

Q5

$$\max. Z = x_1 + 2x_2$$

$$\text{s.t. } x_1 - x_2 \geq -2$$

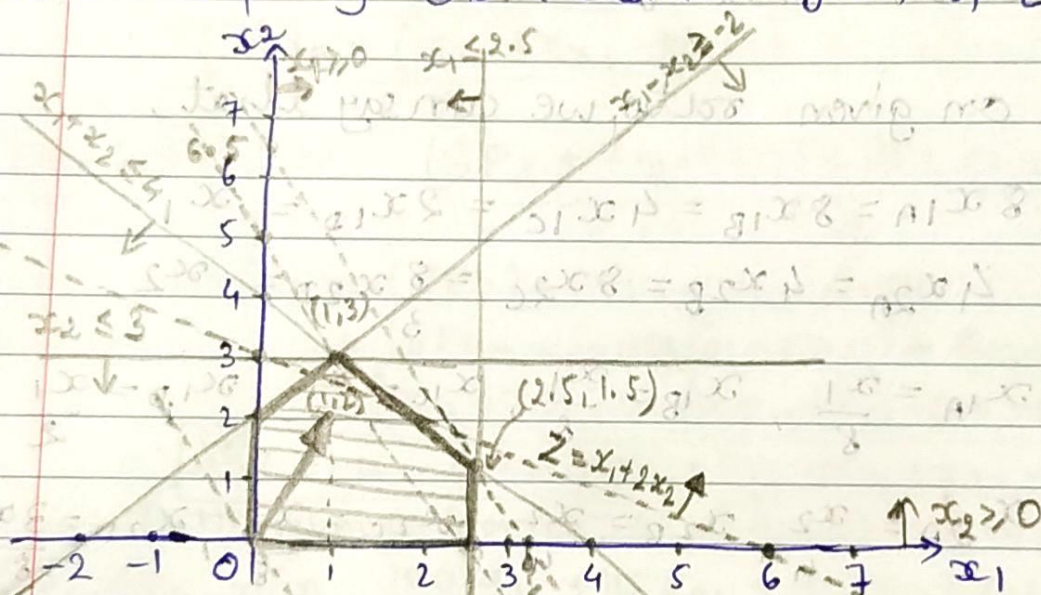
$$x_1 + x_2 \leq 4$$

$$x_1 \leq 2.5$$

$$x_2 \leq 3$$

$$x_1, x_2 \geq 0$$

(a) Graphing the feasible region of LP:



→ As we can say from the graph, the feasible region is bounded, as shaded with hatch lines.

(b) Yes,  $x_2 \leq 3$  is a redundant constraint, as it does not change feasible region.

(c) Solving LP using the graphical method:

$$\text{Obj. fun. } Z = x_1 + 2x_2$$

(max)

here to plot  $Z$  in graph we take,

$$Z = 6 = x_1 + 2x_2$$

∴ we can plot it as shown in graph.



∴ gradient of  $Z \Rightarrow \left( \frac{\partial Z}{\partial x_1}, \frac{\partial Z}{\partial x_2} \right) = (1, 2) \Rightarrow$  It gives dir<sup>n</sup> of increasing  $Z$  value.

∴ We can say that, corner point  $(1, 3)$  will give us max<sup>m</sup>  $Z$  value from this feasible region points.

$$\therefore \boxed{\max Z = x_1 + 2x_2 = 1 + 2(3) = 7}$$

note: If we check,  $Z$  value for all corner-points then also, we can cross-verify that only  $(1, 3)$  point give us maximum  $Z$  value.

note, that we found corner point  $(1, 3)$  by solving  $x_1 - x_2 = (-2)$  and  $x_1 + x_2 = 4$  for their intersection point.

(d) No, there is only one optimal solution. The reason is as we can see from graph, optimal point is where constraints intersect and objective fun<sup>n</sup> touches that point.

→ note that we could have multiple optimal sol<sup>n</sup> only when objective fun<sup>n</sup> is overlapping with one of the constraint, which is not the case here ∴ we have unique corner-point sol<sup>n</sup> here.

(e) Suppose we add the constraint  $2x_1 + x_2 \geq \alpha$  to (LP).

(i) for  $\alpha \leq 0$  constraint will be redundant.

∵ as seen in graph it will not change feasible region in that case.

(ii) for  $\alpha > 5$  and  $\alpha \leq 6.5$ , optimal sol<sup>n</sup> found above as point  $(1, 3)$  will no longer be optimal but some other point will become optimal solution, as for  $5 < \alpha \leq 6.5$  will make  $2x_1 + x_2 \geq \alpha$  constraint go past  $(1, 3)$  corner point and  $(1, 3)$  point will no longer be in the feasible region.

(iii) for  $\alpha > 6.5$ , there will be no feasible region and problem will become infeasible.

note: we found, 6.5 value from point of intersection of  $(2.5, 1.5)$  for  $\alpha$  which gives  $2x_1 + x_2 = 5 + 1.5 = 6.5$ .  
and we found, 5 value from point  $(1, 3)$ :  $2x_1 + x_2 = 2(1) + 3 = 5$ .