

Nonlinear equations

(in one variable)

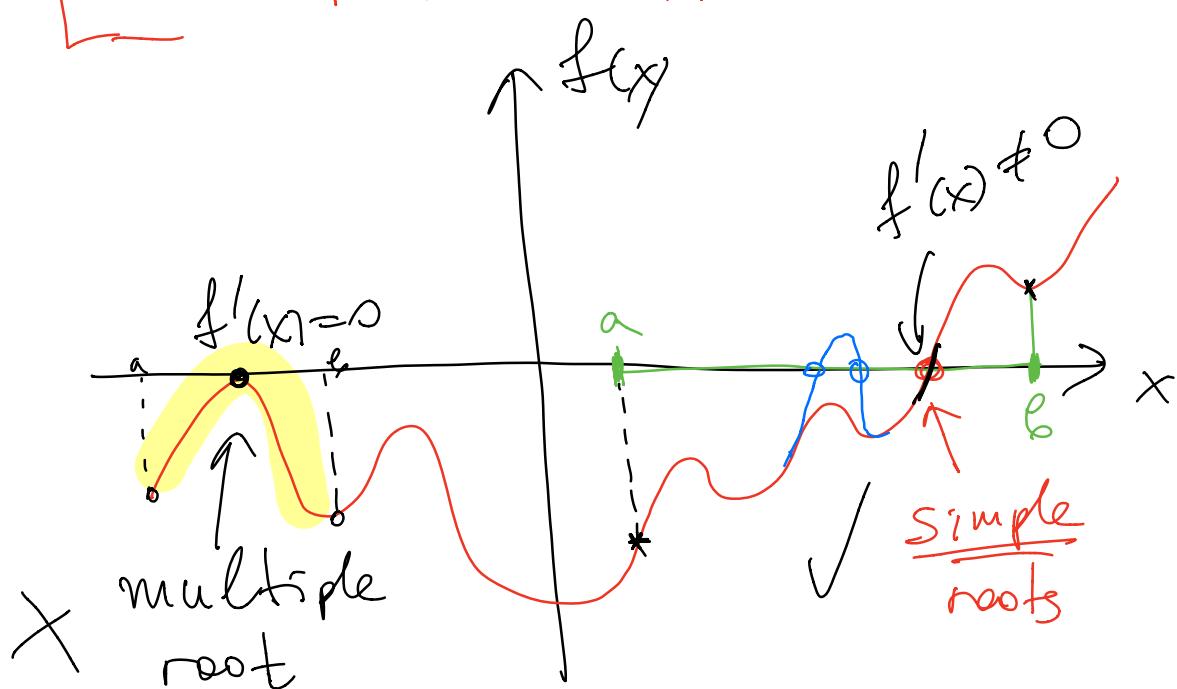
$$x = \sqrt{c} \Leftrightarrow x^2 - c = 0$$

$$\cos x + x^2 - 7 = 0$$

Solve $f(x) = 0 \quad x \in [a, b]$

Find at least one solution

if one exists



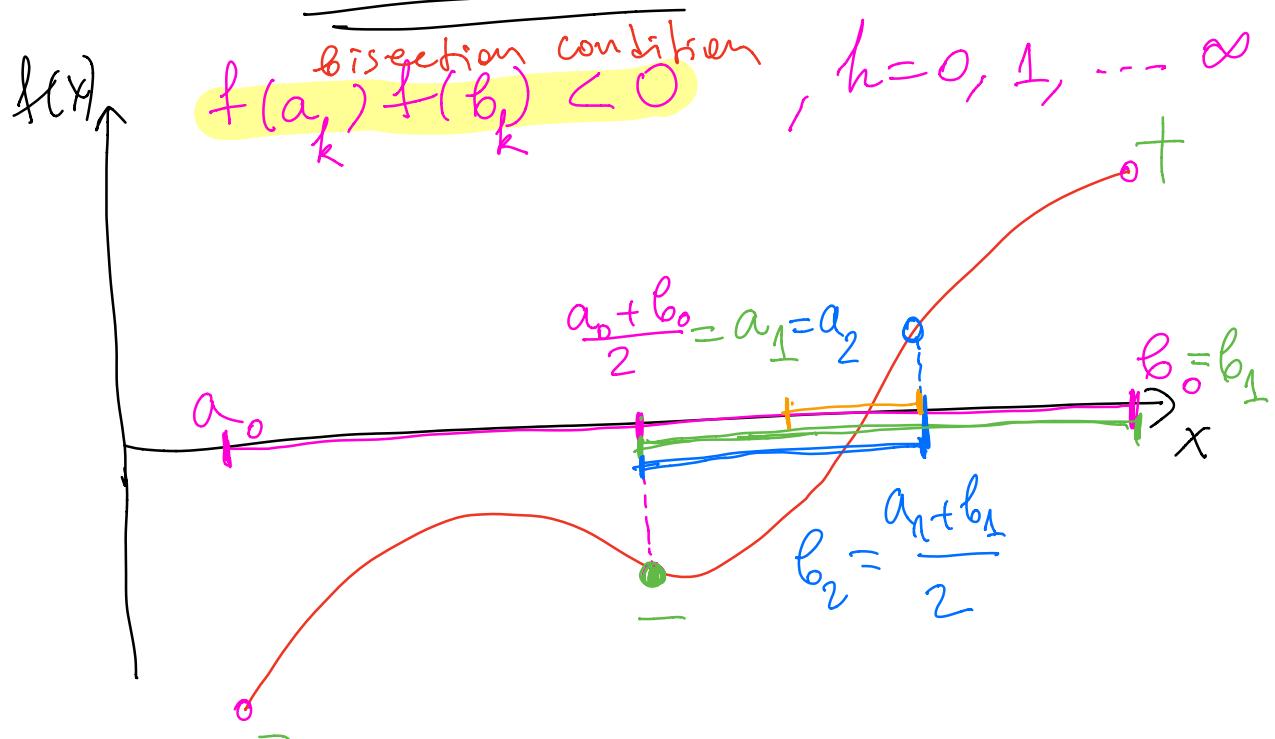
$f(x)$ is continuous on $[a, b]$

if $f(a) \cdot f(b) < 0$

$\Rightarrow \exists x \in [a, b] \text{ s.t. } f(x) = 0$

(Q) Why don't I write
 $f(a) f(b) \leq 0$?

Bisection method



Algorithm

Input: $[a, b]$ s.t. $f(a)f(b) < 0$
 $k \in \mathbb{Z}^+$
 \max

Output: $\left\{ \begin{array}{l} \tilde{x} \text{ s.t. } f(\tilde{x}) \approx 0 \\ [a, b] \text{ s.t. true} \\ \text{root is in } [a, b] \\ x \in [a, b], f(x) = 0 \end{array} \right.$

For $k=0, 1, 2, \dots, k_{\max}$

$$x_k = \frac{a_k + b_k}{2}$$

$$f_k = f(x_k)$$

(actually, we only need
Sign of f_k)

If $f_k \cdot f(a_k) < 0$ then
 left half $[a_k, x_k]$ is a bisection interval,
 $x \in [a_k, x_k]$

$a_{k+1} = a_k ; b_{k+1} = x_k$
else \rightarrow [if $f_k \cdot f(b_k) < 0$]
 (keep tight half)

$b_{k+1} = b_k ; a_{k+1} = x_k$
end if \leftarrow [else if $f_k = 0$ then
 return x_k]
End for

Output : $X \approx x_{k_{\max}} = \frac{a_{k_{\max}} + b_{k_{\max}}}{2}$

and $X \in [a_{k_{\max}}, b_{k_{\max}}]$

We know that [$n = k_{\max}$]

$$|x - x_n| \leq \frac{b_n - a_n}{2}$$

$x_n = \frac{a_n + b_n}{2}$

$$= \frac{\frac{b_0 - a_0}{2}}{2 \cdot 2^n}$$

Absolute error $b - a$

$$e_n = |x - x_n| \leq \frac{b - a}{2^{n+1}}$$

error estimate

Given an error tolerance

$$\epsilon \text{ s.t. } |x - x_n| < \epsilon$$

$$\epsilon \approx \frac{b - a}{2^{n+1}}$$

$$2^{n+1} > \frac{b-a}{\epsilon}$$

$$n+1 = \log_2 \left(\frac{b-a}{\epsilon} \right)$$

How many times evaluate f(x) will evaluate n = $\lceil \log_2 \left(\frac{b-a}{\epsilon} \right) \rceil$

round up / ceil in Matlab

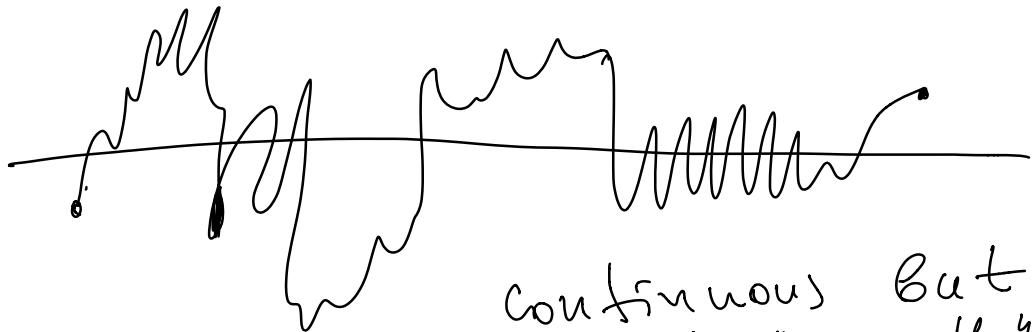
round down $\rightarrow \lfloor \quad \rfloor$ / floor

In Matlab :

$$n_{\text{est}} = \text{ceil}(\log_2((b-a)/\epsilon))$$

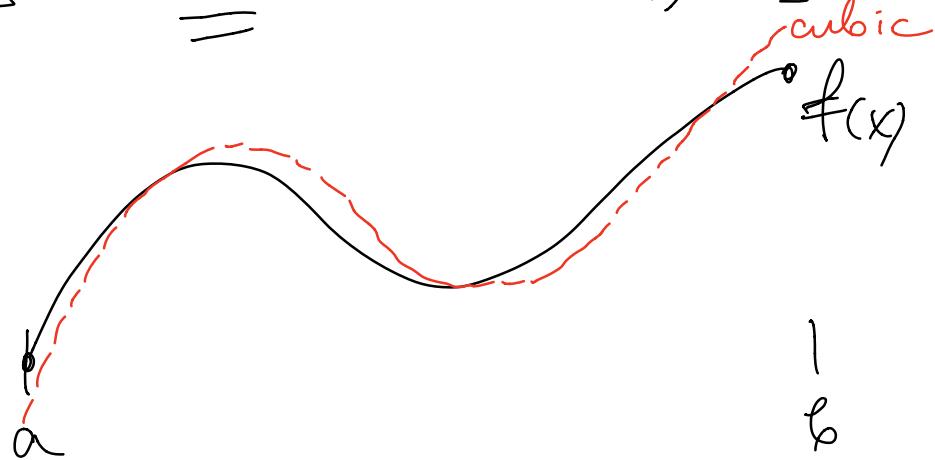
Our job is done

Can we do better?



continuous but
not "smooth"

{ $f(x)$ is "smooth" ^{on $[a, b]$} if it
can be approximated "well"
by a polynomial of low
degree (linear, quadratic, cubic)
over all of $[a, b]$



{ Another (related) is that
 $f(x)$ is smooth if it
is sufficiently differentiable

e^{nx} is a bad example

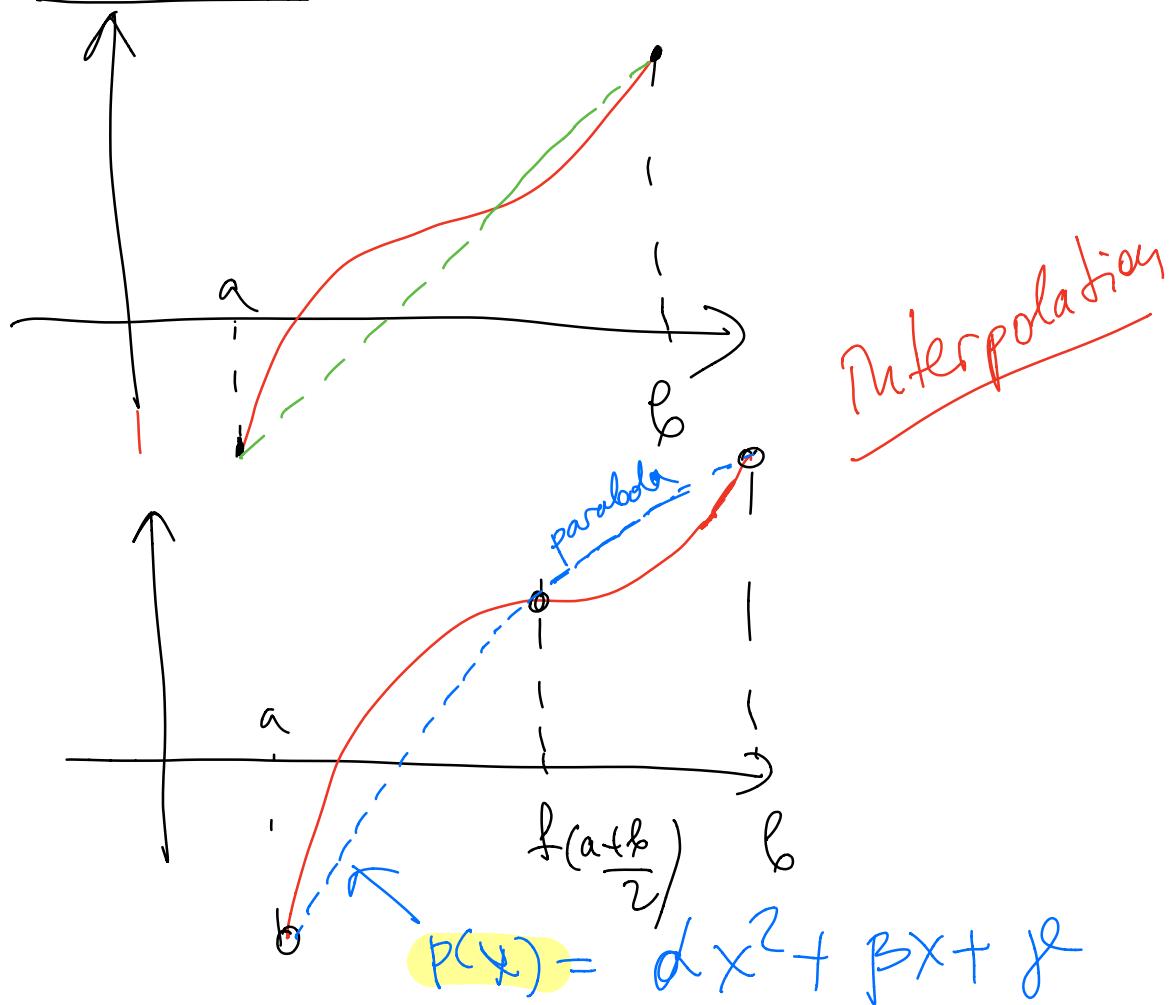
Idea: $p(x) \approx f(x)$
polynomial on $[a, b]$

Solve $p(x) = 0$ instead

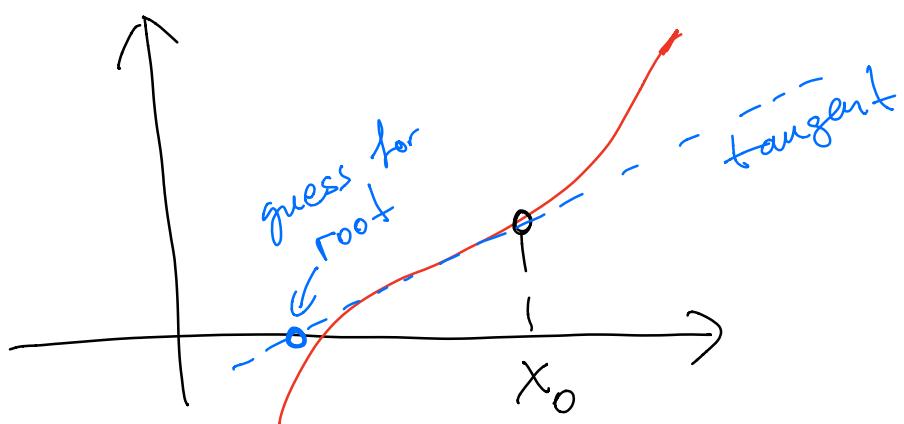
Typically $p(x)$ is linear
or quadratic

How do we find $p(x)$?

Option 1



Option 2 :



Option 2 : Taylor series

$$f(x) = \sum_{n=0}^{\infty} \frac{f^{(n)}(x_0)}{n!} (x-x_0)^n$$

$$\begin{aligned} p(x) &= f(x_0) + f'(x_0)(x-x_0) \\ &\quad + \frac{1}{2} f''(x_0) (x-x_0)^2 + \dots \end{aligned}$$

$$p(x) = \sum_{n=0}^k \frac{f^{(n)}(x_0)}{n!} (x-x_0)^n$$

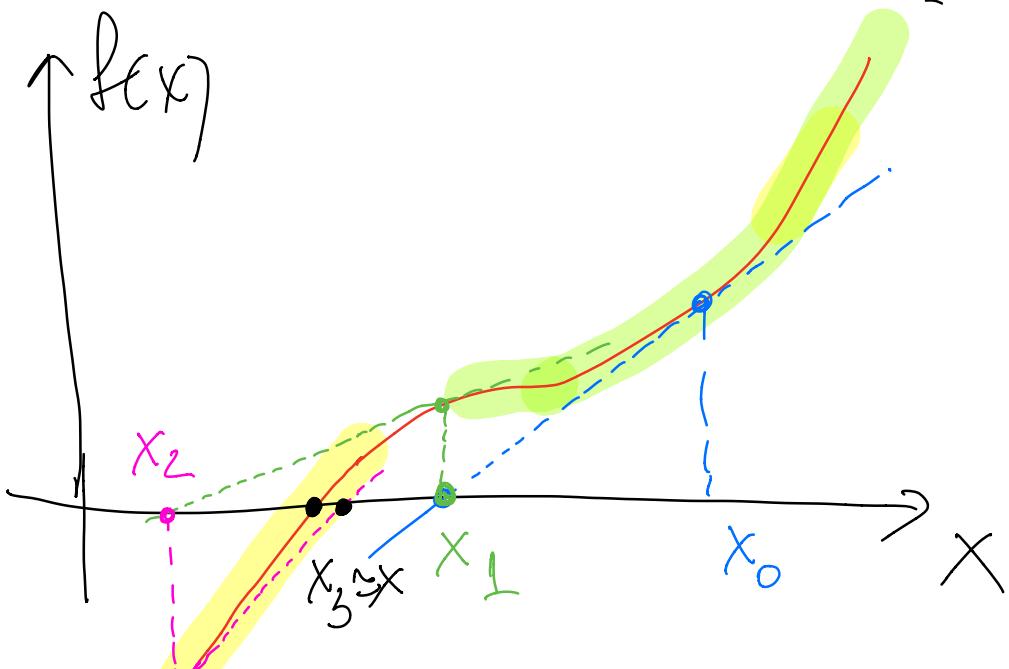
Truncated Taylor series

$$f(x) - p(x) = \frac{f^{(k+1)}(\xi)}{(k+1)!} (x-x_0)^{k+1}$$

Remainder

ξ is between x and x_0

Newton's method



Given x_k , $k = 0, 1, 2, \dots$

Compute x_{k+1}

$$f(x) \approx p_k(x) = f(x_k) + f'(x_k)(x - x_k)$$

$$p_k(x) = 0 \quad \text{solve for } x$$

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$

Newton's method

E.g. $f(x) = x^2 - c$

$$f' = 2x$$

$$x_{k+1} = x_k - \frac{(x_k^2 - c)}{2x_k}$$

$$= \frac{1}{2} \left(x_k + \frac{c}{x_k} \right)$$

Babylonian (\Rightarrow) Newton's

Numerical
Analysis : Include more
terms in Taylor series

$$e_k = x_k - x$$

(sometimes $e_k = |x - x_k|$)

$$f(x) = 0 = f(x_h) + f'(x_h)(x - x_h) \\ + \frac{1}{2} (x - x_h)^2 f''(\xi)$$

ξ is between x and x_h

$$x_k = e_k - x$$

Divide by $f'(x_h)$

$$\left(x_h - \frac{f(x_h)}{f'(x_h)} \right) - x = \frac{1}{2} \underbrace{(x - x_h)^2}_{e_k} \frac{f''(\xi)}{f'(x_h)}$$

$$x_{k+1} - x = e_{k+1}$$

$$e_{k+1} = -\frac{1}{2} \frac{f''(\xi)}{f'(x_k)} e_k^2$$

$$\frac{e_{k+1}}{e_k^2} = -\frac{f''(\xi)}{2f'(x_k)} \rightarrow C$$

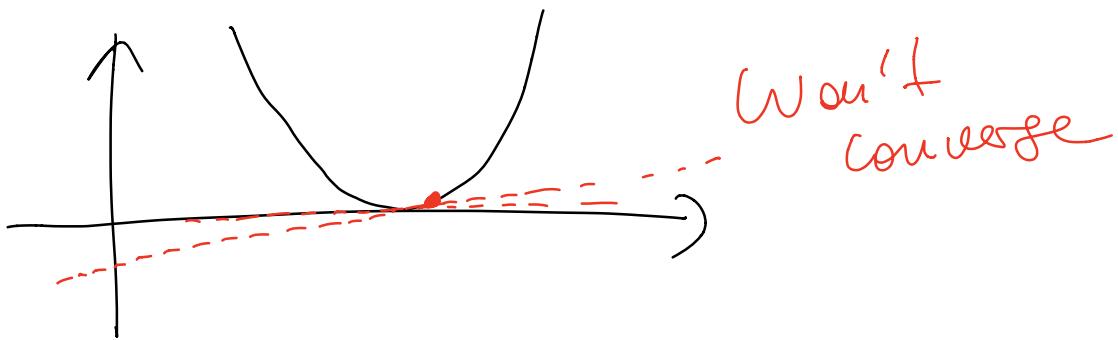
Assume method converges

$$x_k \rightarrow x$$

$$\xi \rightarrow x$$

$$\frac{|e_{k+1}|}{|e_k^2|} \rightarrow \frac{1}{2} \frac{f''(x)}{f'(x)} = C$$

Double root: $f'(x) = 0$



$$|e_{k+1}| < |e_k|$$

We want $|e_{k+1}| \leq \frac{|e_k|}{2}$

$$\left| \frac{f''(\xi)}{2f'(x_k)} \right| e_k^2 < \frac{|e_k|}{2}$$

$$|e_k| < \left| \frac{2f'(x_k)}{2f''(\xi)} \right| < A$$

$$e_0 = x_0 - x$$

$$|x_0 - x| < \left| \frac{2 f'(x_0)}{2 f''(S)} \right|$$

= \Rightarrow

Theorem 1.8

Estimate:

Assumption in Sols:

Practise

$$|x_0 - x| < \left| \frac{f'(x)}{f''(x)} \right|$$

A

$$\left| \frac{f''(y_1)}{f'(y_2)} \right| < A$$

$$\forall y_1, y_2 \in [x-\delta, x+\delta]$$

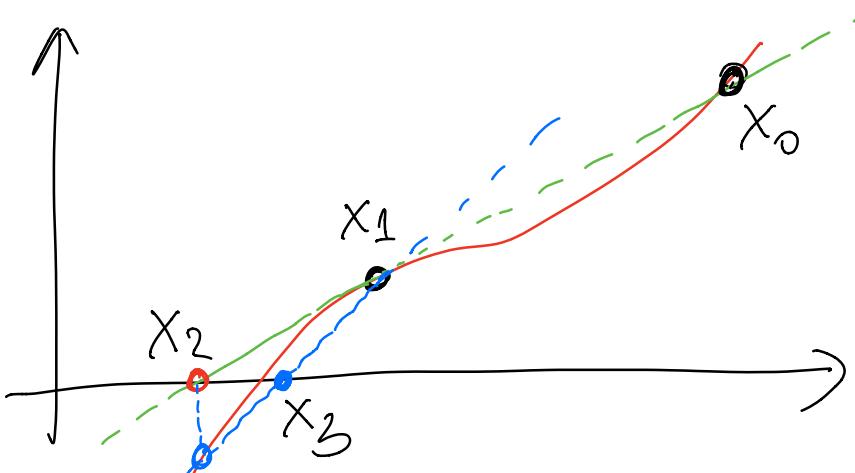
Theorem: If

$$|e_0| = |x_0 - x| < \frac{1}{A}$$

then Newton's method
will converge

Newton's method needs to
be started close to root
and then it will converge
fast

Secant method



$$f(x) = p(x) = f(x_k) + \frac{f(x_{k+1}) - f(x_k)}{(x_{k+1} - x_k)} (x - x_k)$$

$(x_{k+1} - x_k) \rightarrow 0$

if $x = x_k \Rightarrow p(x_k) = f(x_k)$

if $x = x_{k+1} \Rightarrow p(x_{k+1}) = f(x_{k+1})$

Newton $p(x) = f(x_k) + f'(x_k)(x - x_k)$

Conclusion : As $x_k \rightarrow x$

$$x_{k+1} \rightarrow x$$

These two become closer

$$P(x) = 0 \quad \text{solve for } x_{k+2}$$

Secant method

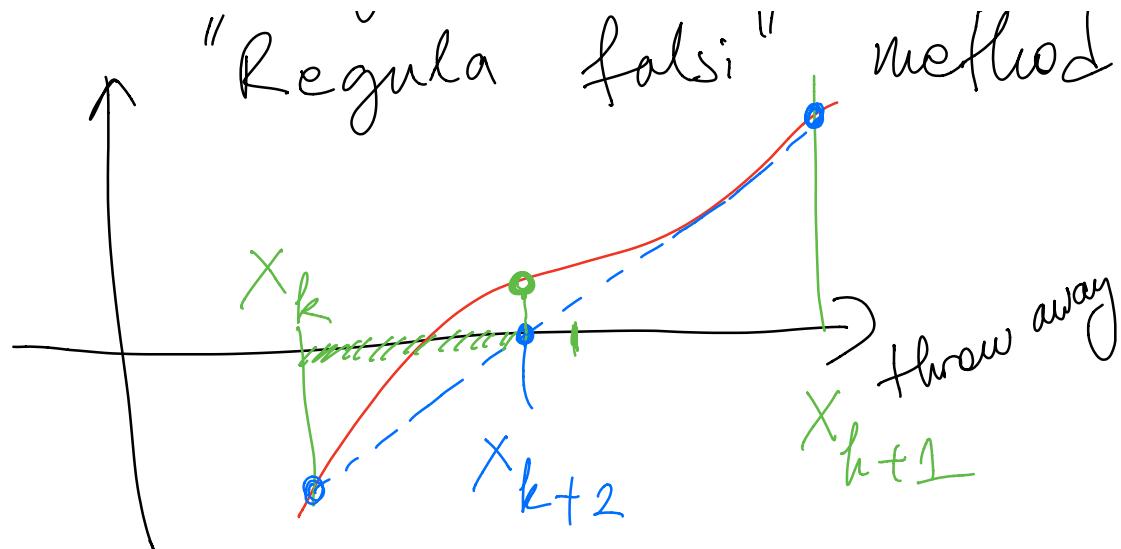
$$x_{k+1} = x_k - \frac{(x_k - x_{k-1})}{\frac{f_k - f_{k-1}}{f_k}}$$

$$f_k = f(x_k)$$

$$\frac{f(x)}{f'(x)}$$

converge to

We can combine
Bisection + secant
or "Safeguarded" secant



Problem 3 in HW 1:
Secant method:

$$|e_{k+1}| \rightarrow \mu |e_k|^q$$

Bisection: $q=1, \mu=1/2$

} Steffensen
Newton: $q=2, \mu = \frac{f''(x)}{2f'(x)}$
(quadratically convergent)

Secant: $q=\frac{1}{2}(1+\sqrt{5}) \approx 1.6$

Fixed-point iteration

$$x_{k+1} = g(x_k)$$

$$x_k \rightarrow x, f(x) = 0$$

$x = g(x)$ ↑
equivalent

Babylonian

$$g(x) = \frac{1}{2}(x + \frac{c}{x})$$

$$x = g(x) \Leftrightarrow x = \sqrt{c}$$

$$\underline{x} = \underline{x^2 - c} + \underline{\frac{x}{x}} = g(x)$$

$$0 = x^2 - c = f(x)$$

$$x_{k+1} = g(x_k)$$

Fixed-point iteration

Example: Relaxation method

$$x_{k+1} = g(x_k) = x_k - \lambda f(x_k)$$
$$g(x) = x - \lambda f(x)$$
$$\lambda \neq 0$$

$$x = g(x) \Rightarrow$$

~~$$x = x - \lambda f(x) = 0$$~~

$$\Rightarrow f(x) = 0$$

Method is consistent

What is a good λ

$$x_{k+1} = x_k - \lambda f(x_k)$$

~~~~~

$$x_{k+1} = g(x_k) \leftarrow \begin{matrix} \text{Taylor series around} \\ x \end{matrix} \quad \begin{matrix} \text{(solution)} \\ \equiv \end{matrix}$$

$$= g(x) + (x_k - x) \overset{|}{g}(\xi)$$

$\xi$  is in-between  $x, x_k$

Express in terms of error

$$\left\{ \begin{array}{l} x_{k+1} = x + e_{k+1} \\ x_k = x + e_k \end{array} \right.$$

$$e_{k+1} = x_{k+1} - x =$$



$$= \tilde{g(x)} + (x_k - x) \tilde{g'(\xi)}$$

$\tilde{x}$   
 [we know  $x = g(x)$ ]  $\Rightarrow f(x) = 0$

$$= \underbrace{(x_k - x)}_{\sim} \tilde{g'(\xi)}$$

$$= e_k \tilde{g'(\xi)}$$

$$\underbrace{|e_{k+1}|}_{\sim} = \underbrace{|g'(\xi)|}_{\sim} \underbrace{|e_k|}_{\sim}$$

We want

$$|e_{k+1}| < |e_k|$$

$$\Rightarrow |g'(\xi)| < 1 \quad (\text{*)})$$

$\xi$  is between  $x$  and  $x_k$

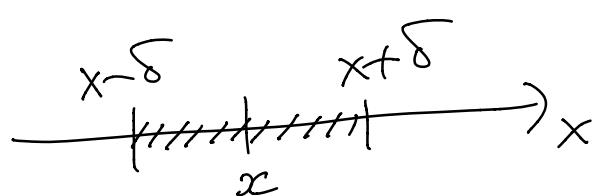
(\*) should be true if  
 $\xi$  is close to root

Theorem:  
if  $g \in C^1$   $\leftarrow$  continuously differentiable

$$|g'(y)| < 1 \quad \forall y \in [x-\delta, x+\delta]$$

where

$$x = g(x)$$

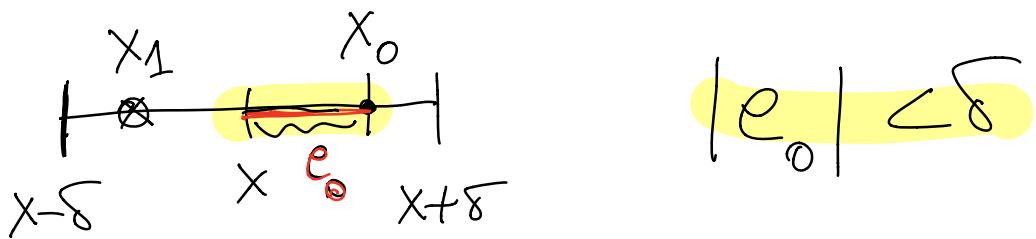


then fixed-point iteration

$$x_{k+1} = g(x_k)$$

will converge to  $x$  if

"Proof":  $x_0 \in [x-\delta, x+\delta]$



$$|e_1| = |g'(z_0)| |e_0|$$

$\uparrow$   
1

$$|e_1| < |e_0| < \delta$$

$$x_1 \in [x-\delta, x+\delta]$$

$$\Rightarrow |e_2| = |g'(z_1)| |e_1|$$

$\uparrow$   
1

$$|e_2| < |e_1| < |e_0|$$

By induction

$$|e_{k+1}| \leq L |e_k|$$

$$0 < L < 1$$

$$L = \max_{y \in [x-\delta, x+\delta]} |g'(y)| < 1$$

$$|e_k| \leq L^k |e_0|$$
$$L^k \rightarrow 0$$

$$\Rightarrow |e_k| \rightarrow 0 \quad \text{Q.E.D.}$$

$$\frac{|e_{k+1}|}{|e_k|} \rightarrow |g'(x)|$$

Since  $x_h \rightarrow x$   
 $\xi \rightarrow x$

Fast convergence we want

$$|g'(x)| << 1$$

Ideally  $g'(x) = 0$   
 convergence is faster than  
 linear  $\Rightarrow$  super linear  
 convergence.

Example : Relaxation

$$x_{k+1} = x_k - \lambda f(x_k)$$

$$g(x) = x - \lambda f(x)$$

$$g'(x) = 1 - \lambda f'(x) = 0$$

$$\lambda = -\frac{1}{f'(x)}$$

Best  $\lambda$ . Fastest  
(super linear) convergence for

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x)}$$

*Doesn't work* we don't know  $x$

c.f. Newton

$$\Rightarrow x_k - \frac{f(x_k)}{f'(x_k)}$$

For Newton's method

$$g(x) = x - \frac{f(x)}{f'(x)}$$

what is  $g'(x)$  at the root  $x$  that satisfies  $x=g(x)$

$$\begin{aligned} g'(x) &= 1 - \frac{f'(x)}{f(x)} + \frac{f(x)f''(x)}{(f'(x))^2} \\ &= \frac{f(x)f''(x)}{(f'(x))^2} \end{aligned}$$

$$\text{If } f(x) = 0 \Rightarrow g'(x) = 0$$

Consistent with Newton's method converging quadratically.