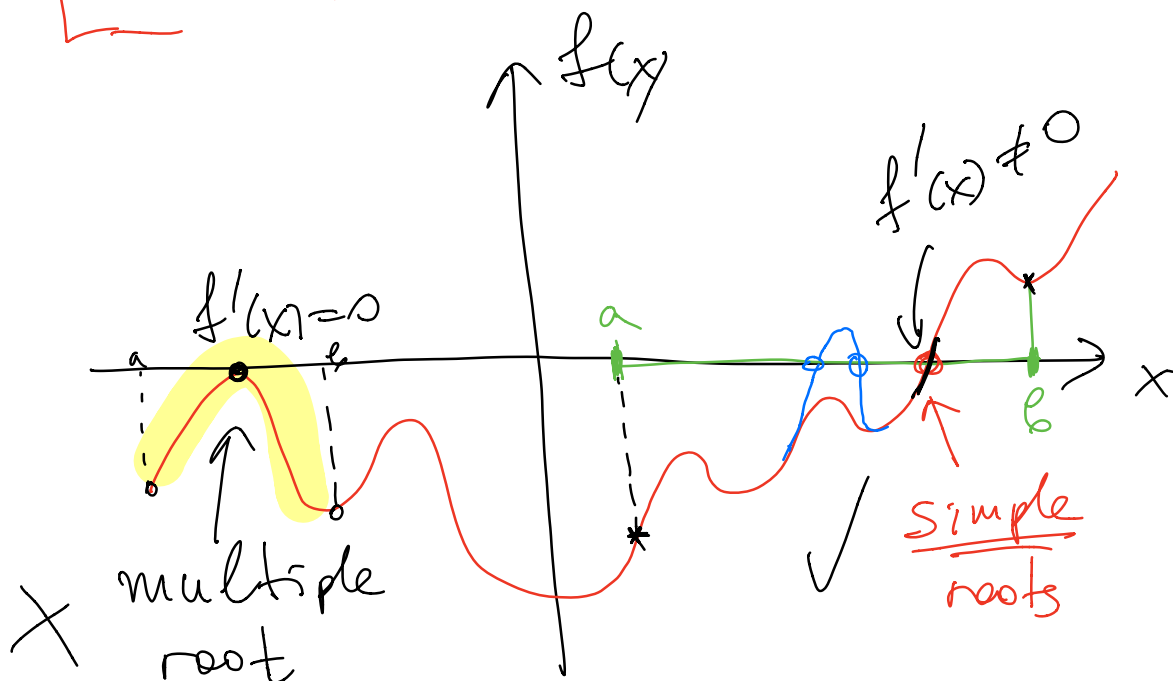


Nonlinear equations (in one variable)

$$x = \sqrt{c} \quad (\Leftrightarrow) \quad x^2 - c = 0$$

$$\cos x + x^2 - 7 = 0$$

Solve $f(x) = 0 \quad x \in [a, b]$
Find at least one solution
if one exists



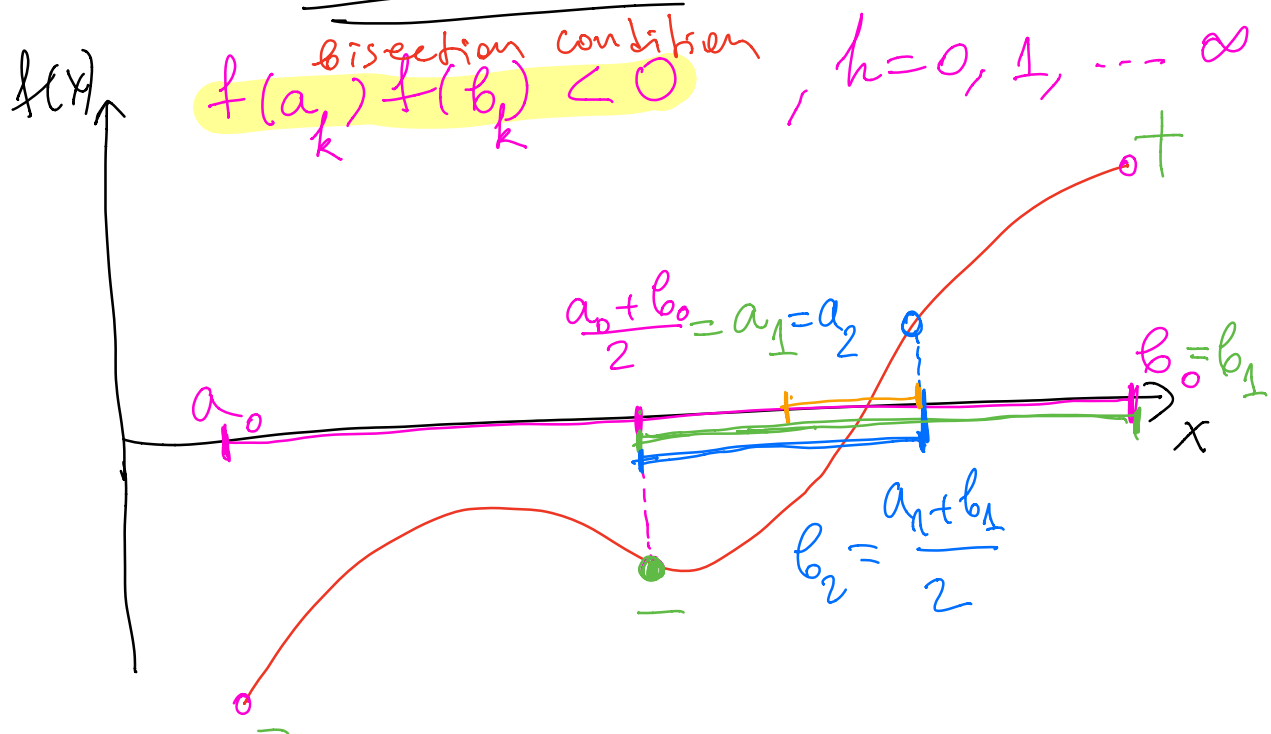
$f(x)$ is continuous on $[a, b]$

$$\exists f \quad f(a) \cdot f(b) < 0$$

$$\Rightarrow \exists x \in [a, b] \text{ s.t. } f(x) = 0$$

(Q) Why don't I write
 $f(a)f(b) \leq 0$?

Bisection method



Algorithm

Input: $[a, b]$ s.t. $f(a)f(b) < 0$
 $k \in \mathbb{Z}^+$
 k_{\max}

Output: $\begin{cases} \tilde{x} \text{ s.t. } f(\tilde{x}) \approx 0 \\ [a, b] \text{ s.t. true root is in } [\tilde{a}, \tilde{b}] \\ x \in [a, b], f(x) = 0 \end{cases}$

For $k = 0, 1, 2, \dots, k_{\max}$

$$x_k = \frac{a_k + b_k}{2}$$

$$f_k = f(x_k)$$

(actually, we only need
Sign of f_k)

If $f_k \cdot f(a_k) < 0$ then
left half $[a_k, x_k]$ is a bisection interval,
 $x \in [a_k, x_k]$

$$a_{k+1} = a_k ; b_{k+1} = x_k$$

else \rightarrow [if $f_k \cdot f(b_k) < 0$]
(keep right half)

$$b_{k+1} = b_k ; a_{k+1} = x_k$$

end if

\rightarrow [else if $f_k = 0$ then
return x_k]

End for

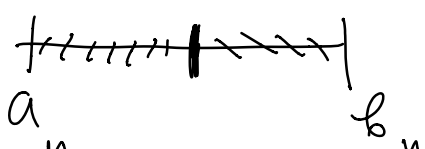
$$\text{Output} : x \approx x_{k_{\max}} = \frac{a_{k_{\max}} + b_{k_{\max}}}{2}$$

$$\text{and } x \in [a_{k_{\max}}, b_{k_{\max}}]$$

We know that $[n = k_{\max}]$

$$|x - x_n| \leq \frac{b_n - a_n}{2}$$

$x_n = \frac{a_n + b_n}{2}$


$$= \frac{b_0 - a_0}{2 \cdot 2^n}$$

Absolute error

$$e_n = |x - x_n| \leq \frac{b - a}{2^{n+1}}$$

error estimate

Given an error tolerance

$$\varepsilon \text{ s.t. } |x - x_n| < \varepsilon$$

$$\varepsilon \approx \frac{b - a}{2^{n+1}}$$

$$2^{n+1} > \frac{b-a}{\epsilon}$$

$$n+1 = \log_2 \left(\frac{b-a}{\epsilon} \right)$$

How many
times f(x)
will evaluate

$$n = \left\lceil \log_2 \left(\frac{b-a}{\epsilon} \right) \right\rceil$$

round
up

/ceil

in Matlab

round
down

→ ⌊

⌋

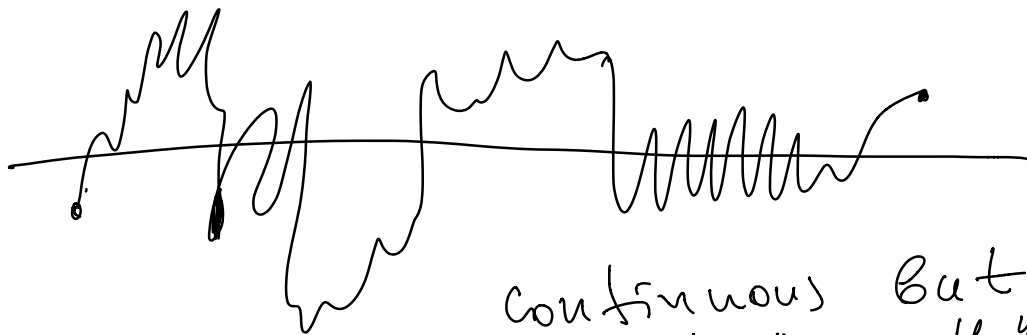
/floor

In Matlab :

$$n_{\text{est}} = \text{ceil}(\log_2((b-a)/\epsilon))$$

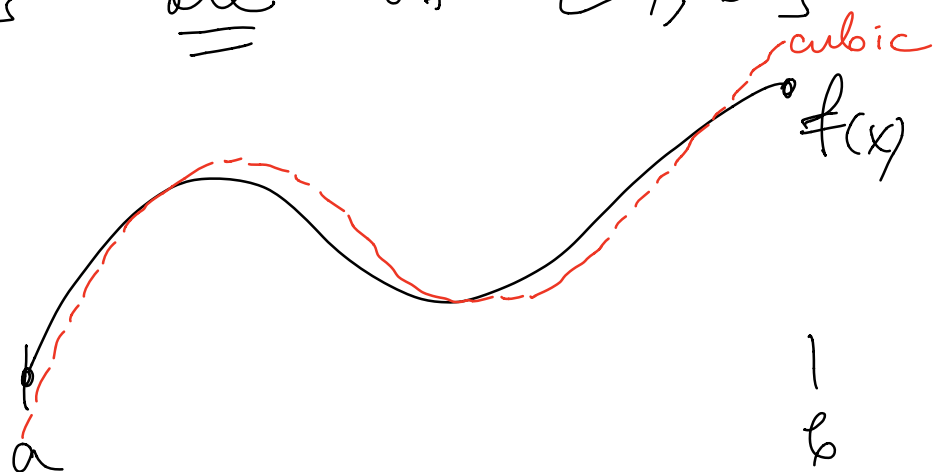
Our job is done

Can we do better?



continuous but
not "smooth"

$f(x)$ is "smooth" ^{on $[a, b]$} if it
can be approximated "well"
by a polynomial of low
degree (linear, quadratic, cubic)
over all of $[a, b]$



{ Another (related) is that
 $f(x)$ is smooth if it
is sufficiently differentiable
 $e^{1/x}$ is a bad example

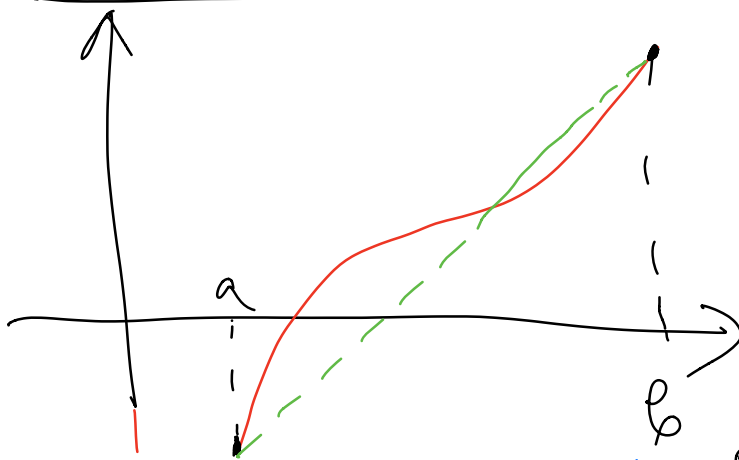
Idea: $p(x) \approx f(x)$
polynomial on $[a, b)$

Solve $p(x) = 0$ instead

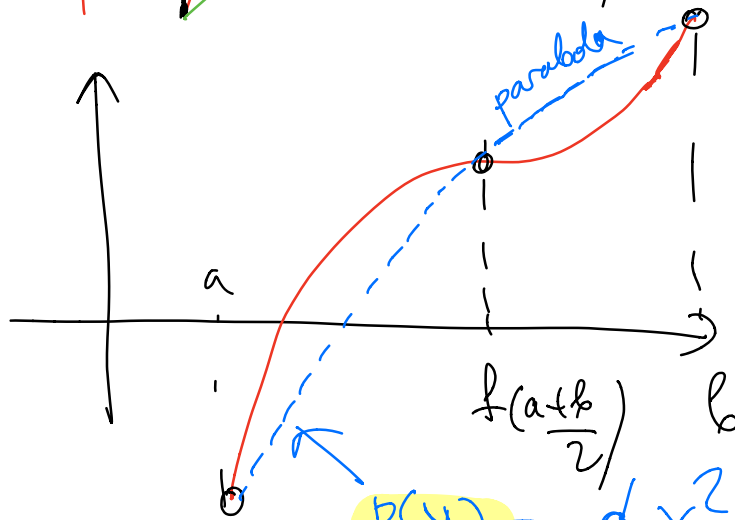
Typically $p(x)$ is linear
or quadratic

How do we find $p(x)$?

Option 1

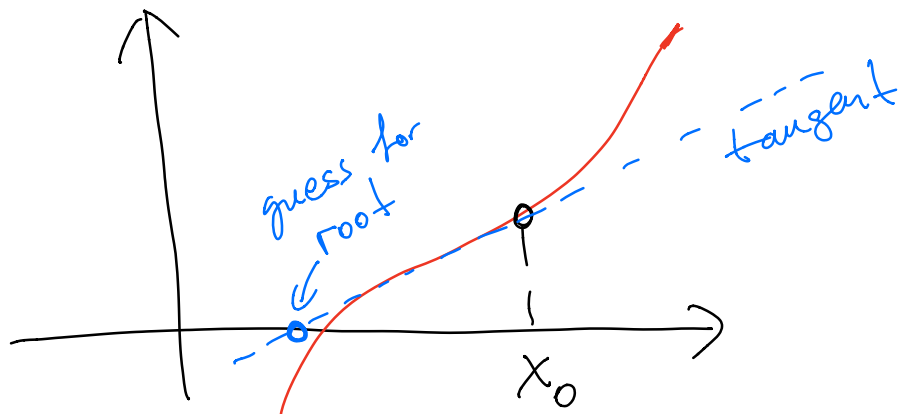


Interpolation



$$p(x) = dx^2 + \beta x + \gamma$$

Option 2 :



Option 2 : Taylor series

$$f(x) = \sum_{n=0}^{\infty} \frac{f^{(n)}(x_0)}{n!} (x-x_0)^n$$

$$p(x) = f(x_0) + f'(x_0)(x-x_0) + \frac{1}{2} f''(x_0)(x-x_0)^2 + \dots$$

$$p(x) = \sum_{n=0}^k \frac{f^{(n)}(x_0)}{n!} (x-x_0)^n$$

Truncated Taylor series

$$f(x) - p(x) = \frac{f^{(k+1)}(\xi)}{(k+1)!} (x-x_0)^{k+1}$$

Remainder

ξ is between x and x_0

