



# Stark 101: Part 3

**FRI Commitment**

# Recap

**Goal : prove a statement on FibonacciSq**

- Trace in 1023 points
- Create *Trace* polynomial (Lagrange interpolation)
- Evaluate and commit on a larger domain

# Recap

- 3 constraints on  $f(x)$ :

$$f(x) - 1 = 0 \text{ , for } x = 1$$

...

- 3 rational functions from the constraints:

$$p_0(x) = \frac{f(x) - 1}{x - g^0}$$

...

# Recap

- Composition **P**olynomial:

$$CP(x) = \alpha_0 \cdot p_0(x) + \alpha_1 \cdot p_1(x) + \alpha_2 \cdot p_2(x)$$

- Prover commits on CP
- Goal - show that CP is a **polynomial**
- CP is a **polynomial**  $\rightarrow$  All constraints satisfied

# What Will We Do?

Goal:

Prove that CP is a **polynomial**

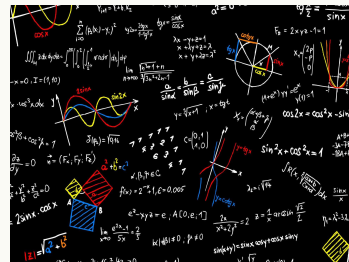


Instead:

Prove that CP is **close** to a **polynomial** of **low degree**

What is close?

What is low degree?

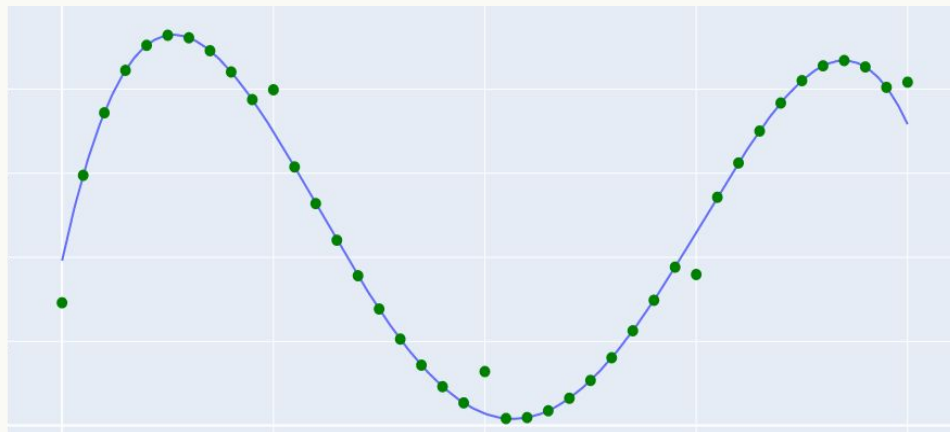


# Proximity to Polynomials

## Distance (def):

Distance between a function  $f: D \rightarrow F$  to a polynomial  $p$ :

$$D(f,p) := \# \text{ points } x \in D \text{ such that } f(x) \neq p(x)$$



$$D(\mathbf{f}, \mathbf{p}) = 5$$

# Proximity to Polynomials

## Distance (def):

Distance between a function  $f: D \rightarrow F$  to a polynomial  $p$ :

$$D(f,p) := \# \text{ points } x \in D \text{ such that } f(x) \neq p(x)$$

## Proximity

A function  $f: D \rightarrow F$  is **close** to a polynomial  $p$  if:  $D(f,p)$  is **small**

FRI

# Fast Reed-Solomon Interactive Oracle Proofs of Proximity

By Ben-Sasson, E., Bentov, I., Horesh, Y., & Riabzev, M.

<https://eccc.weizmann.ac.il/report/2017/134/>



# FRI

Prover convinces verifier:

**“The commitment is close to a low degree polynomial”**

# Without Using FRI



# FRI Operator - The Reduction

# FRI Operator

## Goal:

Prove that a function is close to a polynomial of a bounded degree  $D$

Applying the FRI operator

## New Goal:

Prove that a **new** function is close to a **new** polynomial

Half of the domain size

Degree bound  $D/2$

## FRI Operator - Example

# Before applying FRI operator

- Prove:

A function is close to a polynomial of a bounded degree **1024**

where domain size = **8192**

## FRI Operator - Example

# ~~Before~~ After applying FRI operator

- Prove:

A function is close to a polynomial of a bounded degree ~~1024~~ 512

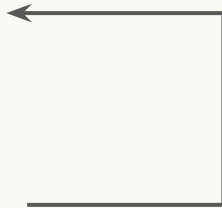
where domain size = ~~8192~~ 4096



# FRI Overview

# FRI - The Protocol

- Receives random  $\beta$
- Computes the next polynomial
- Commit
- Lastly the prover sends the constant



Do it repeatedly until:

$$\mathbf{deg(poly) < 1}$$

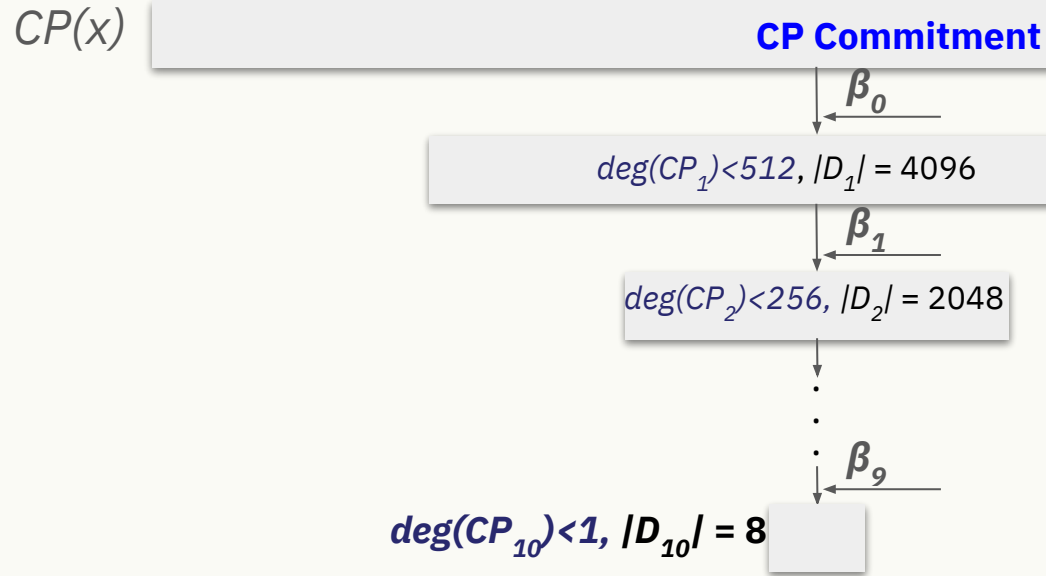
where

**domain size is 8**



# FRI - Illustration

Showing that  $\deg(CP) < 1024$ ,  $|D| = 8192$



# Deep Into FRI



# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = g(x^2) + xh(x^2)$$

- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = g(y) + \beta h(y)$$

- Example:

$$P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3$$

The diagram illustrates the decomposition of the polynomial  $P_0(x)$  into two parts:  $g(x^2)$  and  $xh(x^2)$ . The polynomial is written as  $P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3$ . Arrows point from the even-powered terms to  $g(x^2)$  and from the odd-powered terms to  $xh(x^2)$ .

Term in $P_0(x)$	Term in $g(x^2)$	Term in $xh(x^2)$
$5x^5$	$5x^5$	
$3x^4$	$3x^4$	
$7x^3$		$7x^3$
$2x^2$	$2x^2$	
$x$		$x$
$3$	$3$	

# FRI Operator - How Does it Work?

- Split to even and odd powers

$$P_0(x) = g(x^2) + xh(x^2)$$

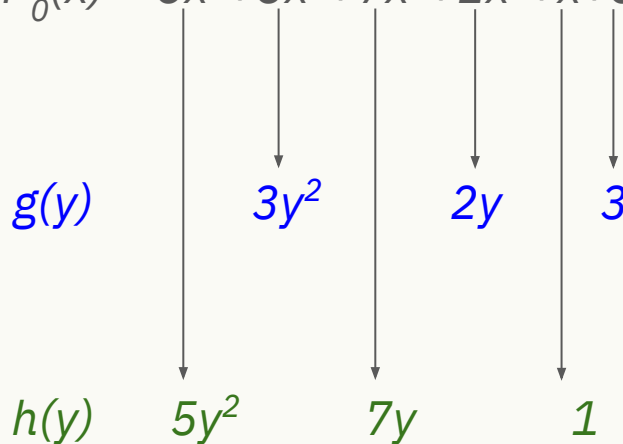
- Get a random  $\beta$

- Consider the new function:

$$P_1(y) = g(y) + \beta h(y)$$

- Example:

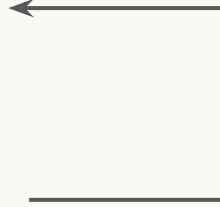
$$P_0(x) = 5x^5 + 3x^4 + 7x^3 + 2x^2 + x + 3$$



- $$P_1(y) = 3y^2 + 2y + 3 + \beta(5y^2 + 7y + 1)$$
$$= (3 + 5\beta)y^2 + (2 + 7\beta)y + 3 + \beta$$

# FRI - The Protocol - Reminder

- Receives random  $\beta$
- Computes the next polynomial
- Commit
- Lastly the prover sends the constant



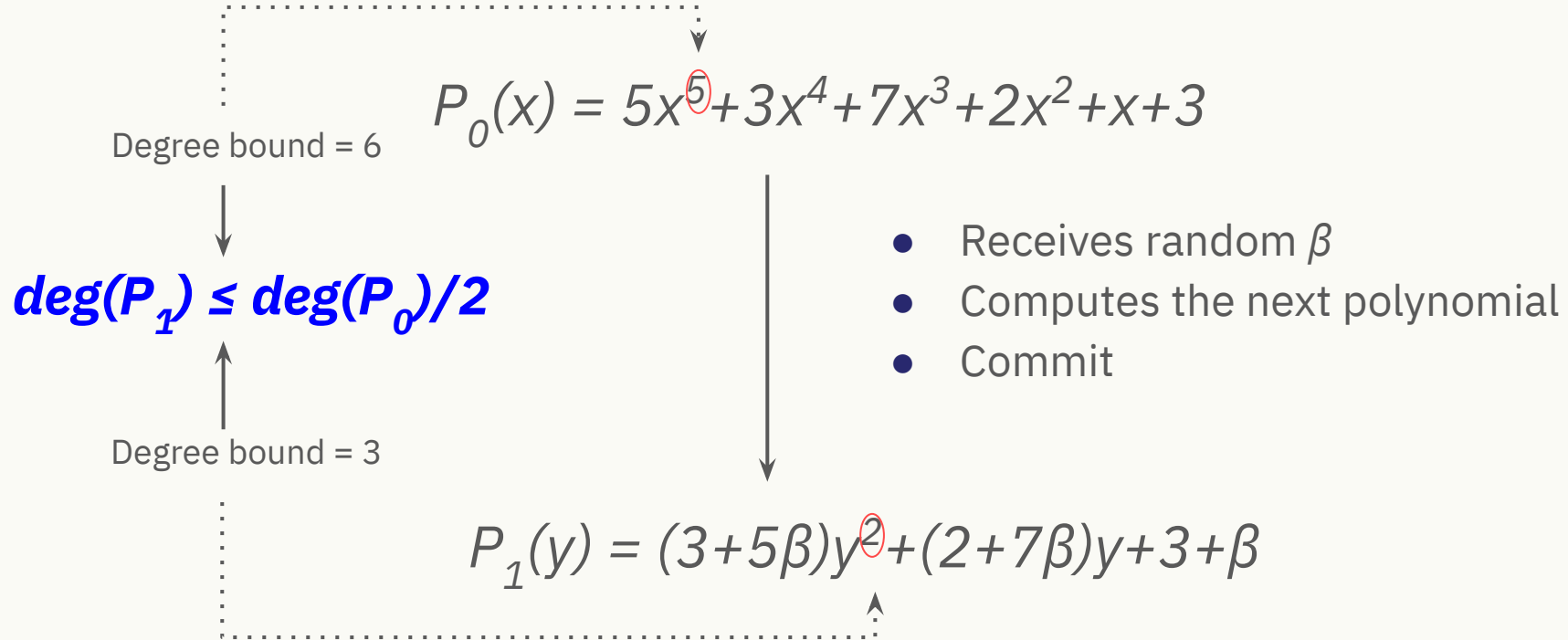
Do it repeatedly until:

$$\mathbf{deg(poly) < 1}$$

where

**domain size is 8**

# FRI - The Protocol - A Single Step



Thank you