

Numerical Integration

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Most integrals cannot be computed analytically using standard functions, and require introducing special functions such as the error function

$$\text{erf}(x) = \frac{2}{\sqrt{\pi}} \int_0^x e^{-t^2} dt$$

How do we compute

$$J = \int_a^b f(x) dx$$

numerically - called numerical integration or quadrature

(1)

(comes from old days when gridded paper was used to count number of squares under a graph)

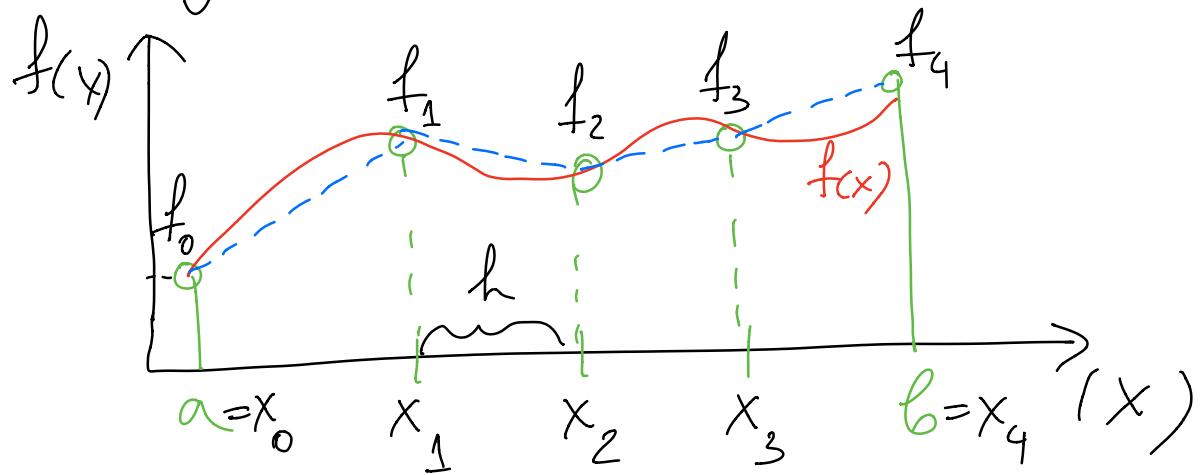
Simple idea:

Approximate function by
(piecewise or global) polynomial
and integrate polynomial

Note : Same works for numerical differentiation, as you saw in Worksheets. But, differentiation is easy so we mostly use numerical differentiation for solving differential equations
(not covered in this course)

(2)

E.g. Piecewise linear interpolation:



$$x_k = a + k \cdot h$$

$$h = \frac{b-a}{n} = \text{grid spacing}$$

Since we approximate area under the graph of $f(x)$ by a sum of areas of trapezoids we call this the **composite trapezoidal rule**:

$$J \approx h \cdot \left(\frac{f_0 + f_1}{2} + \frac{f_1 + f_2}{2} + \frac{f_2 + f_3}{2} + \frac{f_3 + f_4}{2} \right) \quad (3)$$

$$J \approx h \left(\frac{f_0 + f_n}{2} + \sum_{j=1}^{n-1} f_j \right)$$

where $f_j = f(x_j)$

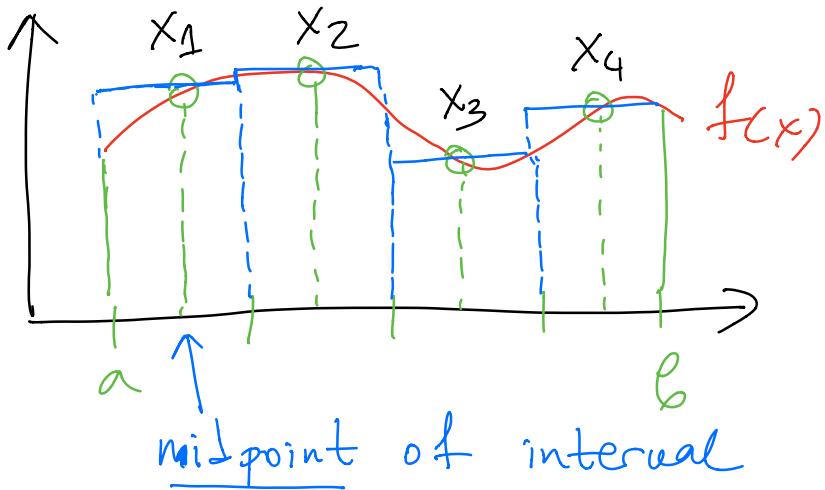
Usually written as

$$J = \int_a^b f(x) dx \approx h \sum_{j=0}^{n-1} f(x_j)$$

Trapezoidal $\rightarrow \frac{h}{2} (f(a) + f(b))$

It is very common to see student codes that miss this endpoint correction. If you don't want a correction at the end points, use instead the midpoint quadrature rule

(4)



Here $x_j = a + (j - 1/2) h, j=1, \dots, n$

$$h = \frac{b-a}{n}$$

and now we approximate the area by a sum of areas of rectangles

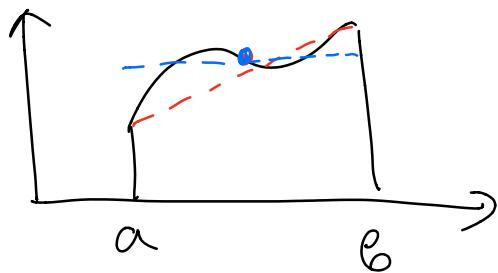
$$J \approx h \sum_{j=1}^n f(x_j) \text{ Midpoint}$$

The trapezoidal & midpoint rules are equally accurate & are the most basic formulas

(5)

How accurate are these quadrature rules?

Let's consider a single interval / trapezoid / rectangle first:



For midpoint rule: Expand $f(x)$ in a Taylor series around the midpoint, with remainder term, to show that

$$E = \int_{x-h/2}^{x+h/2} f(t) dt - h f(x)$$

$$= \frac{h^3}{24} f''(\xi)$$

(6)

where ξ is an inside interval.

Do on your own!

For trapezoidal rule, use formula for error of linear interpolant $P_1(x)$:

$$f(x) - P_1(x) = \frac{1}{2} f''(\xi(x))(x-a)(x-b)$$

$$\xi(x) \in [a, b]$$

$$\Rightarrow E = \left(\int_a^b f(x) dx \right) - \left(\int_a^b P_1(x) dx \right)$$

$$= \frac{1}{2} \int_a^b f''(\xi(x))(x-a)(x-b) dx$$

$$= \frac{f''(n)}{2} \int_a^b (x-a)(x-b) dx \quad (7)$$

where $\eta \in [a, b]$ according
to the mean-value theorem
for integrals. Do integral to get:

$$\left\{ \begin{array}{l} E_{\text{trap}} = -\frac{h^3}{12} f''(\xi) \\ E_{\text{mid}} = \frac{h^3}{24} f''(\xi) \end{array} \right.$$

$\xi \in [a, b], \quad b-a=h$

Note these are similar to
each other.

This was only one rectangle.

For many rectangles, i.e., for
composite trapezoidal rule, in
the worst case scenario all

errors will have the same sign and will add up,

$$|\gamma - \gamma_{\text{trap.}}| = \left| -\frac{h^3}{12} \sum_{k=1}^n f''(\xi_k) \right|$$

where $\xi_k \in [x_{k-1}, x_k]$

Assume $|f''(x)| < M$
for $x \in [a, b]$

$$\begin{aligned} \Rightarrow e_{\text{trap.}} &= |\gamma - \gamma_{\text{trap.}}| \leq \frac{h^3}{12} \cdot n \cdot M \\ &= \frac{h^3}{12} \cdot \frac{(b-a)}{h} \cdot M \end{aligned}$$

$$e_{\text{trap.}} \leq \frac{(b-a)}{12} \cdot M \cdot h^2$$

(9)

error = $O(h^2)$ for both
composite trapezoidal &
composite midpoint rules

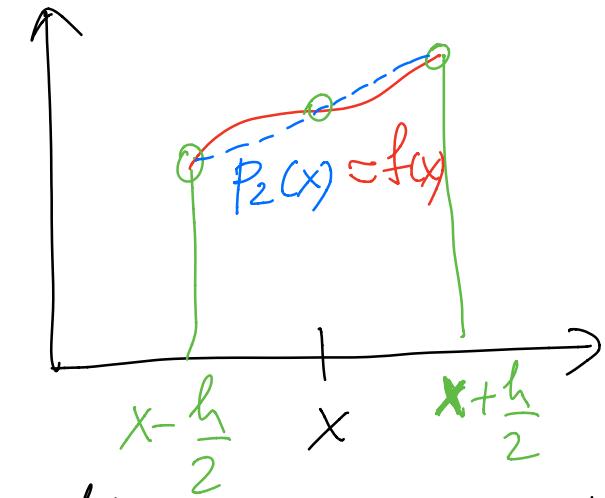
[Composite midpoint / trapezoidal
rule is second-order accurate

This means that if you double
the number of points / intervals,
the error decreases by a factor
of 4!

Note that this is just an
upper bound on error — a much
more precise estimate is
provided by the Euler-MacLaurin
theorem (see Wikipedia if
interested) (10)

To get more accuracy, we can use piecewise quadratic

interpolation : approximate function by a parabola over each interval. Let's take one interval only.



$$\int_{x-h/2}^{x+h/2} f(x) dx \approx \int_{x-h/2}^{x+h/2} P_2(x) dx$$

where P_2 is the quadratic interpolant.

(11)

You did this calculation in a worksheet, to get

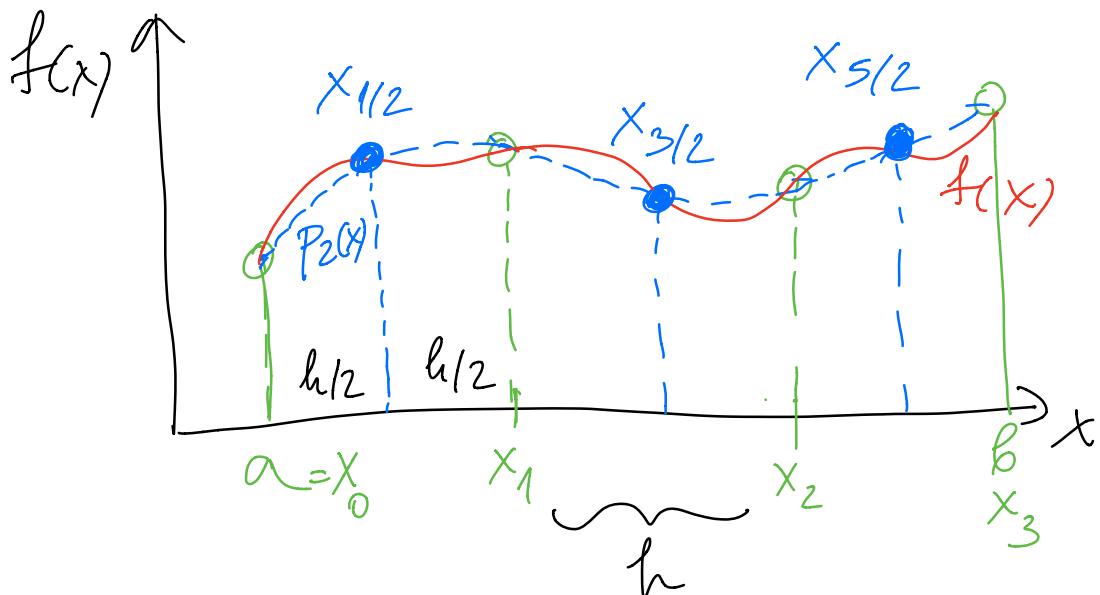
$$\int_{x-h/2}^{x+h/2} f(x) dx = \frac{h}{6} \left[f\left(x-\frac{h}{2}\right) + 4f(x) + f\left(x+\frac{h}{2}\right) \right]$$

This is called **Simpson's rule**.

To make this a composite rule,

split $[a, b]$ into pieces

(piecewise quadratic interpolation)



(12)

$$\begin{aligned}
 J &\approx \frac{h}{6} \left(f_0 + 4f_{1/2} + f_1 \right) \\
 &+ \frac{h}{6} \left(f_1 + 4f_{3/2} + f_2 \right) \\
 &+ \frac{h}{6} \left(f_2 + 4f_{5/2} + f_3 \right)
 \end{aligned}$$

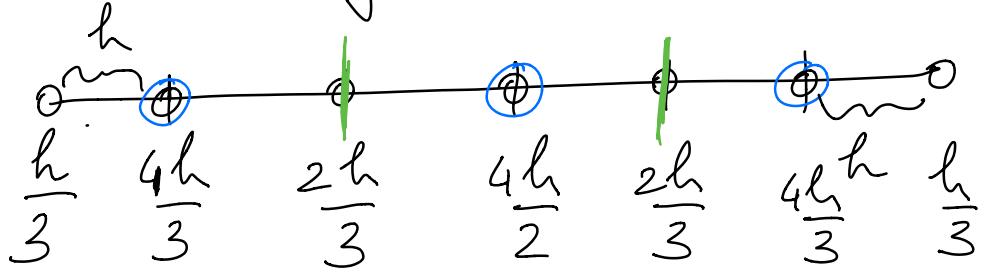
$$\begin{aligned}
 = \frac{h}{6} & \left[f_0 + 4f_{1/2} + 2f_1 + 4f_{3/2} \right. \\
 & \left. + 2f_2 + 4f_{5/2} + f_3 \right]
 \end{aligned}$$

$$\begin{aligned}
 \int_a^b f(x) dx &\approx \frac{h}{6} (f(a) + f(b)) \\
 &+ \frac{h}{3} \sum_{k=1}^{n-1} f(x_k) \\
 &+ \frac{2h}{3} \sum_{k=0}^{n-1} f(x_{k+1/2})
 \end{aligned}$$

Simpson quadrature

(13)

Another way to write it



Give first and last point weight $h/3$, even points weight $4h/3$, and odd ones $2h/3$
(compare to midpoint which gave all points weight h , and trapezoidal which gave first and last point weight $h/2$ and all others weight h).

Simpson's rule is only slightly more complicated *but 4th order accurate* so use it!

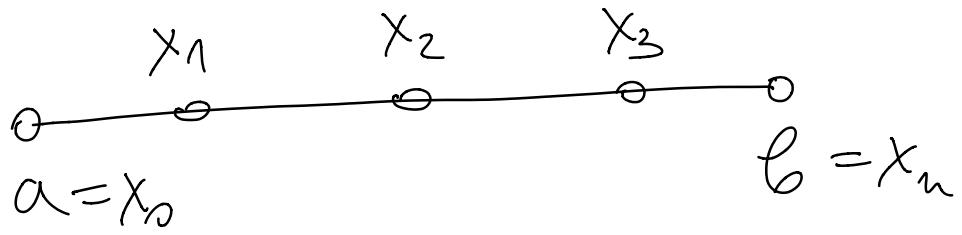
$$\epsilon_{\text{Simp}} = \left| J - J_{\text{Simp}} \right| = \frac{b-a}{2880} h^4 M$$

$$M = \max_{a \leq x \leq b} |f^{(4)}(x)|$$

Double number of points,
16 times smaller error!

We can get even better accuracy by using higher-degree polynomial approximants.

Let's focus on only one interval, but of course one can always split $[a, b]$ into pieces for composite quadrature



Choose interpolation nodes in $[a, b]$ (not necessarily equi-spaced!) and use Lagrange interpolation

$$f(x) \approx p_n(x) = \sum_{j=0}^n f(x_j) L_j(x)$$

Lagrange polynomial

$$\begin{aligned} \Rightarrow \int_a^b f(x) dx &\approx \int_a^b p_n(x) dx \\ &= \int_a^b f_j L_j(x) dx \end{aligned}$$

(16)

$$\int_a^b f(x) dx \approx \sum_{j=0}^n \left(\underbrace{\int_a^b L_j(x) dx}_{\omega_j} \right) f_j$$

$$J = \int_a^b f(x) dx \approx \sum_{j=0}^n w_j f(x_j)$$

↑
 quadrature
 weights

$$(*) \dots w_j = \int_a^b L_j(x) dx$$

↓
 quadrature
 nodes

All quadrature rules have this form, but different weights & nodes. Note that once we choose the nodes we can pre-compute the weights \overrightarrow{w}

(see Newton-Cotes quadrature on Wiki for tables)

(17)

The midpoint, trapezoidal & Simpson's rule were just special cases for $n=0$, $n=1$ and $n=2$. For these small n it is obvious where to place the nodes by symmetry, if we want to include the endpoints.

But what is the best choice of nodes & weights?

Note that once we choose the nodes, the weights are given by (*), so really the question is what are the "best" nodes?

Let us first clarify about what we mean by "best":

We want the most accurate integral for a generic smooth function for a given number of nodes.

And "most accurate" we define (but this choice is not unique or necessarily "optimal") to be the highest order of accuracy:

$$\text{error} = O\left(\frac{1}{n^P}\right)$$

where P is as large as possible for a given number of nodes n

(19)

Now let's observe the following:

A method that is p^{th} order accurate is exact for polynomials of degree \checkmark up to $p-1$ at least

This is a pretty useful way to check (validate) & derive formulas quickly and intuitively.

For example, we can derive Simpson's rule (quadratic interpolation, so exact for quadratic polynomials) which has $p=2+1=3$

[note that Simpson's rule actually has error $\sim 1/n^4$ but this is "luck" and is why we need "at least" in red above] (20)

$$\int_a^b p(x) dx = w_1 p(a) + w_2 p\left(\frac{a+b}{2}\right) + w_3 p(b)$$

← unknown
 w_1, w_2, w_3

for $p(x) \in \mathcal{P}_2$

This is equivalent [make sure you understand why, i.e., can prove both directions of equivalence]

To saying equation above is true for every element of a basis of \mathcal{P}_2 .

If we choose the monomials as a basis

$$p_1 = 1, \quad p_2 = x, \quad p_3 = x^2$$

we get 3 linear equations :

(21)

$$\int_a^b 1 dx = w_1 + w_2 + w_3 = b-a$$

$$\int_a^b x dx = w_1 a + w_2 \frac{a+b}{2} + w_3 b = \frac{b^2 - a^2}{2}$$

$$\begin{aligned}\int_a^b x^2 dx &= w_1 a^2 + w_2 \left(\frac{a+b}{2} \right)^2 + w_3 b^2 \\ &= \frac{b^3 - a^3}{3}\end{aligned}$$

Solve to get Simpson's rule

$$w_1 = w_3 = \frac{b-a}{6}, \quad w_2 = \frac{4(b-a)}{6}$$

Note:

In general, you can check if some formula you derived using polynomial interpolation of degree p is correct by testing it on the monomials $\{1, x, \dots, x^p\}$

Now, observe that the formula

$$(*) \dots w_j = \int_a^b L_j(x) dx$$

ensures that the quadrature rule is exact for all of the Lagrange polynomials up to degree $n-1$ if we have n points [note switch in

notation - n points not $n+1$ points like for interpolation]

Therefore the quadrature rule is exact for all polynomials of degree up to $n-1$ regardless of the choice of nodes.

But we can also choose the n nodes/points at will, so we have n additional degrees of freedom. Therefore we can hope that we can make the quadrature rule exact for polynomials of degree up to

$$(n-1) + n = 2n - 1$$

How can we find the nodes? Just plug all monomials of degree up to $2n-1$ into quadrature to get a nonlinear system of $2n$ equations for the \vec{w} and \vec{x} . $2n$ unknowns. See homework for $n=3$!

It turns out, however, that one can find a solution of this system of equations analytically by using orthogonal polynomials.

Theorem: If x_1, \dots, x_n are the zeros / roots of the Legendre polynomial of degree n , then the Gaussian quadrature rule

$$\int_a^b f(x) dx \approx \sum_{j=1}^n w_j f(x_j)$$

with w_j given by (*) is exact for polynomials of degree $2n-1$ or less

Nodes \vec{x} are called Gaussian nodes and are in $[-1, 1]$ (25)

as we mentioned at the end of the lecture on orthogonal polynomials. So these special points / nodes are good not only for interpolation but also for quadrature.

You can find tables of values or explicit formulas for the Gaussian nodes & weights for $n \leq 5$ on Wikipedia, and a code on the webpage for computing them numerically for any (modest) n .

This concludes our search for the "best" quadrature rule on $[-1, 1]$ - easy to change to $[a, b]$

Proof of theorem:

Let $p \in \mathcal{P}_{2n-1}$ and denote

the Legendre polynomials with

$$\{P_0, P_1, P_2, \dots, P_n\}$$

$$p = q P_n + r$$

where $q(x) \in \mathcal{P}_{n-1}$

$$P_n(x) \in \mathcal{P}_n$$

$$r(x) \in \mathcal{P}_{n-1}$$

residual of "polynomial long division"

Since x_j is a root of P_j ,

$$p(x_j) = q(x_j) P_n(x_j) + r(x_j) = r(x_j)$$

$$\int_{-1}^1 p(x) dx = \int_{-1}^1 q_r(x) P_n(x) dx + \int_{-1}^1 r(x) dx$$

Now, since $P_n(x)$ was constructed (using Gram-Schmidt) to be orthogonal to all $q(x) \in \mathcal{P}_{n-1}$,

$$(q_r, P_n) = \int q_r(x) P_n(x) dx = 0$$

$$\Rightarrow \int p(x) dx = \int r(x) dx = \\ (\text{since } r \in \mathcal{P}_{n-1} \text{ and } *) = \sum w_j r(x_j) \\ = \sum w_j f(x_j)$$

which is Gaussian quadrature.

(28)