



## Final project - Travelling Salesman Problem

Fanoarii BOYER - Jérémy BEAUGEARD - Arthur GUERINEAU - Adrien LE SAUX

CIR3 - Graph Theory - Leandro MONTERO

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# Chapter 1

## Introduction

In this final graph theory project, we try to solve the Travelling Salesman Problem (TSP) using different algorithms and heuristics. Before we get to the code, it is important to consider how we are going to model the graphs.

In the TSP, we work with undirected complete graphs, which means they are very dense. Adjacency lists are useful for sparse graphs, but in our case it will be more appropriate, in terms of complexity, to use adjacency matrix[1].

To implement the adjacency matrix[2] we will use the C++ Boost[3] library. In this project we seek to compare the performance between different solutions to the same problem. As C++ is a low-level programming language it is very fast and will be suitable for our use. Moreover, the boost library is very well known and widely documented.

## Chapter 2

### Real-life situations

# Chapter 3

## Exact algorithm

### 3.1 Pseudo-code

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### 3.2 Time complexity

### 3.3 Optimal Solution

### 3.4 Execution time and performance

# Chapter 4

## Constructive heuristic

### 4.1 Pseudo-code

```
Input:  $G$  is an undirected complete weighted graph with  $n$  vertices  
Output: The optimal path found and the corresponding distance  
1  $finalPath \leftarrow empty$   
2  $finalDistance \leftarrow \infty$   
3 foreach vertex  $V$  of  $G$  do  
4    $path \leftarrow empty$   
5    $distance \leftarrow 0$   
6    $currentVertex \leftarrow V$   
7   while there are still undiscovered vertices do  
8      $minimumWeight \leftarrow \infty$   
9     foreach adjacent vertex  $V'$  not discovered do  
10       $w \leftarrow \text{weight between } V \text{ and } V'$   
11      if  $w < minimumWeight$  then  
12         $minimumWeight \leftarrow w$   
13         $currentVertex \leftarrow V'$   
14      end  
15    end  
16     $path \leftarrow path + currentVertex$   
17    Mark  $currentVertex$  as discovered  
18     $distance \leftarrow distance + minimumWeight$   
19  end  
20   $returnToStart \leftarrow \text{weight from } nextVertex \text{ to } V$   
21   $distance \leftarrow distance + returnToStart$   
22  if  $distance < finalDistance$  then  
23     $finalDistance \leftarrow distance$   
24     $finalPath \leftarrow path$   
25  end  
26 end  
27 return  $finalPath$  and  $finalDistance$ 
```

4.2 Time complexity

4.3 Optimal Solution

4.4 Execution time and performance

# Chapter 5

## Local search heuristic

### 5.1 Pseudo-code

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### 5.2 Time complexity

### 5.3 Optimal Solution

### 5.4 Execution time and performance



# Chapter 6

## GRASP meta-heuristic

### 6.1 Pseudo-code

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### 6.2 Time complexity

### 6.3 Optimal Solution

### 6.4 Execution time and performance

## Chapter 7

## Conclusion

# Bibliography

- [1] graphs - when are adjacency lists or matrices the better choice? Library Catalog:  
[cs.stackexchange.com](https://cs.stackexchange.com).
- [2] The boost graph library - 1.72.0.
- [3] Boost graph library: Adjacency matrix - 1.72.0.