# **CSE160: Computer Networks**

### Lecture #03 - Bits and Bandwidth

2020-09-03



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### Last Time ...

Protocols, layering and reference models

Application		"Chrome"	
Presentation		HTTP	
Session			
Transport	<b>─</b>	TCP	
Network		IP	
Link		Eth own of	
Physical		Ethernet	



Sample Protocol Stack



### **This Lecture**

 Focus: <u>How do we send a</u> <u>message across a wire?</u>

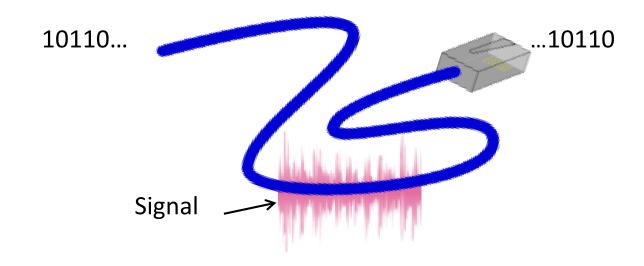
- The physical / link layers:
  - 1. Model of a link
  - 2. Different kinds of media, signals
  - 3. Encoding bits with signals
  - 4. Fundamental Limits

Application
Presentation
Session
Transport
Network
Data Link
Physical



## **Scope of the Physical Layer**

- Concerns how signals are used to transfer message bits over a link
  - Wires etc. carry <u>analog signals</u>
  - We want to send <u>digital bits</u>





### 1. Simple Link Model

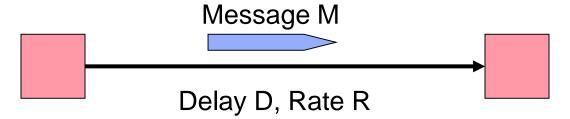


- Abstract model is typically all we will need
  - What goes in comes out altered by the model
- We'll end with an abstraction of a physical channel
  - Rate (or bandwidth, capacity, speed) in bits/seconds
  - Delay in seconds, related to length
- Other parameters that are important:
  - The kind and frequency of errors
  - Whether the media is broadcast or not



## **Message Latency**

How long does it take to send a message?



- Latency is the delay to send a message over a link
  - Transmission delay = message (bits) / rate (bps)
    - Time to put M-bit message "on the wire" (M/R seconds)
  - Propagation delay = length / speed of light in media
    - Time for bits to propagate across the wire (D seconds)
- Combining the two terms we have: L = M/R + D
- Later we will see queuing delay ...



## Relationships

- Latency (sec) =
  - Transmit (sec) + Propagation (sec) + Queue (sec)
- Transmit Time (sec) =
  - Message Size (bits) / Bandwidth (bits/sec)= M/R seconds
- Propagation Delay (sec) =
  - Length (mts) / Speed of signals (mts/sec) =L/(2/3c) = D seconds



### **Metric Units**

The main prefixes we use:

Prefix	Exp.	Prefix	Exp.
K(ilo)	10 <sup>3</sup>	m(illi)	10-3
M(ega)	106	μ(micro)	10 <sup>-6</sup>
G(iga)	109	n(ano)	10-9

- Use powers of 10 for rates, 2 for storage
  - $1 \text{ Kbps} = 1,000 \text{ bps}, 1 \text{ KB} = 2^{10} = 1024 \text{ bytes}$
  - $-1 \text{ Mbps} = 1,000,000 \text{ bps}, 1 \text{ MB} = 2^{20} = 1,048,576 \text{ bytes}$



"B" is for bytes, "b" is for bits, 1B = 8b

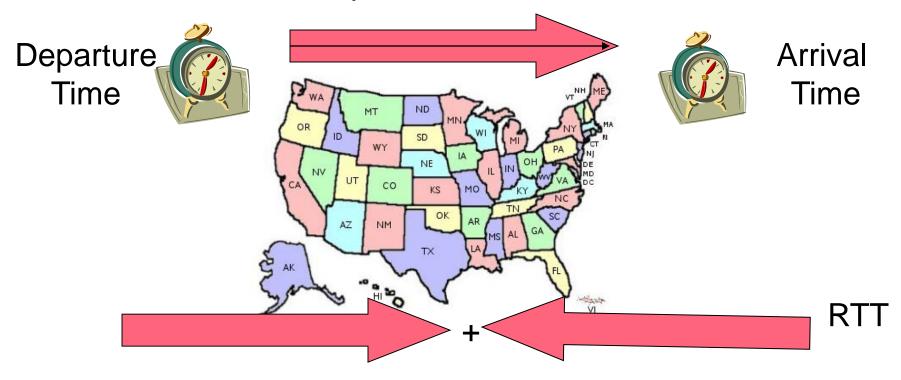
### **One-way Latency**

- "Dialup" with a telephone modem:
  - -D = 10ms, R = 56Kbps, M = 1024 bytes
  - Latency =  $(1024 \times 8)/(56 \times 1,000)$  sec + 10ms = 153ms!
- Cross-country with T3 (45Mbps) line:
  - -D = 50 ms, R = 45 Mbps, M = 1024 bytes
  - Latency =  $(1024 \times 8) / (45 \times 1,000,000)$  sec 50ms = 50ms!
- Either a slow link or long wire makes for large latency
  - Often, one delay components dominates



## **Latency and RTT**

- Latency is typically the one way delay over a link
  - Arrival Time Departure Time



- The round trip time (RTT) is twice the one way delay
  - Measure of how long to signal and get a response

## **Throughput**

- Measure of system's ability to "pump out" data
  - NOT the same as bandwidth
- Throughput = Transfer Size / Transfer Time
  - Eg, "I transferred 1000 bytes in 1 second on a 100Mb/s link"
    - BW?
    - Throughput?
- Transfer Time = SUM OF
  - Time to get started shipping the bits
  - Time to ship the bits
  - Time to get stopped shipping the bits
- Throughput includes all protocols overheads



## Messages Occupy "Space" On the Wire

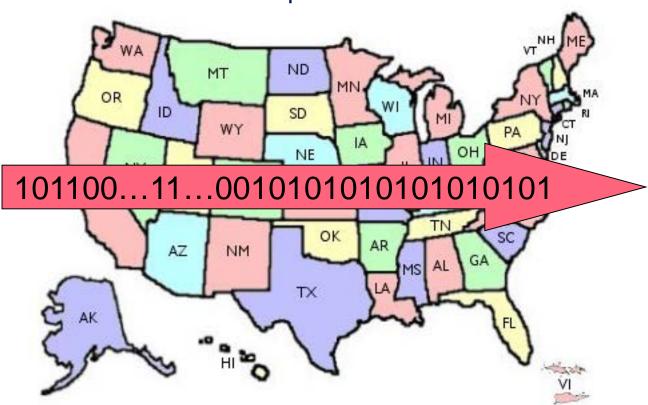
- Consider a 1b/s network.
  - How much space does 1 byte take?
- Suppose latency is 16 seconds.
  - How many bits can the network "store"
  - This is the BANDWIDTH-DELAY product
  - Measure of "data in flight."
  - 1b/s \* 16s = 16b
- Tells us how much data can be sent before a receiver sees any of it.
  - Twice B.D. tells us how much data we could send before hearing back from the receiver something related to the first bit sent.
- SELLA CONTROL OF THE PROPERTY OF THE PROPERTY

– Implications?

### A More Realistic Example

Fiber at home, cross country

$$BD = 50ms * 45Mbps = 2.25 * 10^6 = 280KB$$



- That's quite a lot of data in the network
  - Like a small novel

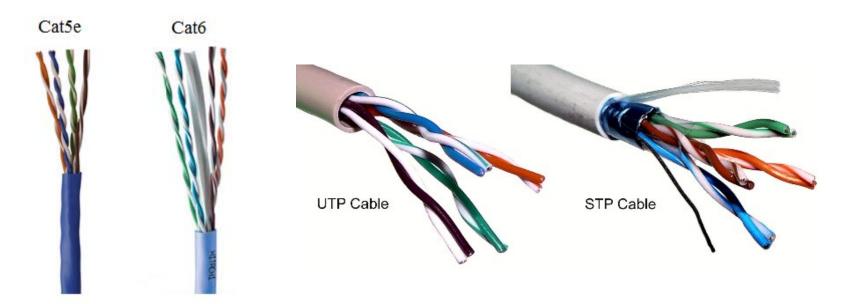
## 2. Types of Media

- Media propagate <u>signals</u> that carry <u>bits</u> of information
- We'll look at some common types:
  - Wires
  - Fiber (fiber optic cables)
  - Wireless



### **Wires – Twisted Pair**

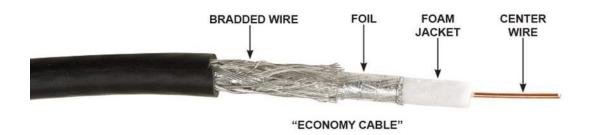
- Very common, used in LANs and telephone lines
  - Four twisted pairs
  - Twists reduce radiated signal
  - Category 5e or 6 (bandwidth and insertion loss)
  - Unshielded (UTP) or Shielded (STP)

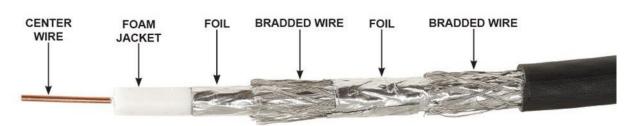




### Wires - Coaxial Cable

Also common. Better shielding for better performance





"QUAD-SHIELDED CABLE"

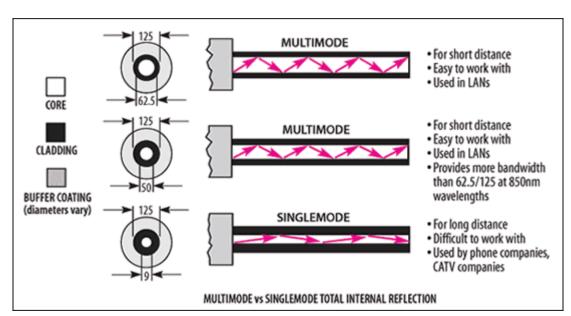
Other kinds of wires too: e.g., electrical power

### **Fiber**

- Long, thin, pure strand of glass
  - light propagated with total internal reflection
  - enormous bandwidth available (terabits)
  - Multi-mode allows many different paths, limited by dispersion
  - Chromatic dispersion if multiple frequencies

Light source (LED, laser)

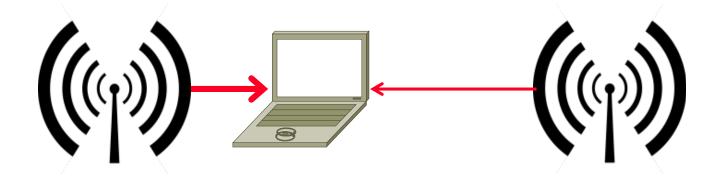
Light detector (photodiode)





### **Wireless**

- Sender radiates signal over a region
  - In many directions, unlike a wire, to potentially many receivers
  - Nearby signals (same freq.) <u>interfere</u> at a receiver; need to coordinate use





### UNITED

### STATES

### **FREQUENCY**

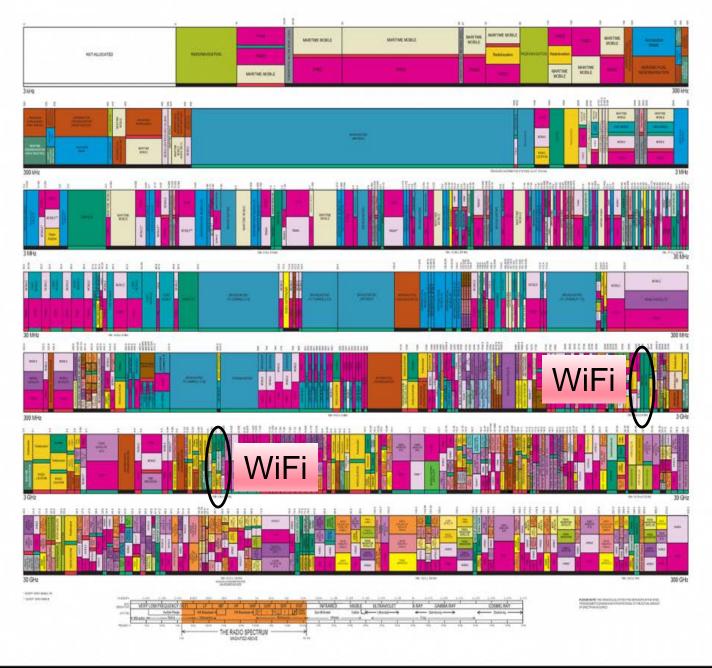
### **ALLOCATIONS**

#### THE RADIO SPECTRUM



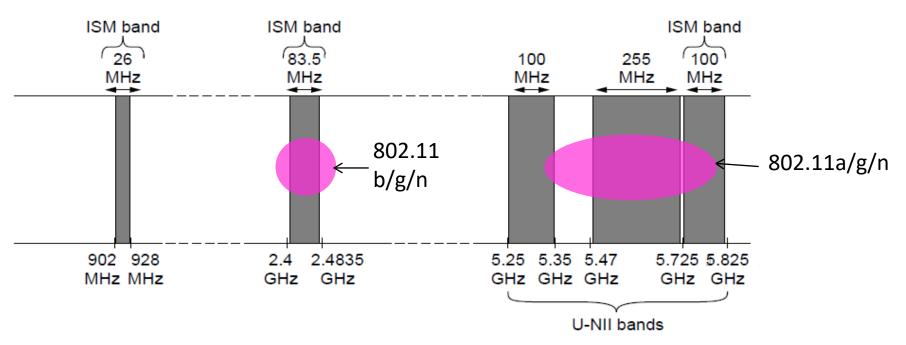
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Postary.	PUED	State Little
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## Wireless (2)

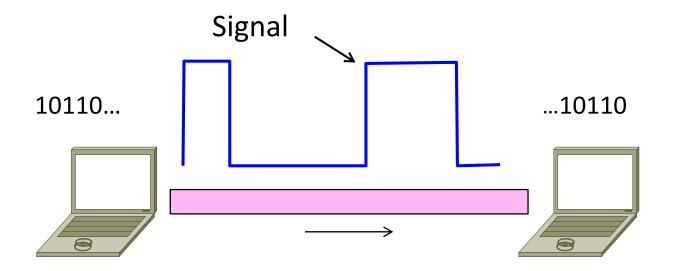
 Microwave, e.g., 3G, and unlicensed (ISM – Industrial, Scientific and Medical) frequencies, e.g. WiFi, are widely used for computer networking





### 3. Signals

- Analog signals encode digital bits
- We want to know what happens as signals propagate over media

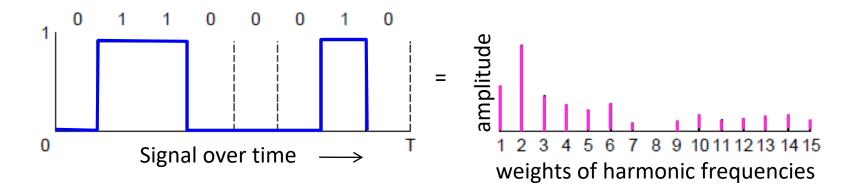




## **Frequency Representation**

 A signal over time can be represented by its frequency components (called Fourier analysis)

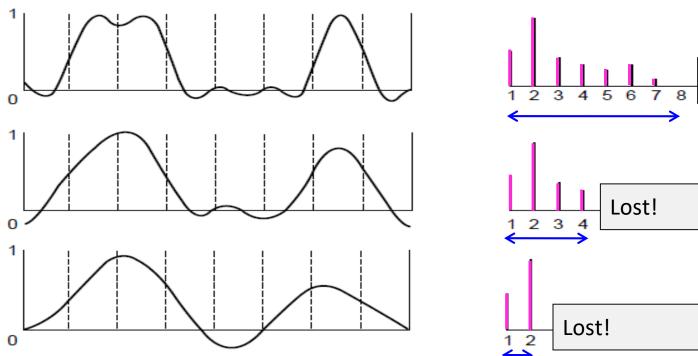
$$g(t) = \frac{1}{2}c + \sum_{n=1}^{\infty} a_n \sin(2\pi n f t) + \sum_{n=1}^{\infty} b_n \cos(2\pi n f t)$$

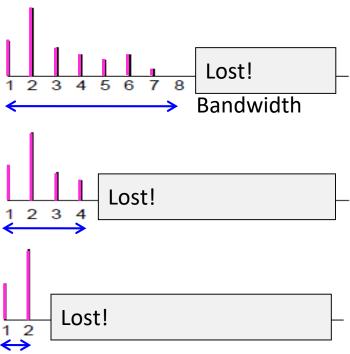




### **Effect of Less Bandwidth**

 Fewer frequencies (=less bandwidth) degrades signal (less rapid transitions)





Which one is good for data communication?

## Signals over a Wire

- What happens to a signal as it passes over a wire?
  - 1. The signal is delayed (propagates at 2/3 c)
  - 2. The signal is attenuated (goes for m to km)
  - 3. Frequencies above a cutoff are highly attenuated
  - 4. Noise is added to the signal (later, causes errors)

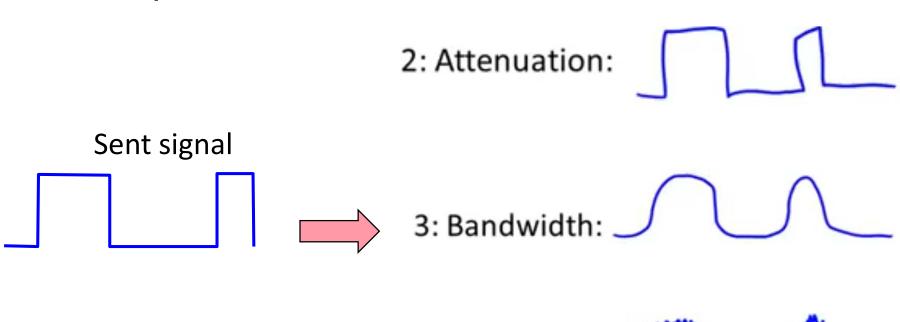
EE: Bandwidth = width of frequency band, measured in Hz

CS: Bandwidth = information carrying capacity, in bits/sec



# Signals over a Wire (2)

Example:

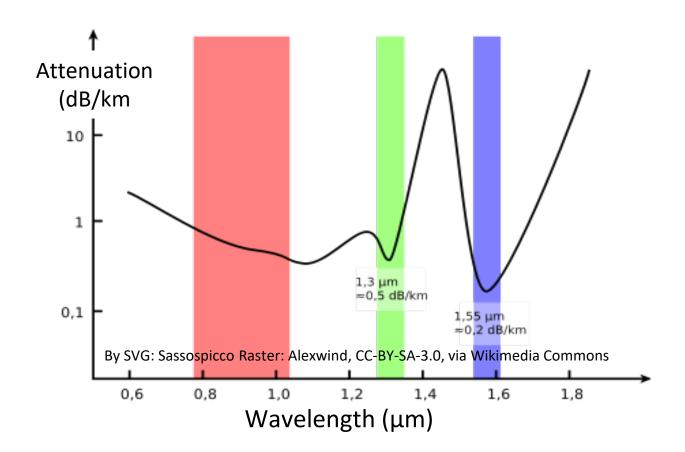


4: Noise:



# Signals over Fiber

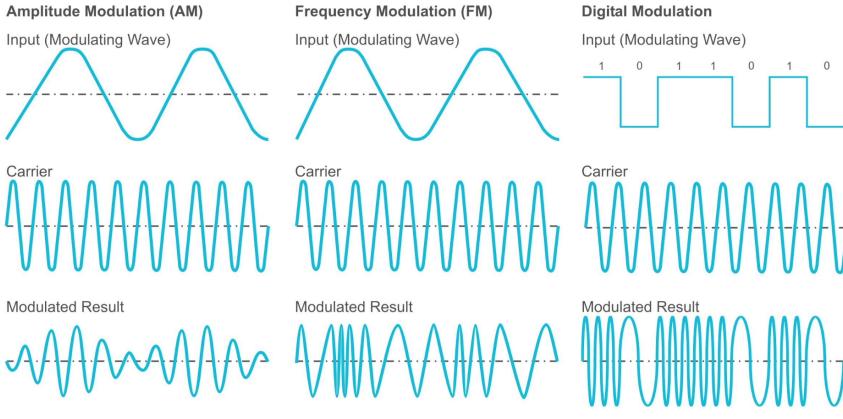
- Light propagates with very low lows in three very wide frequency bands
  - Use a carrier to send information





## **Signals over Wireless**

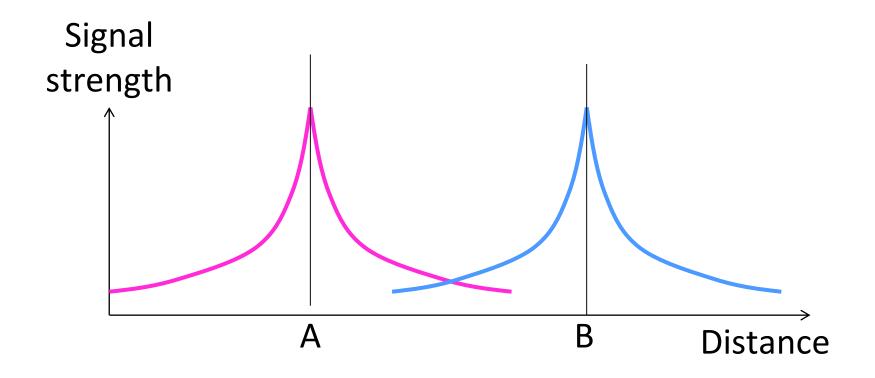
 Signals transmitted on a carrier frequency, like fiber (more later)





# Signals over Wireless (2)

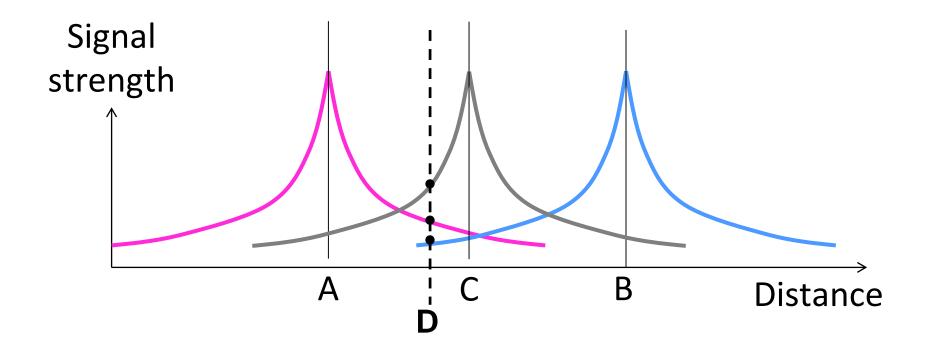
 Travel at speed of light, spread out and attenuate faster than 1/dist<sup>2</sup>





# Signals over Wireless (3)

Multiple signals on the same frequency interfere at a receiver

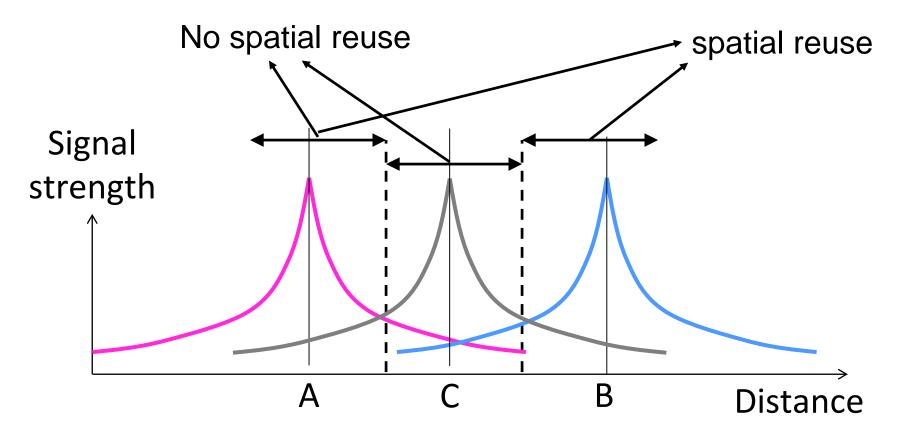


D sees = strong C + weak A, B



# Signals over Wireless (4)

 Interference leads to notion of <u>spatial reuse</u> (of same frequency)





# Signals over a Wireless (5)

- Various other effects too!
  - Wireless propagation is complex, depends on environment

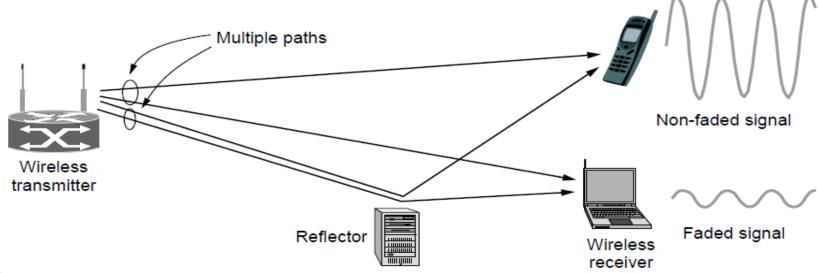
- Some key effects are highly frequency dependent,
  - E.g., <u>multipath</u> at microwave frequencies

 More on this later when discussing Medium Access Control



## Wireless Multipath

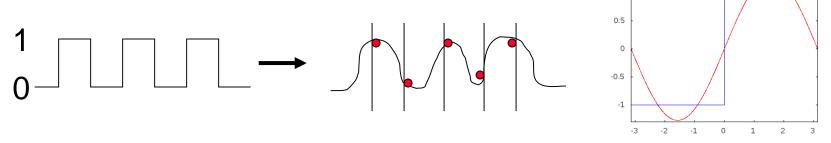
- Signals bounce off objects and take multiple paths
  - Some frequencies attenuated at receiver, varies with location
  - Messes up signal; handled with sophisticated methods





### **Modulation**

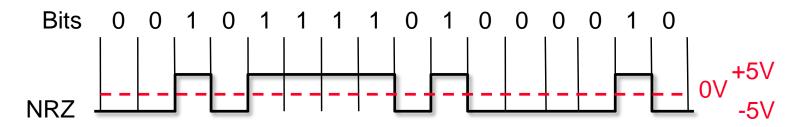
 Generate analog waveform (e.g., voltage) from digital data at transmitter and sample to recover at receiver



- We send/recover symbols that are mapped to bits
  - Signal transition rate = baud rate vs bit rate
- This is baseband transmission ... for more details you need a signals course!

### **Simple Modulation**

- What is the simplest encoding?
  - Use high voltage for 1, zero voltage for zero
  - What is the problem?
- To solve DC bias, NRZ (Non-return to zero)
  - Use high/low voltages, e.g., high=1, low=0
  - Use +/- voltages, e.g., +5V = 1, -5V = 0



- Variation, NRZI (NRZ, invert on 1)
  - Use transition for 1s, no transition for 0s
  - How do we sample a long sequence of 0s?

## **Clock Recovery**

- How do we distinguish consecutive 0s or 1s?
- If we sample at the wrong time we get garbage
- If Tx and Rx have exact clocks no problem
  - But in practice they drift slowly
- This is the Clock Recovery problem
- Possible solutions:
  - Send separate clock signal
    - Expensive
  - Keep messages short
    - Limits data rate
  - Embed clock signal in data signal
    - Other codes

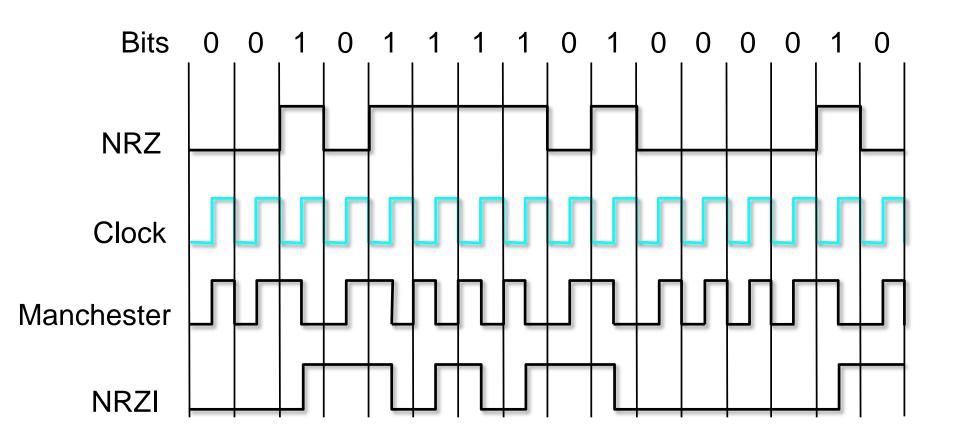


## **Manchester Coding**

- Make transition in the middle of every bit period
  - Low-to-high is 0; high-to-low is 1
  - Signal rate is twice the bit rate
  - Used on 10 Mbps Ethernet
- Advantage:
  - self-clocking: clock is embedded in signal, and we re-sync with a phase-locked loop (PLL) every bit
- Disadvantage:
- STATE OF STA

50% efficiency

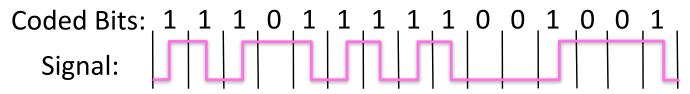
# **Coding Examples**





#### 4B/5B and 8B/10B Codes

- We want transitions \*and\* efficiency... how?
- Solution: map data bits (which may lack transitions) into code bits (which are guaranteed to have them)
- 4B/5B code:
  - $-0000 \rightarrow 11110,0001 \rightarrow 01001,...1111 \rightarrow 11101$
  - Never more than three consecutive 0s back-to-back using NRZI or other codes (e.g. MLT-3)
  - 80% efficiency (used in 100Mbps Ethernet and FDDI)
  - Message bits: 1111 0000 0001



- 8B/10B encoding used for 1GB Ethernet (with NRZ)
  - never more than 5 consecutive 0s or 1s (clock recovery)
  - diff. bw counts of 1s/0s every 20 bits < 2 (DC-balance)</li>



Cerpa, Fall 2020 © UCM

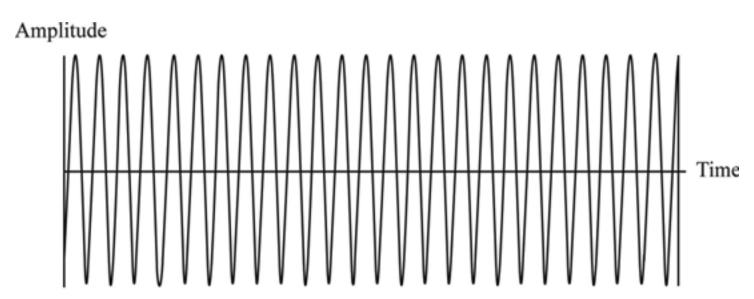
#### **Passband Modulation**

- What we have seen so far is <u>baseband</u> modulation for wires
  - Signal is sent directly on a wire
- These signals do not propagate well on fiber / wireless
  - Need to send at higher frequencies
- Passband modulation carries a signal by modulating a carrier



### **Passband Modulation (2)**

 Carrier is simply a signal oscillating at a desired frequency:



- We can modulate it by changing:
  - Amplitude, frequency, or phase



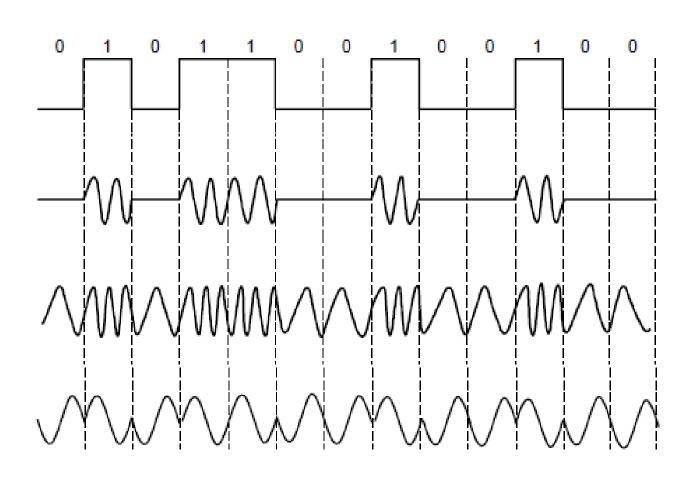
### **Passband Modulation (3)**

NRZ signal of bits

Amplitude shift keying

Frequency shift keying

Phase shift keying





#### 4. Fundamental Limits

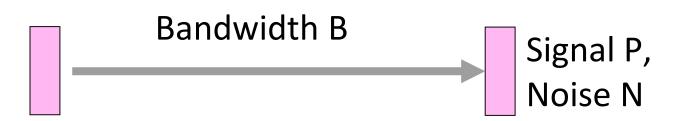
- How rapidly can we send information over a link?
  - Nyquist limit (~1924)
  - Shannon capacity (1948)

 Practical systems are devised to approach these limits



### **Key Channel Properties**

- The bandwidth (B), signal strength (S), and noise strength (N)
  - All measured at the receiver
  - B limits the rate of transitions
  - S and N limit how many signal levels we can distinguish





### **Nyquist Limit**

- The maximum symbol rate is 2B
  - Symbol is a waveform that it's used to represent information (bits)
    - It may represent 1 bit, more than 1bit if you have multiple signal levels

1010101010101010101



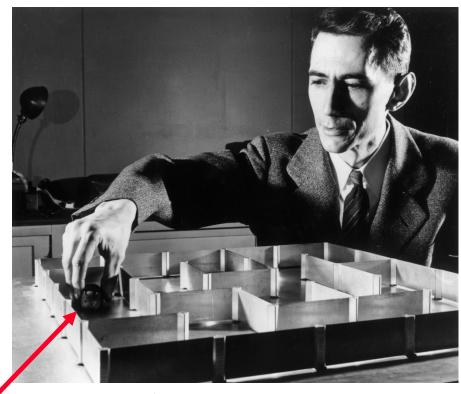
 Thus if there are V signal levels, ignoring noise, the maximum bit rate is:



 $R = 2B \log_2 V \text{ bits/sec}$ 

### **Claude Shannon (1916-2001)**

- Father of information theory
  - "A Mathematical Theory of Communication",
    1948
- Fundamental contributions to digital computers, security, and communications



Credit: Courtesy MIT Museum



Electromechanical mouse that "solves" mazes!

#### **Shannon Limit**

- How many levels we can distinguish depends on S/N
  - Or SNR, the <u>Signal-to-Noise Ratio</u>
  - Note noise is random, hence some errors
- SNR given on a log-scale in deciBels:
  - $-SNR_{dB} = 10log_{10}(S/N)$

S/N	dB
1000	30
100	20
10	10
2	3



### **Shannon Limit (2)**

 Shannon limit is for capacity (C), the maximum information carrying rate of the channel:

$$C = B log_2(1 + S/N) bits/sec$$

- Shannon showed that you can send information up to that rate, but not higher
- Revolutionary at the time: there were many errors in the channels and the only way to cope with them was increasing signal strength
- Shannon showed us that there might be coding schemes that would allow us to transmit the signal



## Putting it all together – DSL

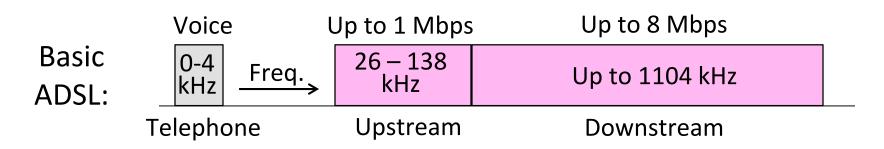
- DSL (Digital Subscriber Line, see §2.1.1) is widely used for broadband; many variants offer 10s of Mbps
- Reuses twisted pair telephone line to the home; it has up to ~2 MHz of bandwidth but uses only the lowest ~4 kHz





### **DSL (2)**

- DSL uses passband modulation
  - Separate bands for upstream and downstream (larger)
  - Modulation called QAM varies both amplitude and phase
  - High SNR, up to 15 bits/symbol, low SNR only 1 bit/symbol





### Wired/Wireless Perspective

- Wires, and Fiber:
  - Engineer link to have requisite SNR and B
  - Can fix data rate

Engineer SNR for data rate

- Wireless
  - Given B, but SNR varies greatly, e.g. up to 60 dB! (that's 1M times)
  - Can't design for worst case, must adapt data rate



Adapt data rate to SNR

### **Key Concepts**

- We typically model links in terms of bandwidth and delay, from which we can calculate message latency
- Different media have different properties that affect their performance as links
- We need to encode bits into signals so that we can recover them at the other end of the channel
- Passband modulation allows us to transmit baseline signals in other higher frequencies for fiber and wireless
- Shannon showed us the limit of what is possible to transmit on any channel

