Data Structures and Algorithms Coursework 2 – 100122248

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Question 1:

Pseudocode

return sortedList

Here is the pseudocode for the formDictionary algorithm:

```
Input: An arrayList from a file which has been called from the function
{\tt readWordsFromCSV.}
Output: A list of the sorted words and frequencies sortedList.
Hashmap wordDictionary
ArrayList words, sortedList
integer i
Set wordSet
FOR i=0 to words.size do
      IF wordDictionary CONTAINS KEY(words.get(i))
            wordDictionary ADD (words.get(i) + frequency)
           wordDictionary ADD (words.get(i))
      END IF
END FOR
FOR (word Object in the wordSet)
     ADD word Object to sortedList
END FOR
Sort the sortedList
```

Formal Analysis

Fundamental Operation

The fundamental operation of the code I have found is after the 1st For Loop and is the comparison if the wordDictionary contains an existing word or not.

IF wordDictionary CONTAINS KEY(words.get(i))

Case description

The Worst Case would be that the wordDictionary have to add an extremely high amount of words, with the Best case being that the wordDictionary would only have to add one word. If the text had a very long length, the time taken to complete the algorithm would be very high.

Run Time Complexity Function and Order of Algorithm

I calculated the Run Time Complexity using the for loop before the fundamental operation:

FOR
$$i = 0$$
 to words, size

Using the summation, I used this formula to work out the Run Time Complexity function, using the "words.size" constant to equal 5.

$$O(n) = \sum_{i=1}^{x}$$

$$O(n) = x$$

From this, we can deduce that the Order of the Algorithm would be O(n) as the highest power of n is n^1 .

Question 2:

Pseudocode

add()

```
Input: String key which is the word going to be input into the Trie
Output: A Boolean addSuccessful to see whether the word has been added
successfully
Boolean addSuccessful
TrieNode temp, next
temp <- root
FOR (character in the key) do
      next <- temp.getOffspring(character)</pre>
      IF(next node is null)
            temp.getOffspring(character)
            next <- temp.getOffspring(character)</pre>
      END IF
      temp <- next
END FOR
      set the end of word
IF(temp.endOfWord is TRUE)
      addSuccessful = TRUE
END IF
return addSuccessful
```

Order of Algorithm

contains()

Order of Algorithm

The algorithm's order is O(n) due to one for loop in the algorithm

outputBreadthFirstSearch()

```
Input: A trie of words starting with TrieNode root
Output: A string of stringBuild of the result of the search.
String stringBuild
TrieNode temp
LinkedList queue
temp <- root
ADD temp to queue
WHILE (queue is not Empty) do
      REMOVE temp from queue
      FOR(node in the temp.offspring)
           IF(node is not NULL)
                 add node to stringBuild
                 add node to queue
            END IF
      END FOR
      temp <- head of the queue
END WHILE
return stringBuild
```

Order of Algorithm

outputDepthFirstSearch()

Order of Algorithm

getSubTrie()

```
Input: A string prefix
Output: A trie subTrie which is formed from the prefix
TrieNode temp
Trie subTrie

temp <- root
FOR(each character in the prefix) do
        temp <- temp.getCharacter(character)
END FOR
        subTrie <- new Trie from root temp
return subTrie</pre>
```

Order of Algorithm

The algorithm's order is O(n) due to one for loop in the algorithm

getAllWords()

```
Input: A trie, a StringBuilder
Output: A List of all the words called allWords within the trie
List allWords
StringBuilder string
TrieNode node
APPEND to the string the node
IF(node is the leaf)
     ADD to allWords(string)
FOR (each TrieNode n in the node offspring)
     IF(n is not NULL)
           getAllWords()
     END IF
END FOR
IF(string length is more than 0)
     delete last character of the string
END IF
```

Order of Algorithm

Question 3:

Pseudocode

```
Input: csv files containing words and the prefix queries
Output: a csv file which contains the query result
DictionaryMaker dictionary
Trie trie, subTrie
HashMap freqValues
ArrayList words, wordlist, prefixList, wordPartList, freqList,
sortedList
Scanner scan
String[] part
String wordPart
int freqPart, wordCounter
List allWords
Set freqSet
WHILE(scanner has a newLine)
     add line to the prefixList
END WHILE
FOR(each word in the wordList)
     SET splitter for string
     wordPart <- part[0]</pre>
     freqPart <- part[1]</pre>
     ADD to wordPartList the wordPart
     ADD to freqList the freqPart
END FOR
FOR (each prefix in the prefixList)
     subTrie <- trie.getSubTrie(prefix)</pre>
     allWords <- subTrie.getAllWords()</pre>
     FOR(each word in allWords)
           freqValues.put(value of frequency in the word, word)
     freqSet <- freqValues.entrySet()</pre>
     FOR(each wordObject in the freqSet)
           ADD to sortedList(wordObject)
     END FOR
     Sort the sortedList
     FOR (each string in the sortedList)
           wordCounter += 1
           print string to File
           IF(wordCounter = 5)
                 break out of loop
           END IF
     END FOR
```

clear the freqValues END FOR

Order of Algorithm The algorithm has multiple for loops in them and a nested for loop, so the Order of the algorithm would be $O(n^2)$.