sentences in propositional logic interpret sentences in propositional logic interpret sentences induction induction induction induction induction induction inducti



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				formulas	
L10	II.n	Propositional Logic: Syntax: Grammars	•	define a grammar for a given normal	-
		for Normal Forms		form	
L11	II.o	Propositional Logic : Soundness	•	state soundness arguments in	T1 Sec.
				general and soundness arguments	1.4.3
				for logics in particular	
L11	II.p	Propositional Logic: Soundness Proof	•	argue that proofs in propositional	T1 Sec.
				logic are sound	1.4.3
L12	II.q	Propositional Logic: Completeness	•	state completeness arguments in	T1 Sec.
				general and completeness	1.4.4
				arguments for logics in particular	
L12	II.r	Propositional Logic: Completeness	•	state completeness arguments for	T1 Sec.
		Proof Overview		propositional logic	1.4.4
L12	II.s	Propositional Logic: Completeness	•	argue that propositional logic is	T1 Sec.
		Proof		complete using structural induction	1.4.4
L13	II.t	Propositional Logic: Conjunctive	•	recognize propositional formulas in	T1 Sec.
		Normal Form (CNF) and Disjunctive		CNF / DNF	1.5.2
		Normal Form (DNF)	•	rewrite propositional formulas in	
				CNF / DNF	
L13	II.u	Propositional Logic: Validity	•	verify whether a propositional	T1 Sec.
				formula is valid or not	1.5.1
L13	II.v	Propositional Logic: Validity of	•	verify validity of a formula in CNF	T1 Sec.
		specific forms (CNF), Algorithm(s) for	•	write a program to verify validity of	1.5.1
		validity		a formula in CNF	
L13	II.w	Propositional Logic: Validity and	•	verify whether a propositional	T1 Sec.
		Satisifiability; Satisfiability of specific		formula is satisfiable or not	1.5.1
		forms (DNF)	•	state the relation between validity	
				and satisifiability of propositional	
114	11 9	Propositional Logic: How Clauses	_	formulas	T1 Sec.
L14	II.x	Propositional Logic: Horn Clauses and Horn Formulas	•	distinguish Horn formulas from general formulas in propositional	1.5.3
		and norn rormalds		logic	1.5.5
			•	identify when a propositional	
				formula can be rewritten in Horn	
				form	

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L14	II.y	Propositional Logic: Satisifiability of Horn Formulas. Algorithm for Satisfiability of Horn Formulas	 verify whether a Horn formula is satisfiable or not write a program to verify validity of a Horn formula 	T1 Sec. 1.5.3
L14	II.z	Propositional Logic: Algorithm for Satisfiability of Horn Formulas	 write a program to verify validity of a Horn formula 	T1 Sec. 1.5.3
L15	III.a	Expressiveness of Propositional Logic – Limitations: Need for predicates	 illustrate statements and arguments that cannot be expressed in propositional logic explain why predicates are more expressive than propositions 	T1 Sec. 2.1
L15	III.b	Expressiveness of Propositional Logic – Limitations: Need for variables and quantification	 illustrate statements that cannot be expressed in propositional logic explain how variables and quantification add expressive power 	T1 Sec. 2.1
L16	III.c	Introduction to Predicate Logic : Predicates vs. Functions and Need for Function Terms	 distinguish between predicates and functions explain when function terms are more expressive than predicates 	T1 Sec. 2.1
L16	III.d	Predicate Logic: Features and Examples	write formulas in predicate logic	T1 Sec. 2.2
L17	III.e	Predicate Logic: Formal Syntax: Grammars and Parse Trees	 write grammar rules for defining terms write grammar rules for defining formulas in Predicate Logic use the grammar rules to parse formulas in Predicate Logic 	T1 Sec. 2.2.1 to 2.2.2
L17	III.f	Predicate Logic: Theorem Proving : Proof Steps; Automating the proof steps.	 explain how proofs in predicate logic can be approached. state issues in automating proofs steps 	-
L17	III.g	Horn Clause Programming: Theorem Proving and Logic Programming, Programming with Horn Clauses	 explain the relation between Theorem Proving and Logic Programming state how Horn Clauses can be used for programming 	-

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L18	III.h	Logic Programming and Prolog : Proof Steps and Term Unification .	•	perform unification of terms	-
L18	III.i	Logic Programming and Prolog: Proof Search, Backtracking	•	apply search and backtracking to execute a Prolog program	
L18	III.j	Programming in Prolog: Examples – Facts and Rules; Queries	•	write simple programs in Prolog	-
L19	III.j	Programming in Prolog: Inductive / Recursive definitions, Lists and Other Data Structures	•	write programs using recursion in Prolog	-
L20	III.k	Programming in Prolog: Problem Solving in Prolog	•	write reasonably large programs to solve complex problems using Prolog	-
L21	III.I	Predicate Logic: Proofs and Proof Rules: Introduction: Need for a Substitution Operation	•	explain / illustrate the need for a substitution operation in applying proof rules of Predicate Logic	T1 Sec. 2.2
L21	III.m	Predicate Logic: Syntax: Free Variables and Bound Variables	•	identify and distinguish between free and bound variables in Predicate Logic formulas	T1 Sec. 2.2.3
L21	III.n	Predicate Logic: Substitution: Definition and Examples	•	apply substitution on Predicate Logic formulas	T1 Sec. 2.2.4
L22	III.o	Predicate Logic: Proof Rules: Equality	•	explain the need for rules for equality use the rules for equality	T1 2.3.1
L22	III.p	Predicate Logic: Natural Deduction: Rules for Universal Quantification	•	apply rules for universal quantification in proofs of predicate logic formulas	T1 Sec. 2.3.1
L22	III.q	Predicate Logic: Natural Deduction: Rules for Existential Quantification	•	apply rules for existential quantification in proofs of predicate logic formulas	T1 Sec. 2.3.1
Self Study	III.s	Predicate Logic: Syntactic Equivalence	•	deduce equivalences of formulas in predicate logic	T1 Sec. 2.3.2



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L23	III.r	Predicate Logic: Introduction to Semantics: Models and Interpretation – Examples Predicate Logic: Semantics: The Model-Checks Relation	•	formally define models for different theories interpret formulas in predicate logic according to given models apply the model-checks relation on predicate logic formulas	T1 Sec. 2.4.1 T1 Sec. 2.4.1
		Widder Checks Neiddlen		predicate logic formulas	
L25	III.t	Predicate Logic: Semantics: Semantic Entailment	•	illustrate and explain semantic entailment in predicate logic	T1 Sec. 2.4.2
L25	III.u	Predicate Logic: Semantics: Validity and Satisifiability	•	state and explain validity and satisfiability in Predicate Logic explain and illustrate issues in verifying satisfiability and validity of formulas in Predicate Logic	T1 Sec. 2.4.3
L25	III.v	Predicate Logic: Soundness and Completeness	•	state and explain claims regarding Soundness and Completeness of Predicate Logic	-
L26	l.n	Proof Techniques: Cantor's Diagonalization – Examples	•	articulate Cantor's diagonalization technique and apply the same on simple examples	-
L27	1.0	Review: Undecidability and the Halting Problem	•	state and explain the notion of undecidability state and explain the Halting problem. explain the use of diagonalization in proving the undecidability of Halting problem	-
L27	III.w	Predicate Logic: <u>Validity is</u> <u>Undecidable</u> : Proof sketch using Diagonalization	•	provide an proof argument for why Validity in Predicate Logic is undecidable using diagonalization	-
L28	III.w	Proof Techniques: Reduction and the use of Reduction to prove computability	•	explain the concept of reduction (between problems) apply reduction to derive simple undecidability proofs of problems	
L28	III.w	Predicate Logic: <i>Validity is</i>	•	provide an explanation for why	T1 Sec.

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		Undecidable: Proof approach using Reduction	Validity in Predicate Logic is undecidable using the idea of a reduction to a known undecidable problem	2.5
L29	III.y	Predicate Logic: Expressiveness: Inexpressible Properties: Example and Proof.	 provide examples of statements that cannot be expressed in First Order Predicate Logic provide an intuitive argument as to why such statements are inexpressible in First Order Predicate Logic provide a rigorous argument as to why reachability in graphs is inexpressible in First Order Predicate Logic 	T1 Sec. 2.6
L29	III.z	Existential Second Order Logic: Expressiveness: Examples(s)	 provide examples of statements that require existential quantification over predicates 	T1 Sec. 2.6.1
L29	III.aa	Universal Second Order Logic: Expressiveness: Example(s)	 provide examples of statements that require universal quantification over predicates 	T1 Sec. 2.6.2
L30	IV.a	Program Verification: Floyd-Hoare Logic: Pre-conditions and Post- Conditions	 describe what pre-conditions and post-conditions are 	T1 Sec. 4.2.2
L30	IV.b	Program Verification: Floyd-Hoare Logic: Rules for Assignment and Sequencing	 write pre-conditions and post- conditions for assignment statements compose pre-conditions and post- conditions over sequential statements 	T1 Sec. 4.2.2 & 4.3.1
L31	IV.b	Program Verification: Floyd-Hoare Logic: Rule for Conditional Statements	 compose pre-conditions and post- conditions over conditional statements 	T1 Sec. 4.2.2 & 4.3.1
L31	IV.c	Program Verification: Floyd-Hoare	-	T1 Sec. 4.2.2 &

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		Logic: Meta-Rules			4.3.1
L31	IV.d	Program Verification: Verification of Straight-Line programs - Examples	•	apply rules of Floyd-Hoare Logic to verify properties of straight line programs	T1 Sec. 4.3.1
L32	IV.e	Program Verification: Floyd-Hoare Logic: Program Variables and Logical Variables	•	identify and distinguish between Program Variables and Logical Variables in verification using Floyd-Hoare Logic illustrate the need for Logical Variables with examples	T1 Sec. 4.3.1
L32	IV.f	Program Verification: Partial Correctness and Total Correctness	•	distinguish between partial correctness arguments and total correctness arguments write termination proofs	T1 Sec. 4.3.1
L32	IV.g	Program Verification: Floyd-Hoare Logic: Loop Invariants	•	illustrate and explain what loop invariants are	T1 Sec. 4.2.2 & 4.3.1
L32	IV.h	Program Verification: Floyd-Hoare Logic: Verifying correctness of Loops: Partial Correctness using Invariants – Examples	•	write loop invariants given simple loops provide correctness arguments using loop invariants	T1 Sec. 4.2.2 & 4.3.1
L33	IV.i	Program Verification: Floyd-Hoare Logic: Proof Rules and Verification Examples	•	provide partial correctness proofs for small programs	T1 Sec. 4.3.1 & 4.3.2
L34	IV.j	Program Verification: Case Study	•	apply Floyd-Hoare logic to verify properties of reasonably large programs	-
L35	IV.k	Program Verification: Approaches and Limitations	•	describe issues and limitations in proving correctness of programs	T1 Sec. 3.1
L35	V.a	Introduction to Dynamic Logics	•	provide examples where "time" needs to be specified explicitly in logical formulas and the form in	T1 Sec. 3.1



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			which it needs to be specified
L35	V.b	Software Modeling: State Machines	• illustrate and explain specification using state machines T1 Sec. 2.7.1
L36	V.c	Software Modeling: State Machines as Models and Model Checking	 describe model checking define state machines for simple problems / computations T1 Sec. 2.7.1
L36	V.d	Linear-time Temporal Logic (LTL) – Introduction	• describe LTL T1 Sec. 3.2
L37	V.e	LTL: Syntax and Formulas	 provide examples of formulas in LTL illustrate internal structure of formulas in LTL
L37	V.f	LTL: Semantics: Transitions and Paths	• interpret LTL formulas T1 Sec. 3.2.2
L38	V.g	LTL: Specifications: Examples	• write LTL specifications for complex scenarios T1 Sec. 3.2.3
L39	V.h	Properties expressible in Temporal Logics – Examples	 define the problem of model checking for a particular case T1 Sec. 3.3.1
L39	V.i	Limitations of Linear Time and LTL:	• state and illustrate the limitations of LTL T1 Sec. 3.4
L40	V.j	Branching Time and Branching Time Temporal Logic (CTL) - Introduction	• read, write, and explain CTL T1 Sec. 3.4
L40	V.k	CTL – Semantics of Temporal Operators, Examples.	 interpret CTL formulas differentiate between interpretation in LTL and that in CTL
L41	V.I	Model Checking in LTL/CTL	perform model checking of properties stated in CTL/LTL T1 Sec. 3.6.1
L42	-	Course Summary	

4. Evaluation



4. a. Evaluation Scheme:

Component	Weight	Date	Remarks
Quizzes (3)	(3 x 20M =) 60M	In Tutorial Sessions	Open Book (scheduled on one- class notice)
Assignment	40M	in Oct. & Nov.	Take Home
Mid-Term Test (90 minutes)	43M	9/10 2:00 - 3:30 PM	(scheduled centrally) Open Book
Comprehensive Exam (120 minutes)	57M	3/12 FN	(scheduled centrally) Open Book
TOTAL	200M	-	-

4. b. Make-up Policy:

- No Make-up will be available for the Assignment under any condition.
- Late submission of an assignment will incur a penalty of 25% per 24 hours after the deadline.
- There will be one make-up (for all three quizzes put together) i.e. a student can take a make-up for at most one quiz. The make-up quiz will be conducted after all the regular quizzes and the coverage for that will be announced later. It is the responsibility of the student to request for a make-up i.e. a make-up is not automatic.
- <u>Prior Permission</u> of the Instructor is usually required to get a make-up for the mid-term test.
- <u>Prior Permission of (Associate) Dean, Instruction</u> is usually required to get a <u>make-up for the comprehensive exam.</u>
- A make-up shall be granted only in genuine cases where in the Instructor's / Dean's judgment the student would be physically unable to appear for the quiz/test/exam. Instructor's / Dean's decision in this matter would be final.

4.c. Fairness Policy:

- Student teams are expected to work on their own on assignments.
- All students are expected to contribute equally within a team. The instructor's assessment regarding the contributions of team members would be final.
- Any use of unfair means in quizzes, assignment, or test/exam will be reported to the Unfair means committee and will be subject to the severest penalty applicable:
 - Unfair means would include copying from other students or from the Web or from other sources of information including electronic devices.
 - All parties involved would be treated equally responsible: <u>allowing others to copy one's</u> work is enabling unfair means and is equally un-acceptable.

Web: www.pilani.bits-pilani.ac.in

5. Consultation / Office Hours:





- Mondays and Wednesdays: 12.15pm to 1.15pm and Thursdays: 2.30pm to 3.30pm in WILP office (1st floor, Library Building, Rear-entrance).
- Or by appointment via email
- **6. Contents and Notices:** All lecture slides will be posted on the course website on Nalanda. Notices concerning this course will be displayed online (on Nalanda) only. If there is a need, email would be used on short notice (12 hours) and only BITS Pilani mail would be used.

Instructor-In- Charge, CS F214.

