

WIKIPEDIA

二分法

Bisection method

The **bisection method** in mathematics is a ^{寻根的方法} root-finding method that repeatedly bisects an interval and then selects a subinterval in which a root must lie for further processing. It is a very simple and robust method, but it is also relatively slow. Because of this, it is often used to obtain a rough approximation to a solution which is then used as a starting point for more rapidly converging methods.^[1] ^{作为更快收敛方法的起始点} The method is also called the **interval halving method**,^[2] the **binary search method**,^[3] or the **dichotomy method**.^[4]

Contents

The method

- Iteration tasks

- Algorithm

Example: Finding the root of a polynomial

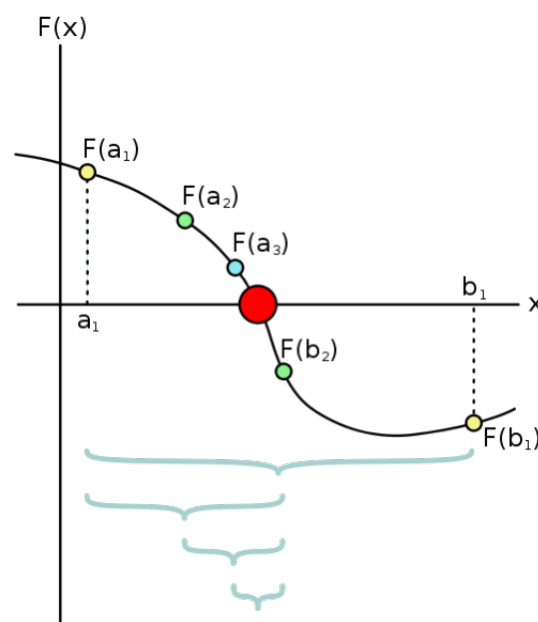
Analysis

See also

References

Further reading

External links



A few steps of the bisection method applied over the starting range $[a_1; b_1]$. The bigger red dot is the root of the function.

The method

The method is applicable for numerically solving the equation $f(x) = 0$ for the real variable x , where f is a continuous function defined on an interval $[a, b]$ and where $f(a)$ and $f(b)$ have opposite signs. In this case a and b are said to bracket a root since, by the intermediate value theorem, the continuous function f must have at least one root in the interval (a, b) .

At each step the method divides the interval in two by computing the midpoint $c = (a+b) / 2$ of the interval and the value of the function $f(c)$ at that point. Unless c is itself a root (which is very unlikely, but possible) there are now only two possibilities: either $f(a)$ and $f(c)$ have opposite signs and bracket a root, or $f(c)$ and $f(b)$ have opposite signs and bracket a root.^[5] The method selects the subinterval that is guaranteed to be a bracket as the new interval to be used in the next step. In this way an interval that contains a zero of f is reduced in width by 50% at each step. The process is continued until the interval is sufficiently small.

Explicitly, if $f(a)$ and $f(c)$ have opposite signs, then the method sets c as the new value for b , and if $f(b)$ and $f(c)$ have opposite signs then the method sets c as the new a . (If $f(c)=0$ then c may be taken as the solution and the process stops.) In both cases, the new $f(a)$ and $f(b)$ have opposite signs, so the method is applicable to this smaller interval.^[6]

Iteration tasks

The input for the method is a continuous function f , an interval $[a, b]$, and the function values $f(a)$ and $f(b)$. The function values are of opposite sign (there is at least one zero crossing within the interval). Each iteration performs these steps:

1. Calculate c , the midpoint of the interval, $c = \frac{a+b}{2}$.
2. Calculate the function value at the midpoint, $f(c)$.
3. If convergence is satisfactory (that is, $c - a$ is sufficiently small, or $|f(c)|$ is sufficiently small), return c and stop iterating.
4. Examine the sign of $f(c)$ and replace either $(a, f(a))$ or $(b, f(b))$ with $(c, f(c))$ so that there is a zero crossing within the new interval.

When implementing the method on a computer, there can be problems with finite precision, so there are often additional convergence tests or limits to the number of iterations. Although f is continuous, finite precision may preclude a function value ever being zero. For example, consider $f(x) = x - \pi$; there will never be a finite representation of x that gives zero. Additionally, the difference between a and b is limited by the floating point precision; i.e., as the difference between a and b decreases, at some point the midpoint of $[a, b]$ will be numerically identical to (within floating point precision of) either a or b .

Algorithm

The method may be written in pseudocode as follows:^[7]

```

INPUT: Function  $f$ , endpoint values  $a, b$ , tolerance  $TOL$ , maximum iterations  $NMAX$ 
CONDITIONS:  $a < b$ , either  $f(a) < 0$  and  $f(b) > 0$  or  $f(a) > 0$  and  $f(b) < 0$ 
OUTPUT: value which differs from a root of  $f(x)=0$  by less than  $TOL$ 

 $N \leftarrow 1$ 
While  $N \leq NMAX$  # limit iterations to prevent infinite loop
     $c \leftarrow (a + b)/2$  # new midpoint
    If  $f(c) = 0$  or  $(b - a)/2 < TOL$  then # solution found
        Output( $c$ )
        Stop
    EndIf
     $N \leftarrow N + 1$  # increment step counter
    If  $\text{sign}(f(c)) = \text{sign}(f(a))$  then  $a \leftarrow c$  else  $b \leftarrow c$  # new interval
EndWhile
Output("Method failed.") # max number of steps exceeded
  
```

Example: Finding the root of a polynomial

Suppose that the bisection method is used to find a root of the polynomial

$$f(x) = x^3 - x - 2.$$

First, two numbers a and b have to be found such that $f(a)$ and $f(b)$ have opposite signs. For the above function, $a = 1$ and $b = 2$ satisfy this criterion, as

$$f(1) = (1)^3 - (1) - 2 = -2$$

and

$$f(2) = (2)^3 - (2) - 2 = +4.$$

Because the function is continuous, there must be a root within the interval [1, 2].

In the first iteration, the end points of the interval which brackets the root are $a_1 = 1$ and $b_1 = 2$, so the midpoint is

$$c_1 = \frac{2 + 1}{2} = 1.5$$

The function value at the midpoint is $f(c_1) = (1.5)^3 - (1.5) - 2 = -0.125$. Because $f(c_1)$ is negative, $a = 1$ is replaced with $a = 1.5$ for the next iteration to ensure that $f(a)$ and $f(b)$ have opposite signs. As this continues, the interval between a and b will become increasingly smaller, converging on the root of the function. See this happen in the table below.

Iteration	a_n	b_n	c_n	$f(c_n)$
1	1	2	1.5	−0.125
2	1.5	2	1.75	1.6093750
3	1.5	1.75	1.625	0.6660156
4	1.5	1.625	1.5625	0.2521973
5	1.5	1.5625	1.5312500	0.0591125
6	1.5	1.5312500	1.5156250	−0.0340538
7	1.5156250	1.5312500	1.5234375	0.0122504
8	1.5156250	1.5234375	1.5195313	−0.0109712
9	1.5195313	1.5234375	1.5214844	0.0006222
10	1.5195313	1.5214844	1.5205078	−0.0051789
11	1.5205078	1.5214844	1.5209961	−0.0022794
12	1.5209961	1.5214844	1.5212402	−0.0008289
13	1.5212402	1.5214844	1.5213623	−0.0001034
14	1.5213623	1.5214844	1.5214233	0.0002594
15	1.5213623	1.5214233	1.5213928	0.0000780

After 13 iterations, it becomes apparent that there is a convergence to about 1.521: a root for the polynomial.

Analysis

The method is guaranteed to converge to a root of f if f is a continuous function on the interval $[a, b]$ and $f(a)$ and $f(b)$ have opposite signs. The absolute error is halved at each step so the method converges linearly, which is comparatively slow.

Specifically, if $c_1 = \frac{a+b}{2}$ is the midpoint of the initial interval, and c_n is the midpoint of the interval in the n th step, then the difference between c_n and a solution c is bounded by^[8]

$$|c_n - c| \leq \frac{|b - a|}{2^n}.$$

This formula can be used to determine in advance the number of iterations that the bisection method would need to converge to a root to within a certain tolerance. The number of iterations needed, n , to achieve a given error (or tolerance), ϵ , is given by: $n = \log_2\left(\frac{\epsilon_0}{\epsilon}\right) = \frac{\log \epsilon_0 - \log \epsilon}{\log 2}$,

where $\epsilon_0 = \text{initial bracket size} = b - a$.

Therefore, the linear convergence is expressed by $\epsilon_{n+1} = \text{constant} \times \epsilon_n^m$, $m = 1$.

See also

- Binary search algorithm
- Lehmer–Schur algorithm, generalization of the bisection method in the complex plane
- Nested intervals

References

- ↑ Burden & Faires 1985, p. 31
 - ↑ "Archived copy" (<https://web.archive.org/web/20130519092250/http://siber.cankaya.edu.tr/NumericalComputations/ceng375/node32.html>). Archived from the original (<http://siber.cankaya.edu.tr/NumericalComputations/ceng375/node32.html>) on 2013-05-19. Retrieved 2013-11-07.
 - ↑ Burden & Faires 1985, p. 28
 - ↑ "Dichotomy method - Encyclopedia of Mathematics" (https://www.encyclopediaofmath.org/index.php/Dichotomy_method). *www.encyclopediaofmath.org*. Retrieved 2015-12-21.
 - ↑ If the function has the same sign at the endpoints of an interval, the endpoints may or may not bracket roots of the function.
 - ↑ Burden & Faires 1985, p. 28 for section
 - ↑ Burden & Faires 1985, p. 29. This version recomputes the function values at each iteration rather than carrying them to the next iterations.
 - ↑ Burden & Faires 1985, p. 31, Theorem 2.1
- Burden, Richard L.; Faires, J. Douglas (1985), "2.1 The Bisection Algorithm", *Numerical Analysis* (3rd ed.), PWS Publishers, ISBN 0-87150-857-5

Further reading

- Corliss, George (1977), "Which root does the bisection algorithm find?", *SIAM Review*, **19** (2): 325–327, doi:10.1137/1019044 (<https://doi.org/10.1137/1019044>), ISSN 1095-7200 (<https://www.worldcat.org/issn/1095-7200>)
- Kaw, Autar; Kalu, Egwu (2008), *Numerical Methods with Applications* (https://web.archive.org/web/20090413123941/http://numericalmethods.eng.usf.edu/topics/textbook_index.html) (1st ed.), archived from the original (http://numericalmethods.eng.usf.edu/topics/textbook_index.html) on 2009-04-13

External links

- Weisstein, Eric W. "Bisection" (<http://mathworld.wolfram.com/Bisection.html>). *MathWorld*.
 - Bisection Method (https://web.archive.org/web/20060901073129/http://numericalmethods.eng.usf.edu/topics/bisection_method.html) Notes, PPT, Mathcad, Maple, Matlab, Mathematica from Holistic Numerical Methods Institute (<https://web.archive.org/web/20060906070428/http://numericalmethods.eng.usf.edu/>)
-

Retrieved from "https://en.wikipedia.org/w/index.php?title=Bisection_method&oldid=842231149"

This page was last edited on 21 May 2018, at 03:10.

Text is available under the [Creative Commons Attribution-ShareAlike License](#); additional terms may apply. By using this site, you agree to the [Terms of Use](#) and [Privacy Policy](#). Wikipedia® is a registered trademark of the [Wikimedia Foundation, Inc.](#), a non-profit organization.