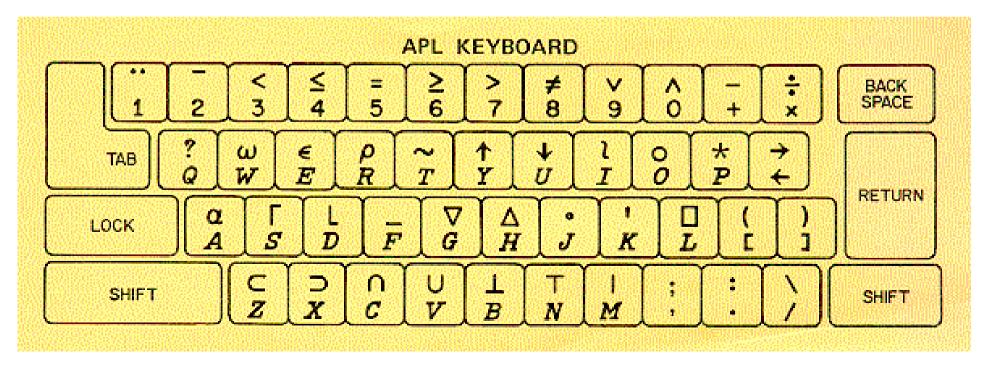
APL: The Greatest Programming Language You Never Heard Of





Ken Iverson



Prime Numbers

$$(2 = + \neq 0 = (\iota X) \circ . | \iota X) / \iota X$$

lota

```
1 X
1 2 3 4 5 6 7 8 9 10
```

Rho

```
3 3 ρ ι 9
1 2 3
4 5 6
7 9 9
```

Select

0 1 0 1 0 1 0 1 0 / ι 8 1 3 5 7

Outer Product

```
3 4 5 °.+ 1 2 3 4
4 5 6 7
5 6 7 8
6 7 8 9
```

Identity Matrix

```
(\iota X) \circ = \iota X
1 0 0 0 0 0 0 0 0
0 1 0 0 0 0 0 0 0
0 0 1 0 0 0 0 0 0
0 0 0 1 0 0 0 0 0
000010000
 0 0 0 0 1 0 0 0
   0 0 0 0 1 0 0 0
   0 0 0 0 0 1 0 0
   0 0 0 0 0 0 1 0
   0 0 0 0 0 0 0 1
```

Residue Matrix

```
(\iota X) \circ . | \iota X
0 \ 0 \ 0 \ 0 \ 0 \ 0
1 0 1 0 1 0 1 0 1 0
1 2 0 1 2 0 1 2 0 1
1 2 3 0 1 2 3 0 1 2
1 2 3 4 0 1 2 3 4 0
1 2 3 4 5 0 1 2 3 4
1 2 3 4 5 6 0 1 2 3
1 2 3 4 5 6 7 0 1 2
1 2 3 4 5 6 7 8 0
1 2 3 4 5 6 7 8 9 0
```

Integral Divisibility Matrix

```
0 = (\iota X) \circ . | \iota X
1 1 1 1 1 1 1 1 1 1
0 1 0 1 0 1 0 1 0 1
0 0 1 0 0 1 0 0 1 0
0 0 0 1 0 0 0 1 0 0
0 0 0 0 1 0 0 0 0 1
  0 0 0 0 1 0 0 0 0
    0 0 0 0 1 0 0 0
        0 0 0 1 0 0
    0 0 0 0 0 0 1 0
   0 0 0 0 0 0 0 1
```

Number of Integral Divisors

$$+\neq 0 = (\iota X) \circ . | \iota X$$

1 2 2 3 2 4 2 4 3 4

Exactly Two Integral Divisors

$$2 = + \neq 0 = (\iota X) \circ . | \iota X$$

0 1 1 0 1 0 1 0 0 0

Prime Numbers

$$(2 = + \neq 0 = (\iota X) \circ . | \iota X) / \iota X$$

2 3 5 7

More APL Examples

Leap Year Test

$$(0 = 400 \mid X) \lor (0 \neq 100 \mid X)^0 = 4 \mid X$$

Test for Duplicate Elements

$$^{\prime\prime}(X\iota X) = \iota \rho X$$

Standard Deviation

$$((+/(X - (+/X) \div \rho X) * 2) \div \rho X) * .5$$

lota

```
(define iota
  (lambda (n)
    (letrec
       ((loop
         (lambda (n acc)
           (if (= n 0))
               acc
               (loop (sub1 n)
                      (cons n acc))))))
       (loop n '())))
> (iota 8)
(1 2 3 4 5 6 7 8)
```

Select

Select Explained

```
> (map cons '(1 2 3) '(a b c))
((1 . a) (2 . b) (3 . c))
> (map (lambda (x) (pred (car x))))
'((1 . a) (2 . b) (3 . c))
((#t . a) (#f . b) (#t . c))
> (filter (lambda (x) (pred (car x)))
'((1 . a) (2 . b) (3 . c)))
((1 . a) (3 . c))
> (map cdr '((1 . a) (3 . c)))
(a c)
```

Tally



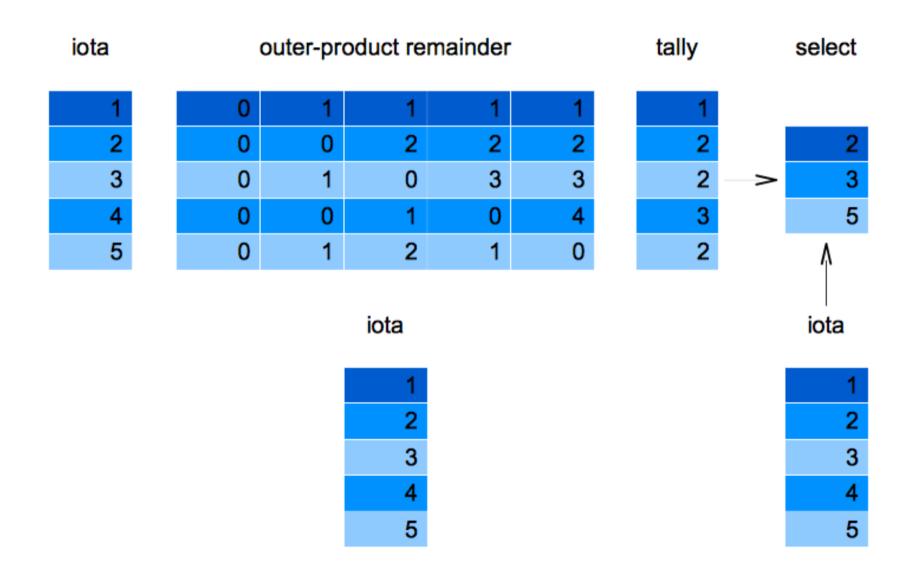
Outer Product



Prime Numbers

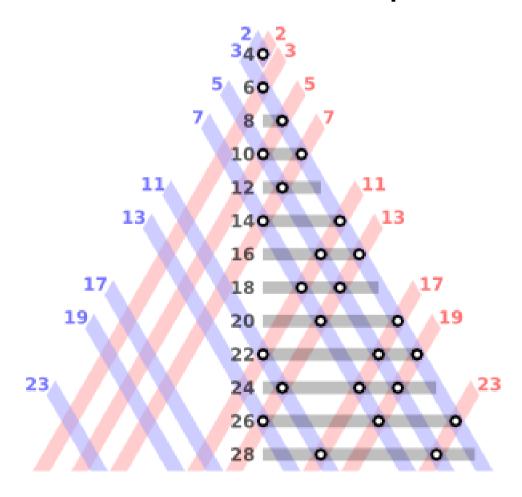
```
(define primes
  (lambda (n)
    (let ((ls (iota n)))
       ((select (lambda (x) (= x 2))
        (map (tally zero?)
             ((outer-product remainder) ls ls))
       ls))))
> (primes 10)
(2 \ 3 \ 5 \ 7)
```

Prime Numbers Explained



Goldbach's Conjecture

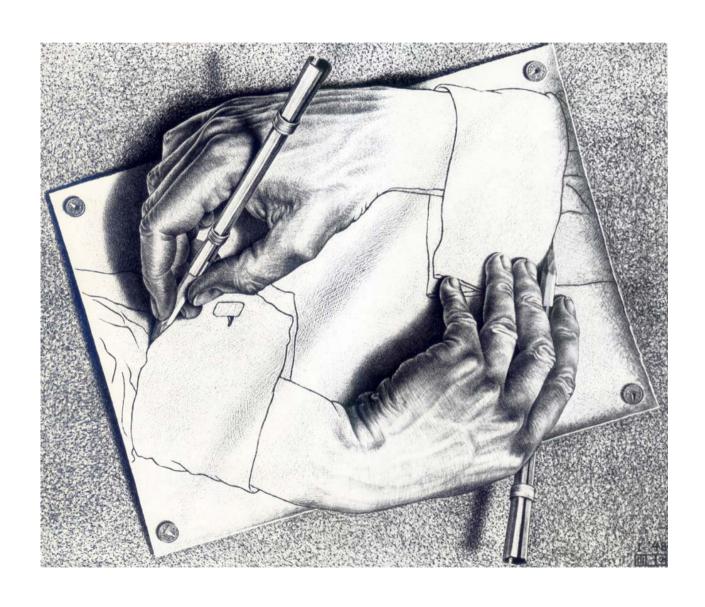
Every even integer greater than 2 can be expressed as the sum of two primes.



Goldbach's Conjecture

```
(define goldbach
  (lambda (n)
    (let ((primes (primes n)))
      (apply append
        (apply append
         (outer-product
           (lambda (x y))
             (if (= n (+ x y))
                 (list (list x y))
                 '()))
           primes
           primes))))))
> (goldbach 98)
((1979)(3167)(3761)(6137)(6731)(7919))
```

Quines



C Quine

```
char data[] = \{35, 105, 110, 99, 108, 117, 100, 101, 32, \dots\};
#include <stdio.h>
main() {
  int i;
  printf("char data[] = {");
  for (i = 0; i < sizeof(data); i++) printf("%d,",data[i]);
  printf("};\n\n");
  for (i = 0; i < sizeof(data); i++) printf("%c", data[i]);
```

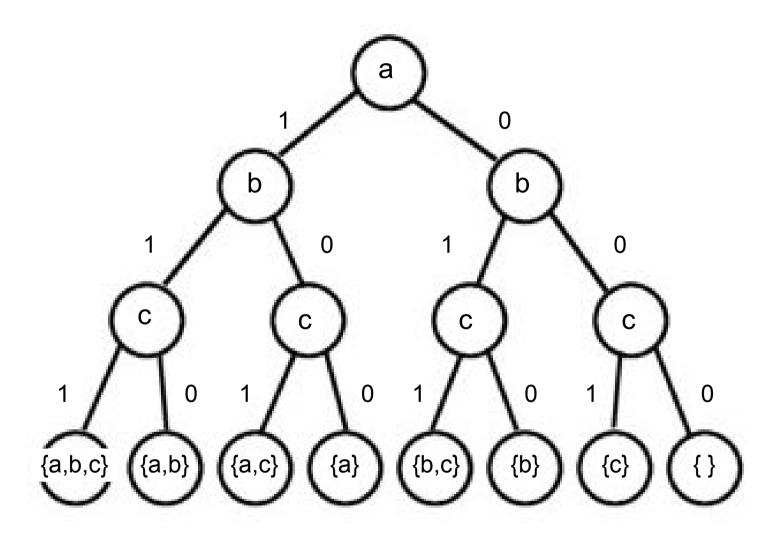
Scheme Quine

```
((lambda (x) (list x (list 'quote x)))
'(lambda (x) (list x (list 'quote x))))
```

APL Quine

 $1 + 22 \rho 11 \rho''' 1 + 22 \rho 11 \rho'''$

Powerset



Results which Grow Fast

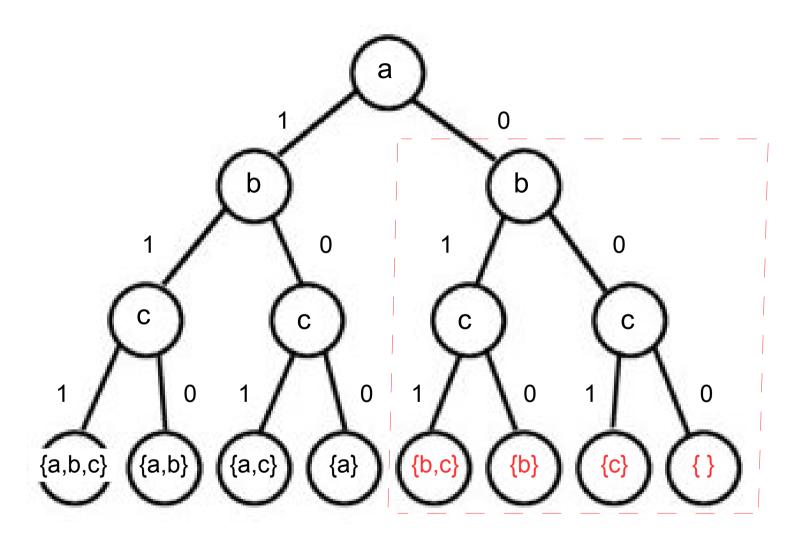
- Space complexity gives a lower bound on time complexity.
- A result of size $O(2^n)$ cannot be computed in less than $O(2^n)$ time!
- To grow this fast, a recursive function must call itself twice in every step.

Opening an Oyster





Powerset



Powerset

```
> (powerset (cdr '(a b c)))
          ((b c) (b) (c) ()
          > (map (lambda (x) (cons 'a x))
                 (powerset (cdr '(a b c))))
          ((a b c) (a b) (a c) (a))
(define powerset
  (lambda (xs)
    (if (null? xs)
        '(())
        (let ((half (powerset (cdr xs))))
          (append (map (lambda (x) (cons (car xs) x))
                        half)
                   half))))
```

Make Change



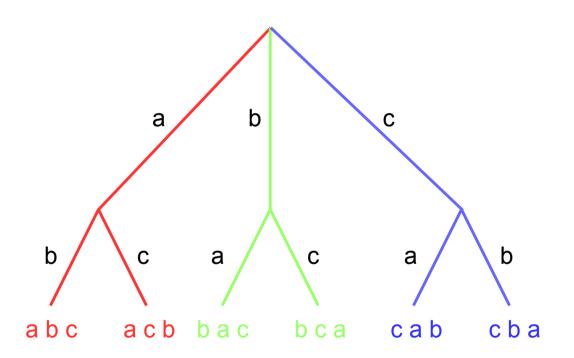
Make Change Explained

```
>(define half (powerset '(25 10 10 5 5 5 1 1 1)))
> half
((25\ 10\ 10\ 5\ 5\ 1\ 1\ 1)
 (25\ 10\ 10\ 5\ 5\ 5\ 1\ 1)
 (25\ 10\ 10\ 5\ 5\ 5\ 1\ 1)
 (1 1)
 (1)
 (1 1)
 (1)
 (1)
 ())
```

Make Change Explained

```
>(map (lambda (ls) (apply + ls)) half)
(63 62 62 61 62 61 61 60 58 57 ... 1 2 1 1 0)
>((select (lambda (x) (= x 57)))
  (map (lambda (ls) (apply + ls)) half)
  half)
((25\ 10\ 10\ 5\ 5\ 1\ 1)
 (25\ 10\ 10\ 5\ 5\ 1\ 1)
 (25 10 10 5 5 1 1))
```

Permutations of {a,b,c}



Permutations

```
> (permutations '(b c))
((b c) (c b))
> (permutations (delete 'a '(a b c)))
((b c) (c b))
>(map (lambda (p) (cons 'a p))
      (permutations (delete 'a '(a b c))))
((a b c) (a c b))
```

Results which Grow Even Faster

- Space complexity gives a lower bound on time complexity.
- A result of size O(n!) cannot be computed in less than O(n!) time!
- To grow this fast, a recursive function must call itself n times in step n.
- It can only do this by mapping itself across a list of size n.

Permutations

```
a b c c c b c a b a
```

```
> (map (lambda (x))
        (map (lambda (p) (cons x p))
              (permutations (delete x '(a b c))))
      '(a b c))
(((a b c) (a c b)) ((b a c) (b c a)) ((c a b) (c b a)))
> (apply append
   (map (lambda (x))
          (map (lambda (p) (cons x p))
                (permutations (delete x '(a b c))))
        '(a b c))
((a b c) (a c b) (b a c) (b c a) (c a b) (c b a))
```

Permutations

