



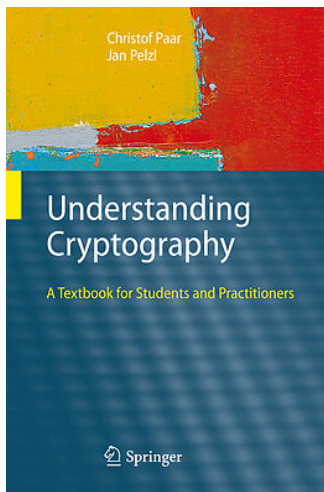
HA NOI UNIVERSITY OF SCIENCE AND TECHNOLOGY  
SCHOOL OF INFORMATION AND COMMUNICATION TECHNOLOGY

# Introduction to Cryptography and Security

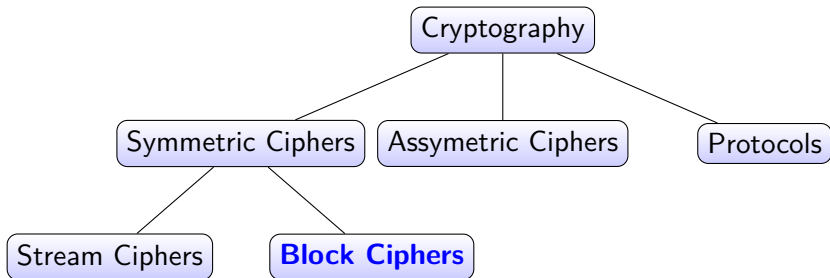
## Modes of Operation

# Textbook

<https://www.crypto-textbook.com>



# Cryptography



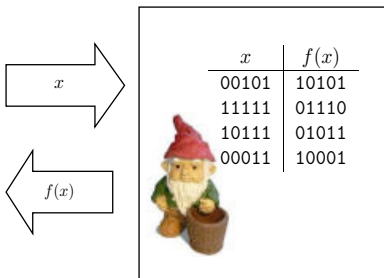
## Ideal block cipher

- In fact, we consider AES or 3DES as an **ideal block cipher**
- That is, for each key  $k$ , the mapping

$$F_k(x) = \text{Enc}(k, x)$$

is an independent random permutation from  $X$  onto itself.

# Random Permutation



upon receiving the  $i$ th query  $x_i \in \mathcal{X}$  from  $\mathcal{A}$  do:  
  if  $x_i = x_j$  for some  $j < i$   
    then  $y_i \leftarrow y_j$   
  else  $y_i \xleftarrow{R} \mathcal{X} \setminus \{y_1, \dots, y_{i-1}\}$   
  send  $y_i$  to  $\mathcal{A}$

# Modes of Operations

**Question:** How to encrypt long messages? (using AES)

**Answer:** There are several ways of encrypting long messages with a block cipher:

- Electronic Code Book (ECB)
- Cipher Block Chaining (CBC)
- Output Feedback (OFB)
- Cipher Output Feedback (CFB)
- Counter Mode (CTR)

**Modes of Operation:** One-time key and **many-time key**.

# Example applications

Many-time key

## File systems

- Same AES key used to encrypt many files.

## IPSec

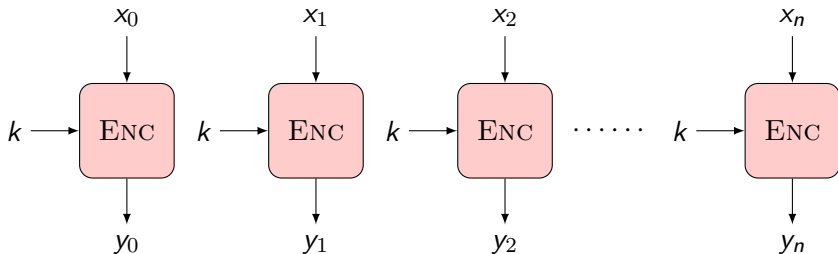
- Same AES key used to encrypt many packets.

# Outline

- 1 Electronic Codebook Mode (ECB)
- 2 Cipher Block Chaining Mode (CBC)
- 3 Stream cipher

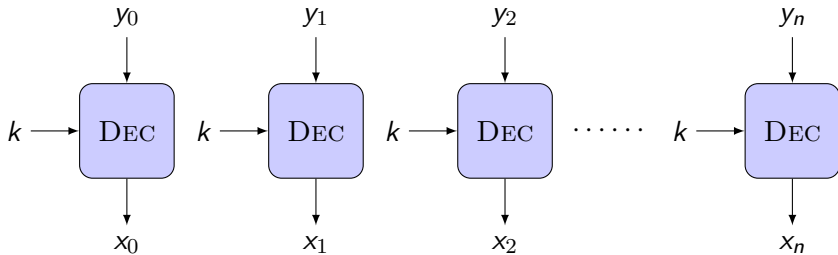


## ECB (Electronic code book)

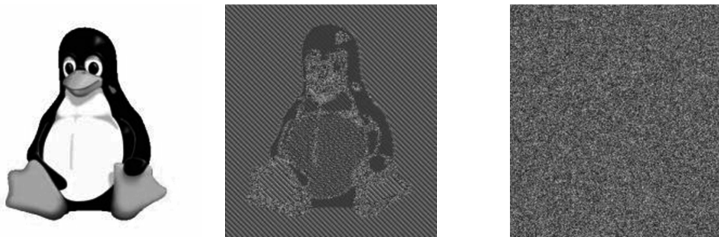


- The messages are partitioned into  $b$ -bit blocks, where  $b$  is the block size.
- If the length of the message is not a multiple of  $b$  bits, it must be padded 10..0 to a multiple of  $b$  bits prior to encryption.
- The **padding operation** is invertible.

## ECB: Decryption



## ECB is not secure



**Figure:** The middle figure is an encryption using ECB mode; the figure on the right is an encryption using a secure mode.

- **Problem:** If  $x_i = x_j$  then  $y_i = y_j$ .
- ECB is secure if the message is random (eg., the keys).

## Example: Electronic Bank Transfer

| Block # | 1                 | 2                    | 3                   | 4                      | 5            |
|---------|-------------------|----------------------|---------------------|------------------------|--------------|
|         | Sending<br>Bank A | Sending<br>Account # | Receiving<br>Bank B | Receiving<br>Account # | Amount<br>\$ |

- ① **Assumption:** Each of the fields has exactly the size of the block cipher width (for example 128 bits)
- ② **Assumption:** The encryption key  $k_{AB}$  between the two banks  $A$  and  $B$  does not change too frequently.

## Attacker Oscar

| Block # | 1                 | 2                    | 3                   | 4                      | 5            |
|---------|-------------------|----------------------|---------------------|------------------------|--------------|
|         | Sending<br>Bank A | Sending<br>Account # | Receiving<br>Bank B | Receiving<br>Account # | Amount<br>\$ |

- 1 He opens one account at bank *A* and one at bank *B*.
- 2 He sends \$1.00 transfers from his account at bank *A* to his account at bank *B* repeatedly.
- 3 He observes the ciphertexts going through the communication network

$$B_1 \parallel B_2 \parallel B_3 \parallel B_4 \parallel B_5$$

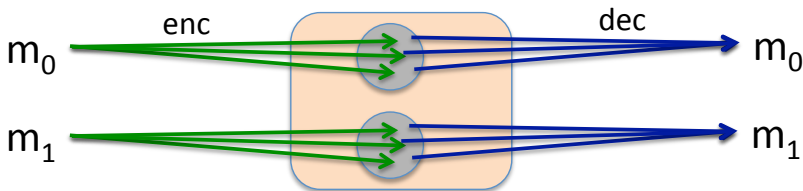
and he stores blocks  $B_1, B_3, B_4$ .

- 4 For all transfers that are made from  $B_1$  to  $B_3$ , he replaces block 4 with  $B_4$ .

# Outline

- 1 Electronic Codebook Mode (ECB)
- 2 Cipher Block Chaining Mode (CBC)
- 3 Stream cipher

## Randomized encryption



- $\text{Enc}(k, m)$  is a randomized algorithm.
- Given the same plaintext message twice, encryption must produce different outputs.
- Ciphertext must be longer than plaintext
- Roughly speaking:  $\text{CT-size} = \text{PT-size} + \text{"\# random bits"}$

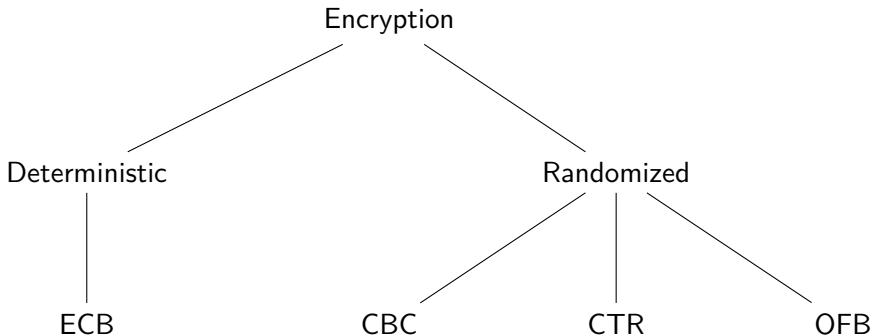
## Exercise

Write the Dec() function for the following Enc().

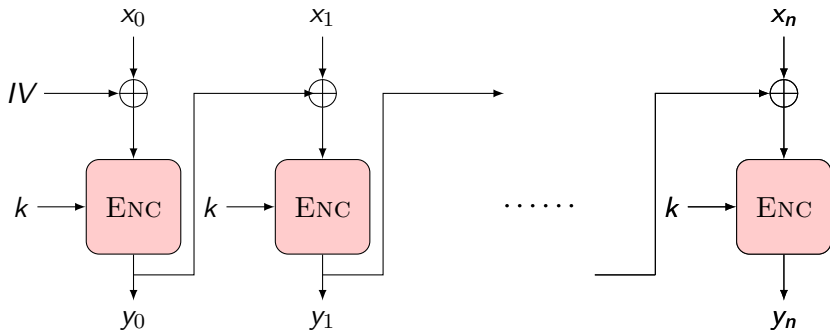
$$\text{Enc}(k, m) := \begin{cases} r = \text{random}() \\ c = \text{AES}(k, r) \oplus m \\ \text{output } (r, c) \end{cases}$$



# Types of Encryption



## Cipher Block Chaining mode (CBC)



**Algorithm.** Choose IV (“initialization value”) randomly, use each  $y_i$  is “IV” for  $M_{i+1}$ . Transmit IV with ciphertext:

$$IV || y_0 || y_1 || \dots || y_n$$

## How to use the $IV$ ?

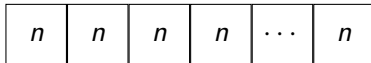
- $IV$  does not need to be kept a secret
- But it must be “nonce” = “number used only once”

### Example

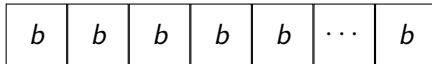
- 1 True random number
- 2 Counter value (must be stored by Alice)
- 3  $ID_A || ID_B || \text{time}$

## A CBC technicality: PKCS5 padding

- The value is the number of bytes that need to be added.
- Padding  $n$  byte, for  $n > 0$

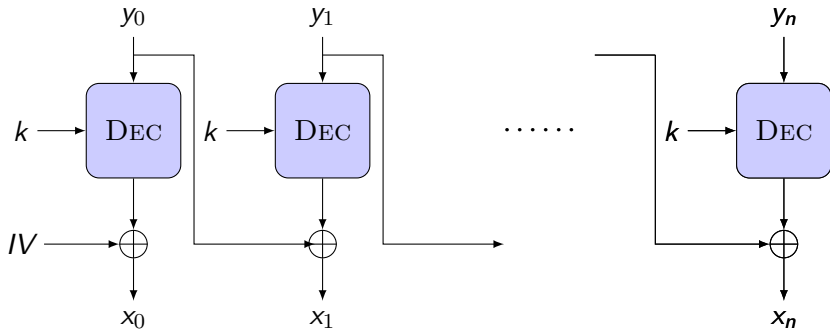


- if no padding is needed, we add a dummy block:



where  $b$  is block size (in bytes).

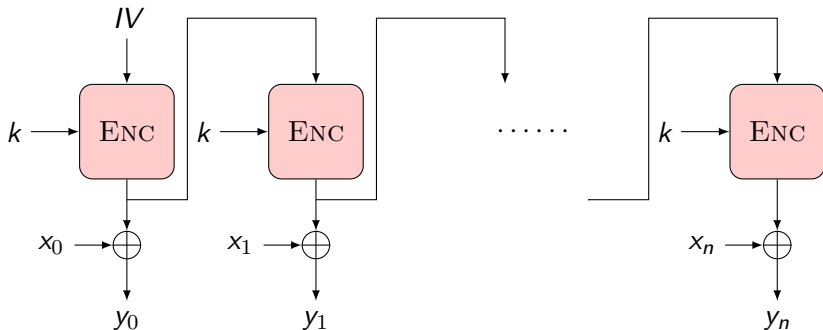
## CBC: Decryption



# Outline

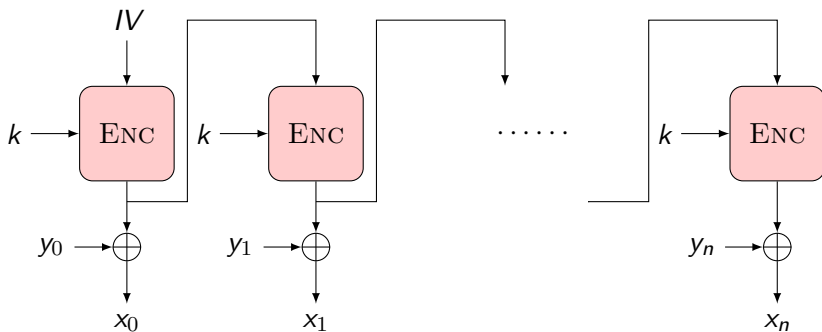
- 1 Electronic Codebook Mode (ECB)
- 2 Cipher Block Chaining Mode (CBC)
- 3 Stream cipher

## Output Feedback Mode (OFB)



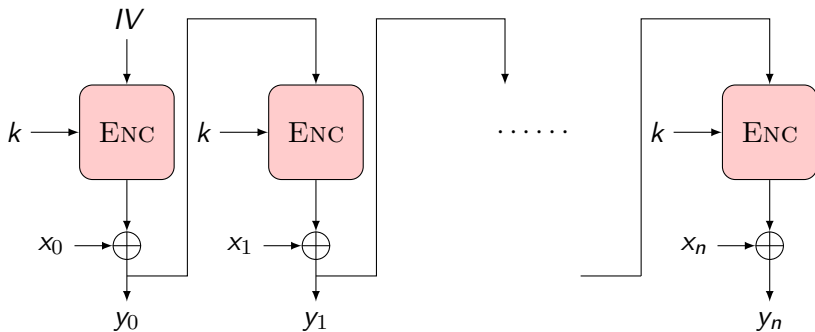
**Algorithm.** Similar to CBC mode. Use a random IV transmitted with the ciphertext.

## OFB: Decryption



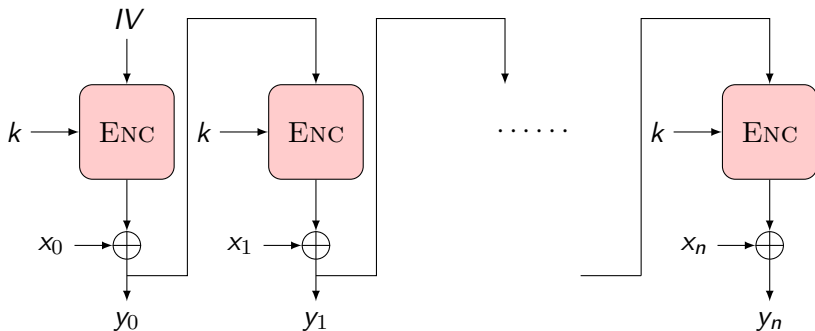


## Cipher Feedback Mode (CFB)



## Exercise

What is the decryption for CFB mode?



## Counter Mode (CTR)

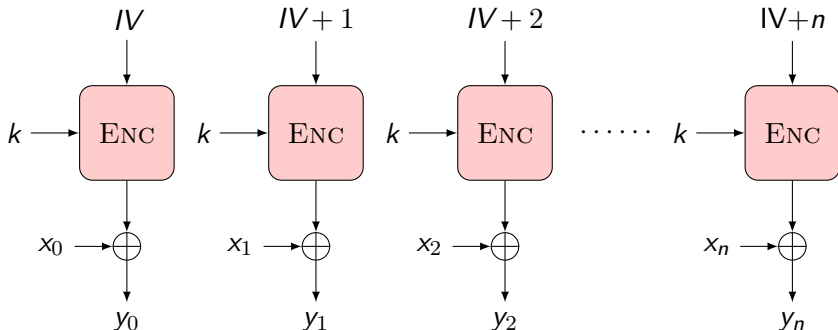
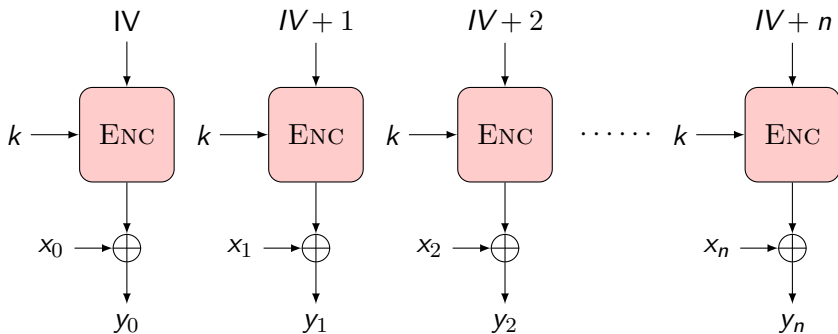


Figure: Use a random IV transmitted with the ciphertext.

## Exercise

What is the decryption for CTR mode?



## Exercise

- Let  $m$  be a message consisting of  $\ell$  AES blocks (say  $\ell = 100$ ).
- Alice encrypts  $m$  using CBC mode and transmits the resulting ciphertext to Bob.
- Due to a network error, ciphertext block number  $\ell/2$  is corrupted during transmission. All other ciphertext blocks are transmitted and received correctly.
- Once Bob decrypts the received ciphertext, how many plaintext blocks will be corrupted?

## Exercise

- Let  $m$  be a message consisting of  $\ell$  AES blocks (say  $\ell = 100$ ).
- Alice encrypts  $m$  using randomized counter mode and transmits the resulting ciphertext to Bob.
- Due to a network error, ciphertext block number  $\ell/2$  is corrupted during transmission. All other ciphertext blocks are transmitted and received correctly.
- Once Bob decrypts the received ciphertext, how many plaintext blocks will be corrupted?



25  
SOICT

VIỆN CÔNG NGHỆ THÔNG TIN VÀ TRUYỀN THÔNG  
SCHOOL OF INFORMATION AND COMMUNICATION TECHNOLOGY

Thank you!



[soict.hust.edu.vn/](http://soict.hust.edu.vn/)



[fb.com/groups/soict](https://fb.com/groups/soict)

